

Lattices

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Mathematical Foundations of Cryptography – WT 2020/21

Outline

Vector Spaces

Definition and Properties

- Fundamental Domain
- Volume

Short Vectors in Lattices

- Computational Problems
- Minkowski's and Hermite's Theorem

Lattice Reduction Algorithms

- Babai's Algorithm
- The Two-Dimensional Case

Literature

The slides are based on the following sources

- **An Introduction to Mathematical Cryptography**, Hoffstein, Jeffrey, Pipher, Jill, Silverman, J.H.
- **A Decade of Lattice Cryptography**, Chris Peikert
- **The LLL Algorithm**, Phong Q. Nguyen, Brigitte Vallée (Eds.)

Many graphics are based on graphics from Maria Eichlseder.

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3/4 resp. 2/3 candidates for NIST Post-Quantum Cryptography Standardization are lattice-based (in the category PKE resp. signature).

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- Strong security guarantees from worst-case hardness.
- Construction of versatile and powerful cryptographic objects
 - Fully Homomorphic Encryption
 - Attribute-Based Encryption

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- A **linear combination** of the vectors v_1, \dots, v_k is any vector of the form

$$w = \alpha_1 v_1 + \dots + \alpha_k v_k, \text{ with } \alpha_1, \dots, \alpha_k \in \mathbb{R}.$$

The collection of all such linear combinations is called the **span** of $\{v_1, \dots, v_k\}$.
e.g.

$$w = 2 \cdot \begin{pmatrix} 1 \\ 2 \\ 3 \end{pmatrix} + 3 \cdot \begin{pmatrix} -3 \\ 4 \\ 5 \end{pmatrix}$$

Vector Spaces

- A set of vectors $v_1, \dots, v_k \in V$ is **linearly independent** if

$$\alpha_1 v_1 + \dots + \alpha_k v_k = \mathbf{0} \Rightarrow \alpha_1 = \dots = \alpha_k = 0.$$

What about

$$\begin{pmatrix} 2 \\ 6 \\ -5 \end{pmatrix}, \begin{pmatrix} 3 \\ 9 \\ -7.5 \end{pmatrix}, \text{ and } \begin{pmatrix} 2 \\ -7 \\ 5 \end{pmatrix}, \begin{pmatrix} 3 \\ -1 \\ 1 \end{pmatrix}$$

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- Any two bases for V have the same number elements.

Length and Angle 1

- The **dot product** of $v = (x_1, \dots, x_m)$, $w = (y_1, \dots, y_m) \in V$ is the quantity

$$v \cdot w = x_1 y_1 + \dots + x_m y_m.$$

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- The **length**, or **Euclidean norm**, of v is the quantity

$$\|v\| = \sqrt{x_1^2 + \dots + x_m^2}.$$

$$v = \begin{pmatrix} 2 \\ 1 \\ 2 \end{pmatrix}; \|v\| = ?$$

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- Let α be the **angle** between v and w , then

$$v \cdot w = \|v\| \|w\| \cos(\alpha).$$

Gram Matrix

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Example: Let $v_1 = (2, 3)$, $v_2 = (1, 4)$.

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$$\text{vol}(v_1, v_2) = \sqrt{\det G} = \sqrt{25} = 5$$

Gram-Schmidt Algorithm

Theorem (Gram-Schmidt Algorithm)

Let v_1, \dots, v_n be a basis for a vector space $V \subset \mathbb{R}^m$. The following algorithm creates an orthogonal basis v_1^*, \dots, v_n^* for V :

$$v_1^* \leftarrow v_1$$

for $i = 2..n$ **do**

for $j = 1..i - 1$

$$\mu_{i,j} \leftarrow \frac{v_i \cdot v_j^*}{\|v_j^*\|^2}$$

$$v_i^* = v_i - \sum_{j=1}^{i-1} \mu_{i,j} v_j^*$$

Definition and Properties

Lattices

Definition (Lattice)

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- discrete

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Lattice: Example

In other words, let $v_1, \dots, v_n \in \mathbb{R}^n$ be a set of linearly independent vectors. The lattice generated by v_1, \dots, v_n is the set of linear combinations of v_1, \dots, v_n with coefficients in \mathbb{Z} ,

$$L = \{a_1v_1 + \dots + a_nv_n : a_1, \dots, a_n \in \mathbb{Z}\}.$$

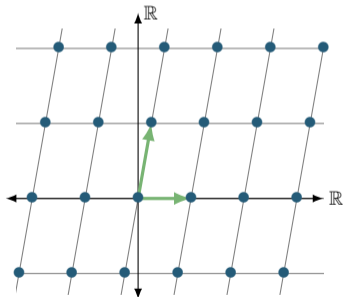
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Example:

$$v_1 = \begin{pmatrix} 1 \\ 0 \end{pmatrix}, v_2 = \begin{pmatrix} 1/4 \\ \sqrt{2} \end{pmatrix}$$



Fundamental Domains

Definition (Fundamental Domain)

Let L be a lattice of dimension n and let v_1, \dots, v_n be a basis for L . The **fundamental domain** is the set

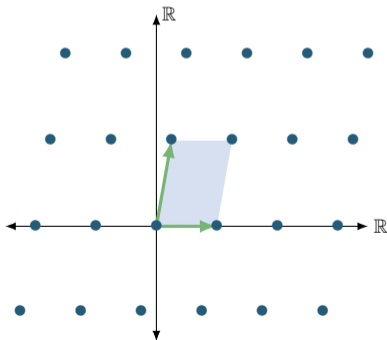
$$F = [0, 1)v_1 + \dots + [0, 1)v_n.$$

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First, compute Gram matrix:

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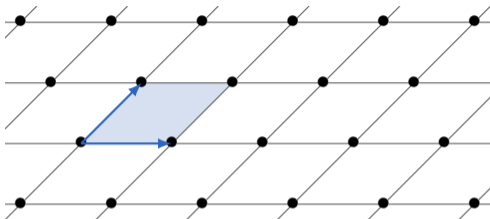
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Therefore,

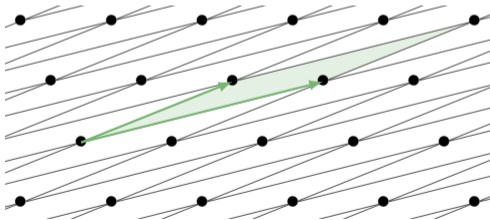
$$\text{vol}(L) = \sqrt{\det G} = \sqrt{2}$$

Same Lattice?

$$\mathbf{v}_1 = \begin{pmatrix} 3 \\ 0 \end{pmatrix}, \mathbf{v}_2 = \begin{pmatrix} 2 \\ 2 \end{pmatrix}$$



$$\mathbf{v}'_1 = \begin{pmatrix} 8 \\ 2 \end{pmatrix}, \mathbf{v}'_2 = \begin{pmatrix} 5 \\ 2 \end{pmatrix}$$



Volume: Task

Task: Compute the volumes V resp. V' of the fundamental domains corresponding to v_1, v_2 respectively v'_1, v'_2 .

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$$G = \begin{pmatrix} 3 & 0 \\ 2 & 2 \end{pmatrix} \begin{pmatrix} 3 & 2 \\ 0 & 2 \end{pmatrix} = \begin{pmatrix} 9 & 6 \\ 6 & 8 \end{pmatrix}.$$

$$G' = \begin{pmatrix} 8 & 2 \\ 5 & 2 \end{pmatrix} \begin{pmatrix} 8 & 5 \\ 2 & 2 \end{pmatrix} = \begin{pmatrix} 68 & 44 \\ 44 & 29 \end{pmatrix}.$$

Therefore $V = \sqrt{G} = \sqrt{36} = 6 = \sqrt{36} = \sqrt{G'} = V'$.

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Proposition

Every fundamental domain for a given lattice L has the same volume.

Short Vectors in Lattices

Computational Problems

$\lambda_1(L)$... length of shortest nonzero vector in L .

- **Shortest Vector Problem (SVP):** Find a shortest nonzero vector v in L , i.e. $\|v\| = \lambda_1(L)$.
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Example: Given the lattice generated by v_1, v_2

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and given the vector $w = (2, 3)^T$. What is a shortest nonzero vector of L ? Which vector is closest to w ?

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$$\begin{pmatrix} -1 \\ 2 \end{pmatrix} \text{ and } \begin{pmatrix} 2 \\ 2 \end{pmatrix}$$

How long is the shortest vector?

Theorem (Minkowski's Theorem)

Let $L \subset \mathbb{R}^n$ be a lattice of dimension n . Let $S \subset \mathbb{R}^n$ be convex, closed and symmetric. Suppose that $\text{vol}(S) \geq 2^n \text{vol}(L)$, then

$$S \cap L \neq \{0\}.$$

S... hypercube in \mathbb{R}^n centered at 0 with length $2 \text{vol}(L)^{1/n}$, then $\text{vol}(S) = 2^n \text{vol}(L)$. Applying Minkowski's theorem leads to:

Corollary (Hermite's Theorem)

Every lattice L of dimension n contains a nonzero $v \in L$ satisfying

$$\|v\| \leq \sqrt{n} \text{vol}(L)^{\frac{1}{n}}.$$

Lattice Reduction Algorithms

Babai's Closest Vertex Algorithm

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Orthogonality Defects

Definition (Hadamard Ratio)

We define the **Hadamard ratio** of the basis $B = \{v_1, \dots, v_n\}$ to be the quantity

$$H(B) = \left(\frac{\text{vol}(L)}{\|v_1\| \cdots \|v_n\|} \right)^{\frac{1}{n}} \in (0, 1].$$

(the closer to 1, the more orthogonal)

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$$H(B) = \left(\frac{6}{\sqrt{9}\sqrt{8}} \right)^{\frac{1}{2}} \approx 0.84.$$

Recap

Lattice: Basis, Fundamental Domain, Volume

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Reduction: Babai's Algorithm

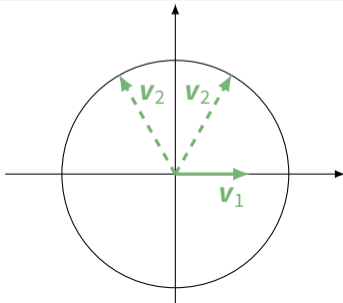
Lagrange-Reduced

Definition (Lagrange-reduced)

Let L be a two-dimensional lattice. A basis (v_1, v_2) of L is said to be **Lagrange-reduced** if and only if

$$\|v_1\| \leq \|v_2\| \quad \text{and} \quad |v_1 \cdot v_2| \leq \frac{\|v_1\|^2}{2}.$$

Optimal: $\lambda_1(L) = \|v_1\|$



Lagrange's Reduction Algorithm

Input: A basis (u, v) of a 2-dimensional lattice L .

Output: A Lagrange-reduced basis of L .

Lagrange's Reduction Algorithm

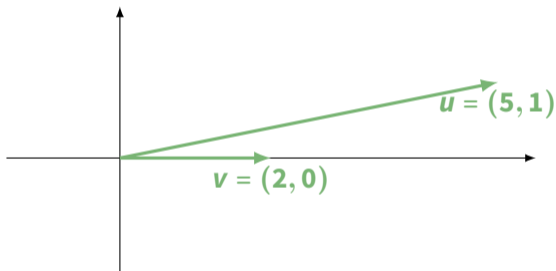
Input: A basis (u, v) of a 2-dimensional lattice L .

Output: A Lagrange-reduced basis of L .

```
if  $\|u\| > \|v\|$  then  
    swap  $u$  and  $v$   
while  $\|v\| > \|u\|$  do  
     $r \leftarrow u - qv$  where  $q = \left\lfloor \frac{u \cdot v}{\|v\|^2} \right\rfloor$   
     $u \leftarrow v$   
     $v \leftarrow r$   
return  $(u, v)$ 
```

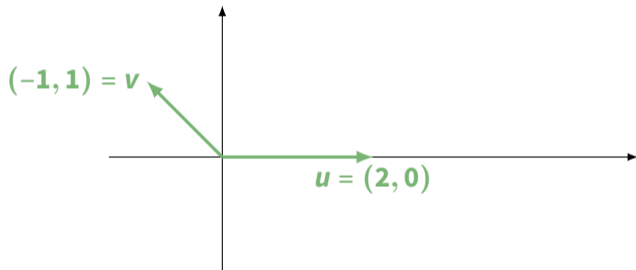
Lagrange Reduction: Example

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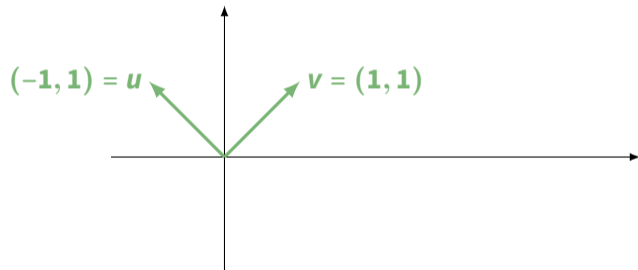
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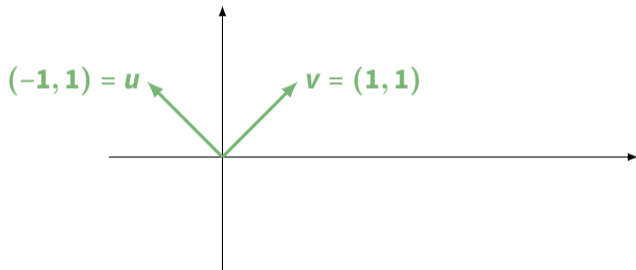
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Task: Solve SVP for the lattice generated by

$$v_1 = (66586820, 65354729)^T, v_2 = (6513996, 6393464)^T.$$