



Chapter 8 - Software-based Power Attacks

Attacking CPUs with Power Side Channels from Software

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3rd April 2025

CPU Power Management

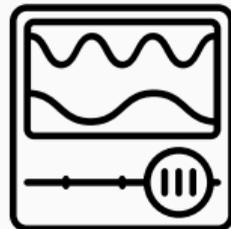
- CPU power management is **complex**
- In order to **save power**, you can ...



Shut down resources



Reduce **voltage**



Reduce **frequency**



- Therefore, the CPU requires:
 - Thermal Management
 - Platform Power Limiting
 - Power/Performance Budgeting
- Domains: PKG, CORE, MC
- **Intel Running Average Power Limit (RAPL)** provides:



power limiting



energy reading



- **Linux:** accessed via **powercap** framework
`/sys/devices/virtual/powercap/intel-rapl`
- **macOS** and **Windows:** Intel driver needs to be installed

Intel RAPL: Properties



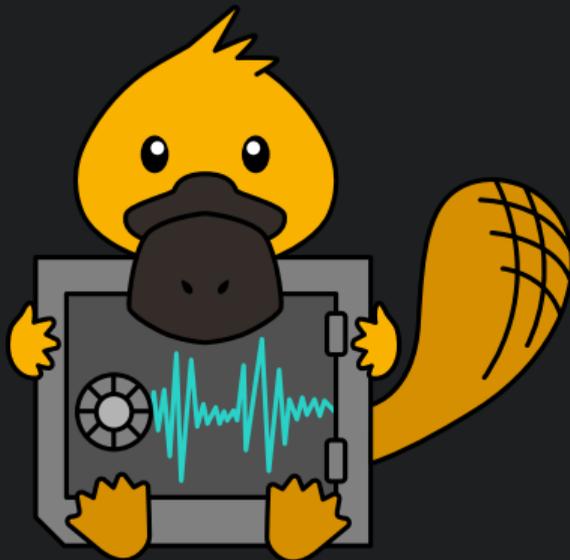
Unprivileged power meter



No physical access

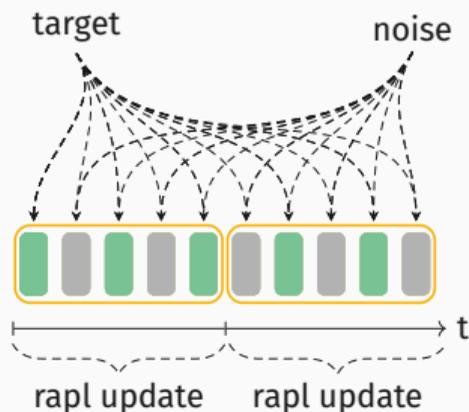


Low refresh rate

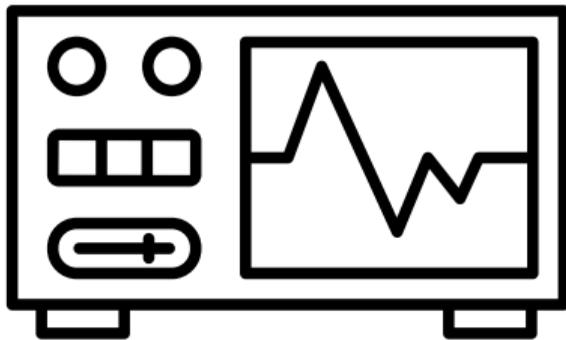


Platypus Attack

RAPL: Measurement Techniques



- Measure an **instruction** by
 - executing it **once**
 - executing it **repeatedly**
 - padding it with **known** instructions
 - **reissue** the instruction after an interrupt



What can we do with this?

Distinguishing Instructions

- Measure the **energy consumption** of **different instructions**

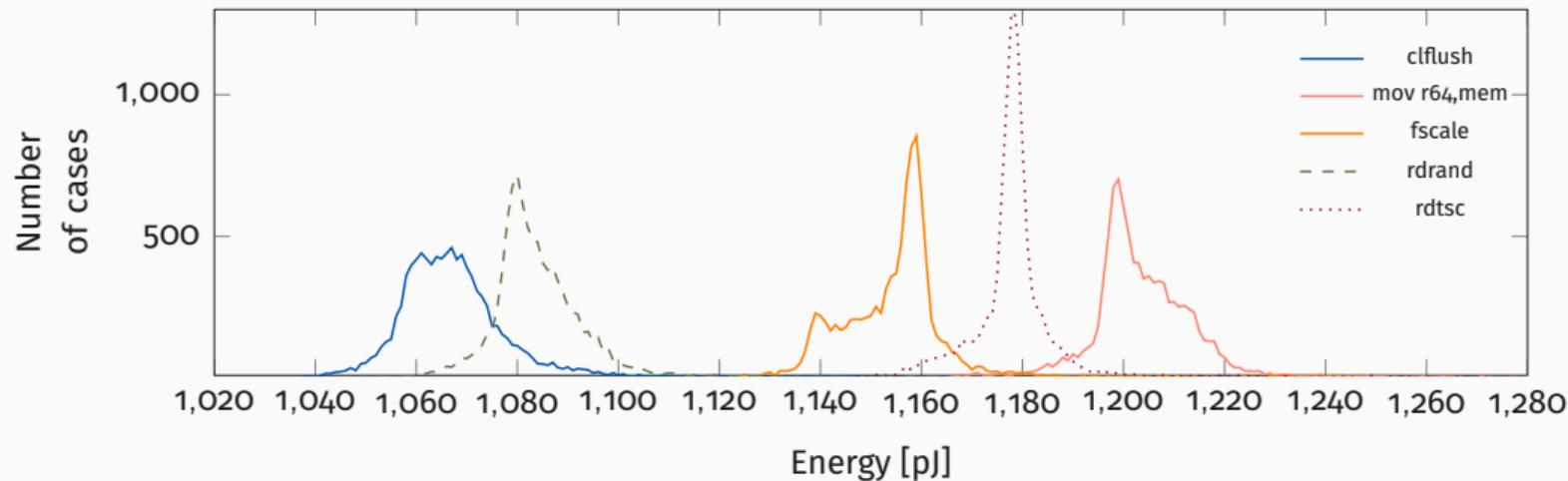
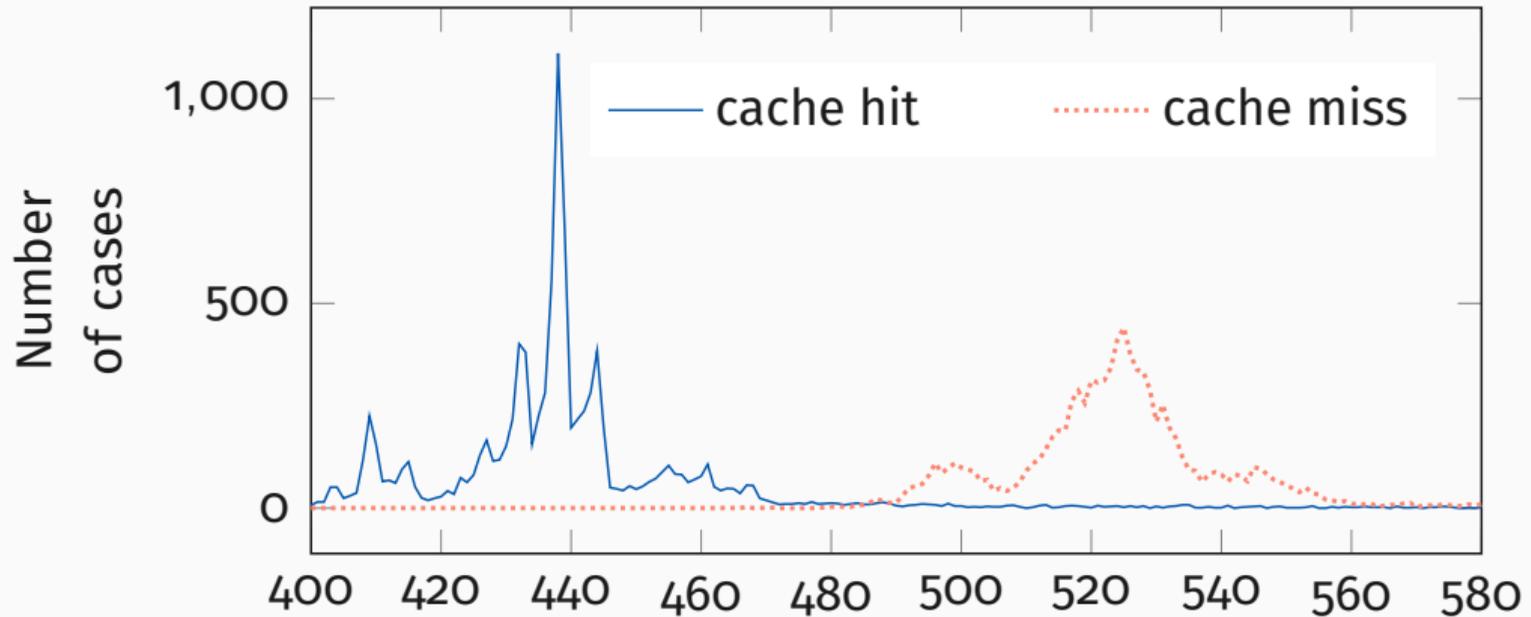


Figure 1: A histogram of the power consumption of various instructions on the i7-6700K (desktop) system.

Distinguishing Load Targets

- Measure the **energy consumption** of **different load targets**



Distinguishing Operands

- Measure the **energy consumption** of **different operands**

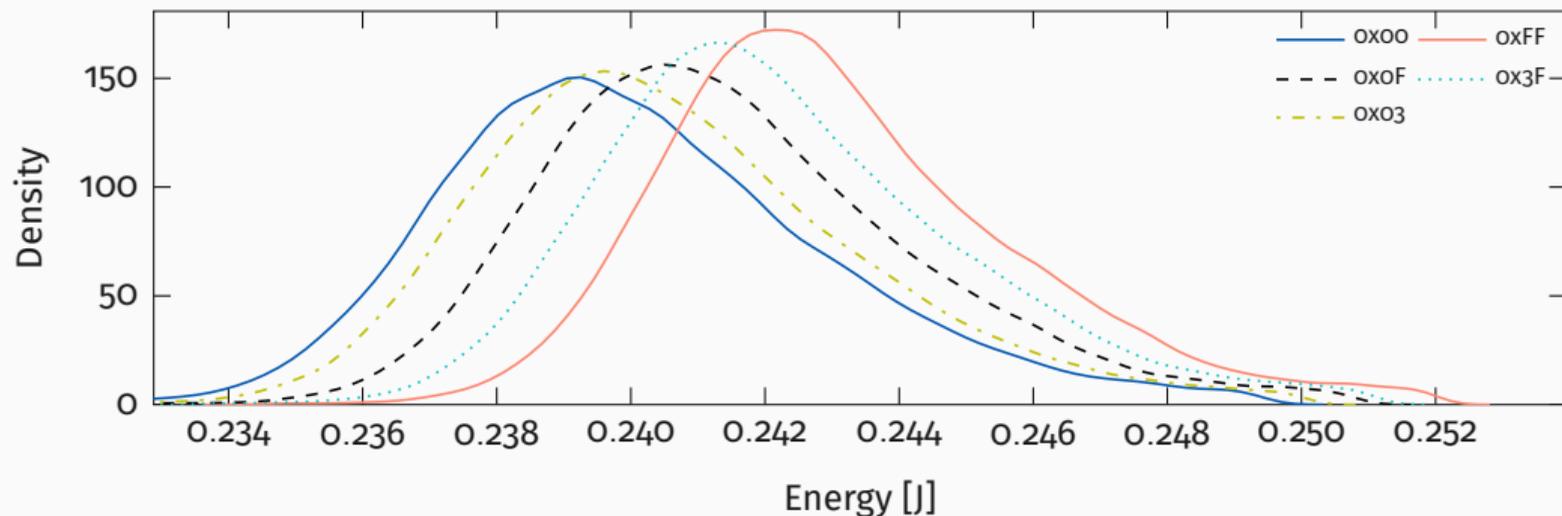


Figure 3: Measured energy consumption of the `imul` instruction with one operand fixed to 8 and the other varying in its Hamming weight.



Let's exploit this!



- **Hidden** communication channel
- Leveraging the **power** side channel



- 2 Processes, Sender and Receiver
 - **Send a 1:** Perform energy-consuming instructions
 - **Send a 0:** Idle
 - Receiver measures **power consumption**
- **Deduces transmitted bit**

Covert Channel

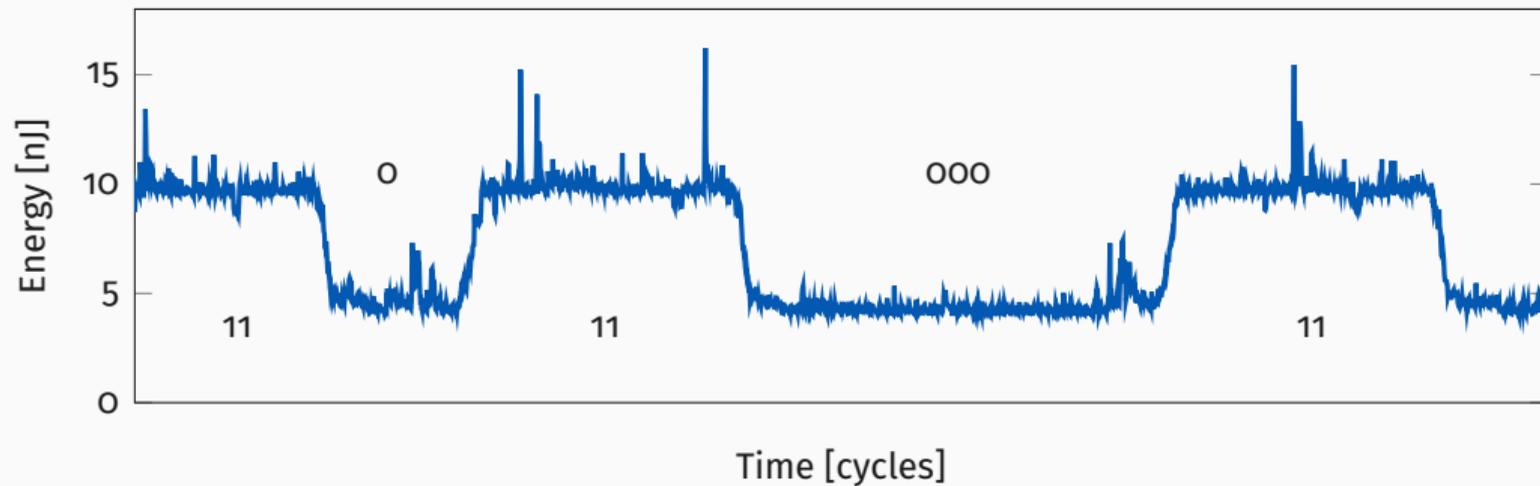


Figure 4: Transmission of bits 1101100011 using the time-less covert channel.



- Kernel Address Space Layout Randomization (KASLR)
- **Exploit energy consumption differences** between
 - Mapped addresses
 - Unmapped addresses
- **Valid address translations** are cached in the **TLB**

Breaking KASLR

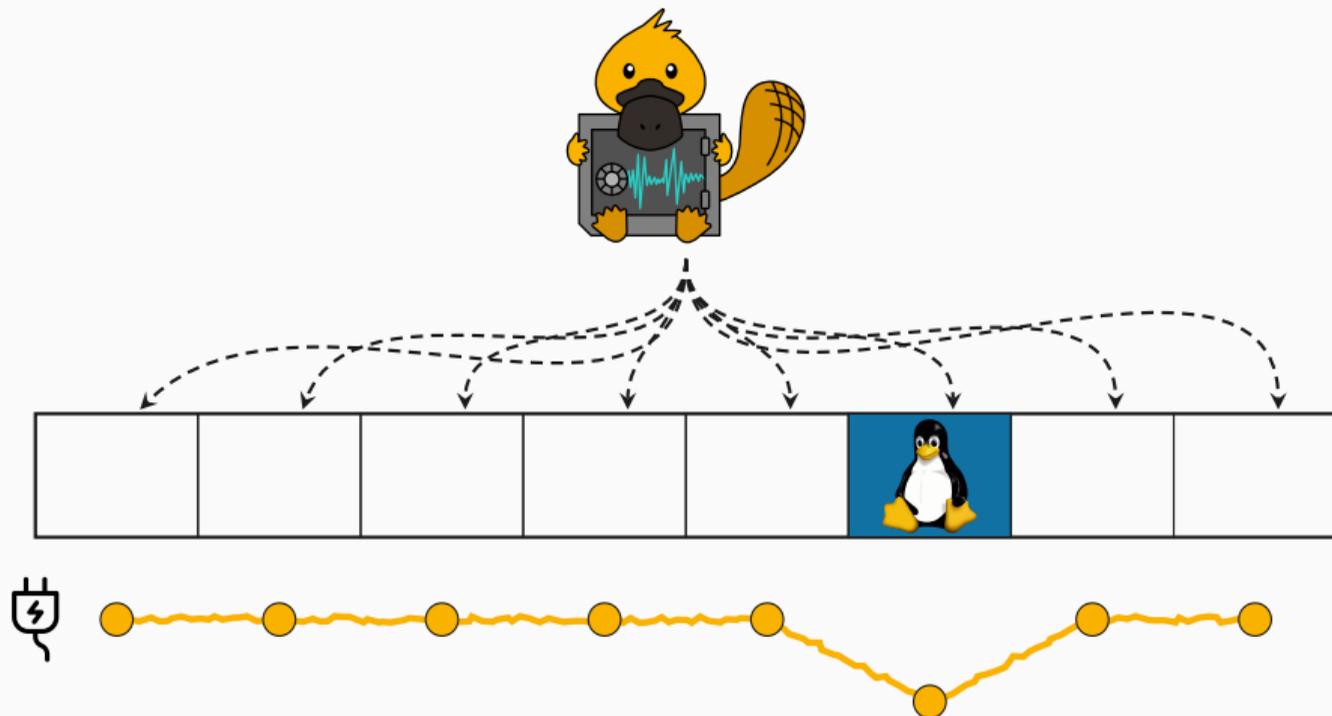
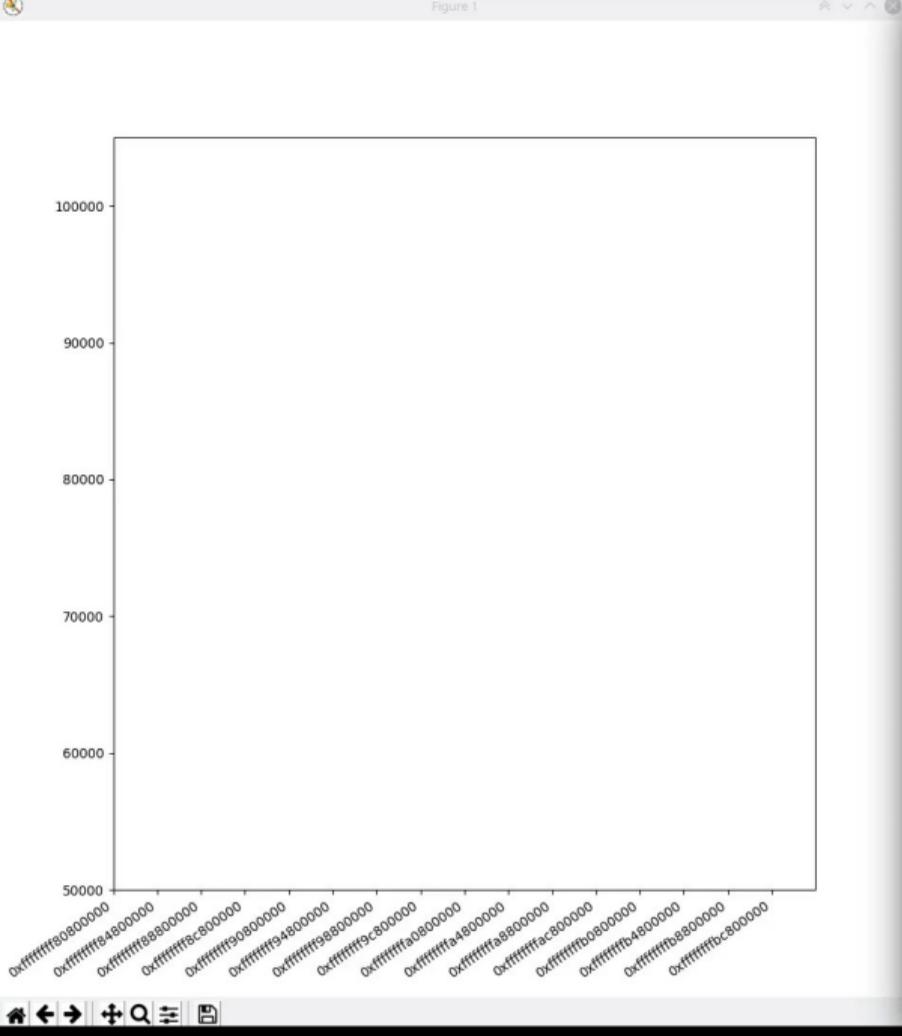
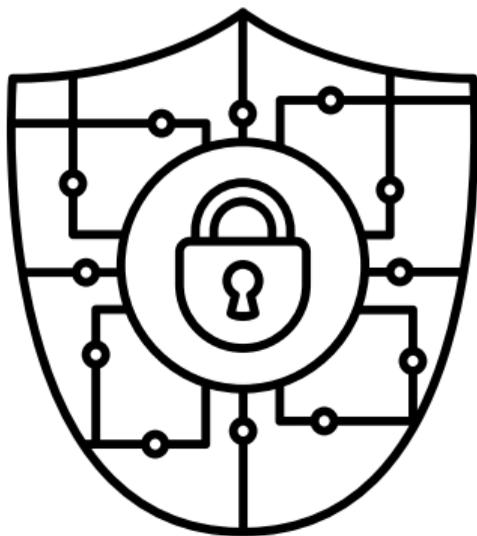


Figure 5: Repeated Page-table walks for unmapped pages require more power



```
File Edit View Bookmarks Settings Help
kaslr : zsh — Konsole
michael@hp /tmp/kaslr %
```



Attacking Intel SGX: RSA Key Recovery



- Instruction-set extension
- **Integrity** and **confidentiality** in **untrusted environments**
- **Enclaves** offer **protected areas of memory**
- **Operating system** can be **compromised**

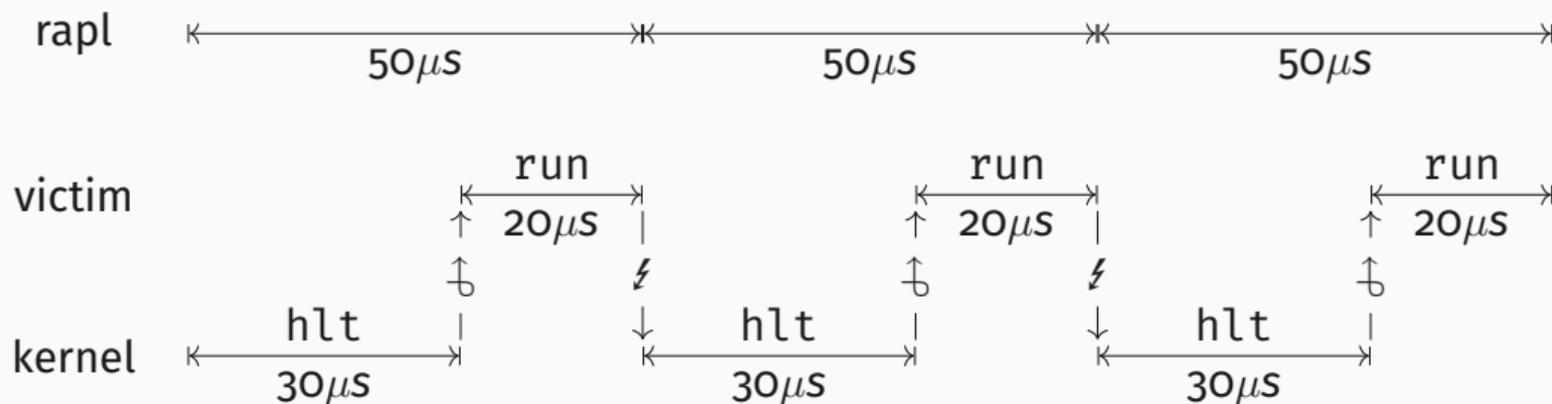


- **More power** as an evil operating system
- Hook the SGX Enclave exit point
- **Directly** read out the **RAPL values** from the MSR
- No operating system overhead!
- Interrupt victim often to **increase** resolution

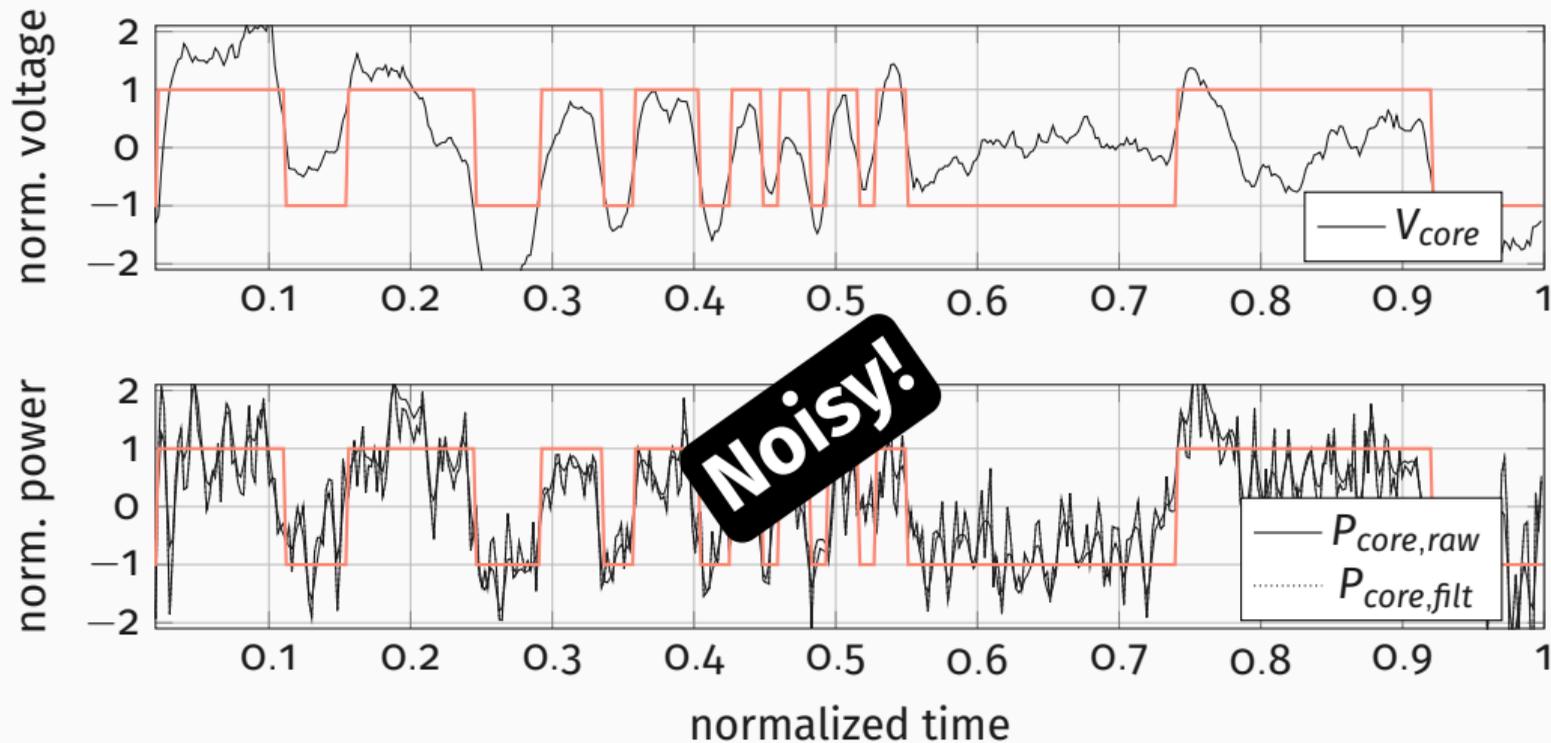
Halt Delay



- RAPL domains have a nearly **fixed** update interval
- Delay the interrupt return with the halt delay in the ISR
- Reduces the **execution time** of the victim in the current interval



SPA Attack - Results



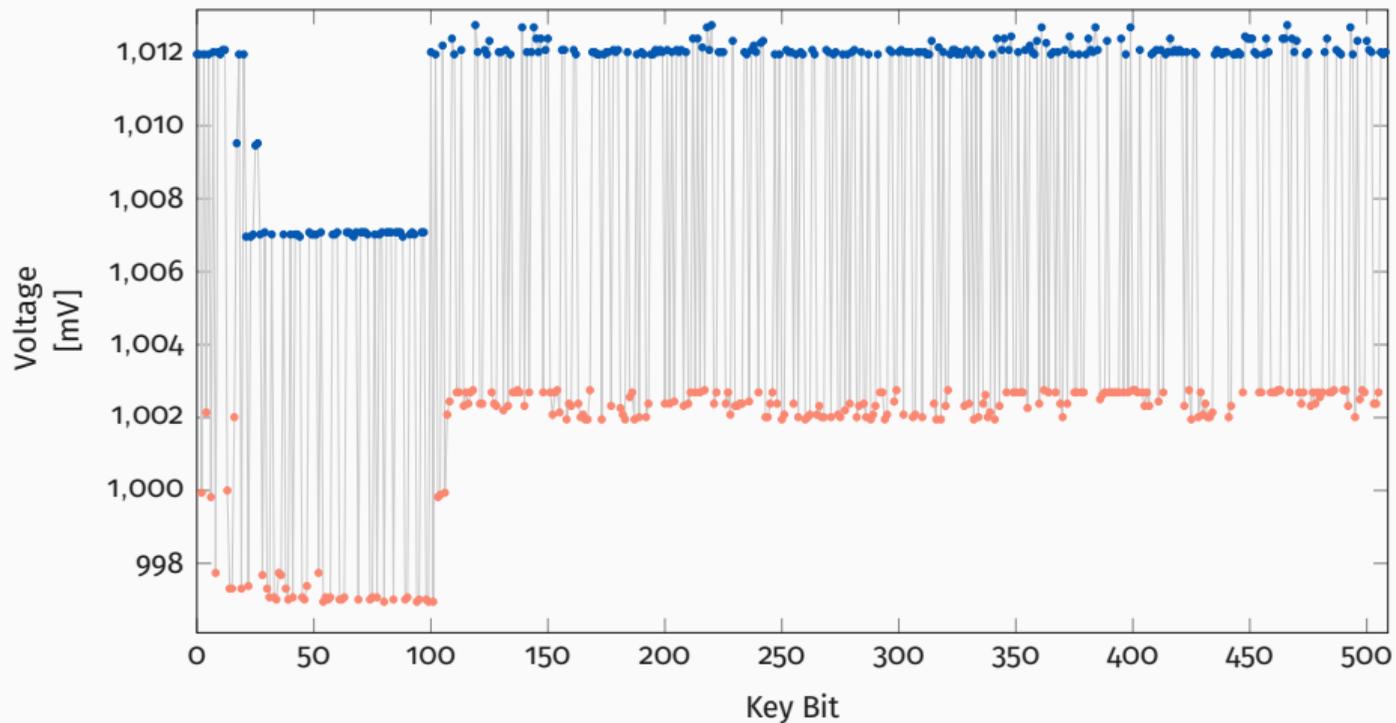


- **SGX-step** is an open-source Linux kernel framework
- Configure **APIC** timer interrupts
- **Single** and **zero-step** enclave execution



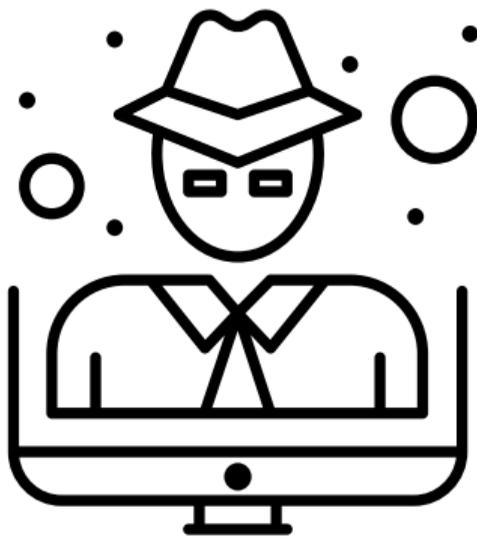
- **Combine Intel RAPL** with **SGX-step**
- Measure the energy consumption of **single instructions**

Attacking mbed TLS





- Time per key bit increases **linearly** based on the index
- **3h 31m** for a **512 bit**
 - **52 minutes** for finding target instruction
- Record 3 samples per key bit
 - This could be extend to a **single** trace attack



Crypto Attacks from User Space



- **Difficult** to measure parts without SGX-step
- Can **measure** over the **overall execution**

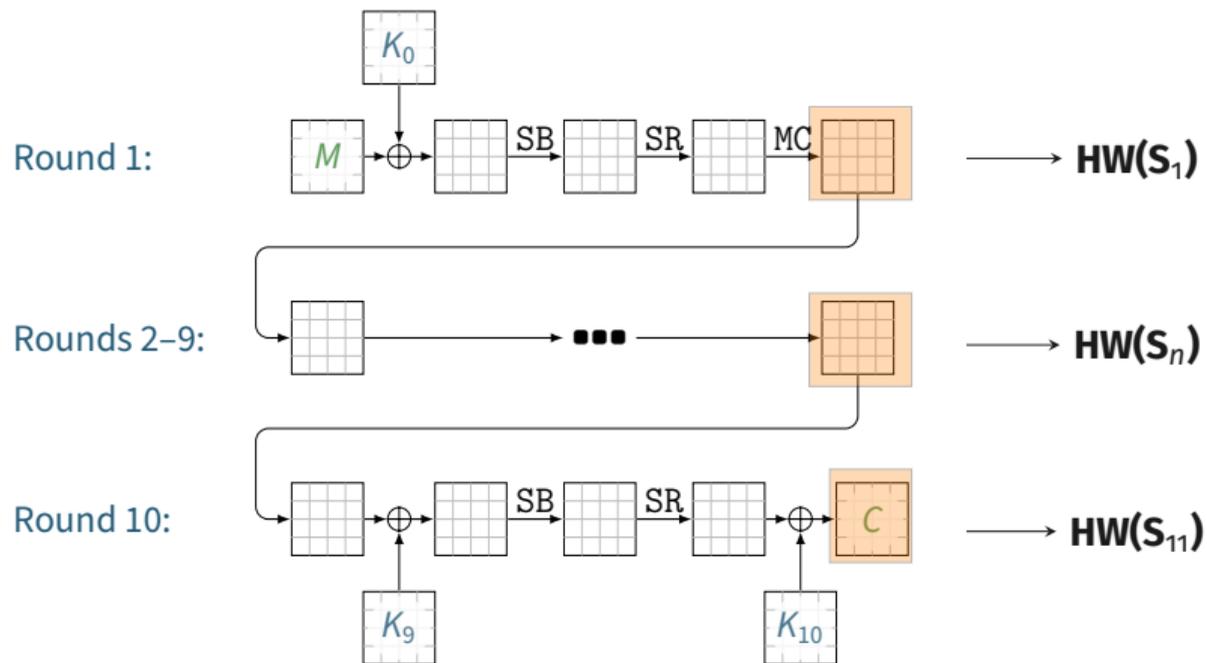
- Building a power consumption **model** of the device:

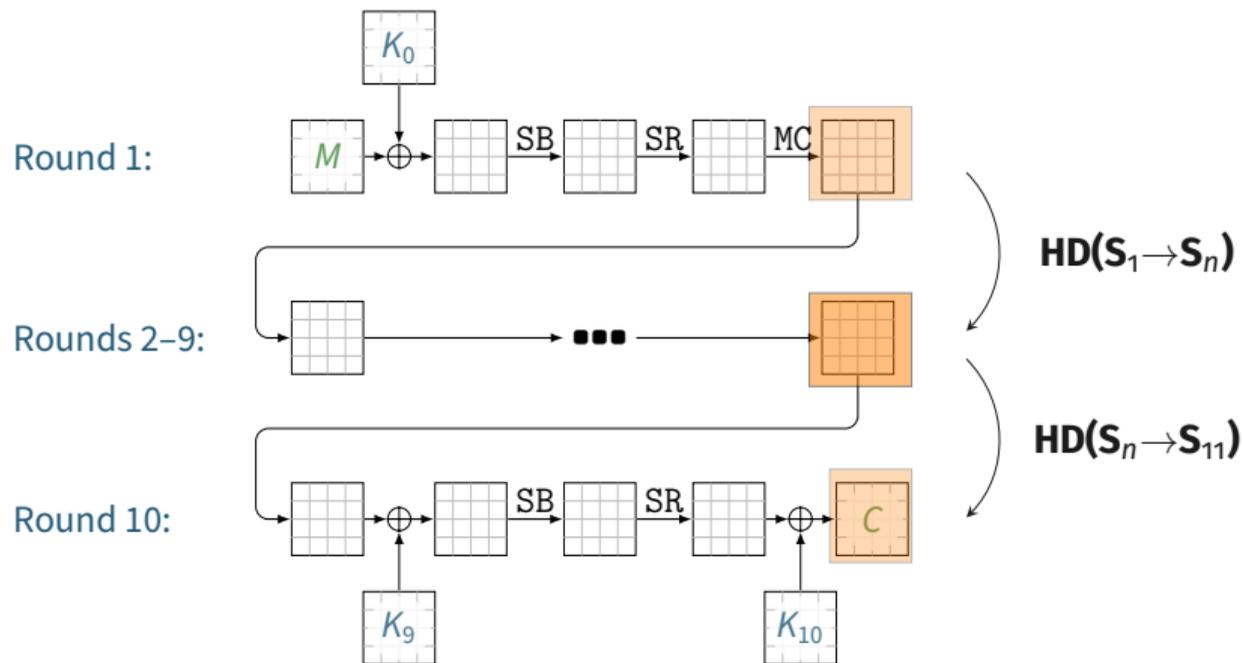


Hamming Weight
Number of bits set



Hamming Distance
Bits flipping between operations







- **AES-NI**: Side-channel resilient instruction-set extension
- Target **AES-NI** in a scenario where we can trigger encryption/decryption of many blocks
 - Disk encryption/decryption
 - TLS
 - (Un)sealing SGX enclave state

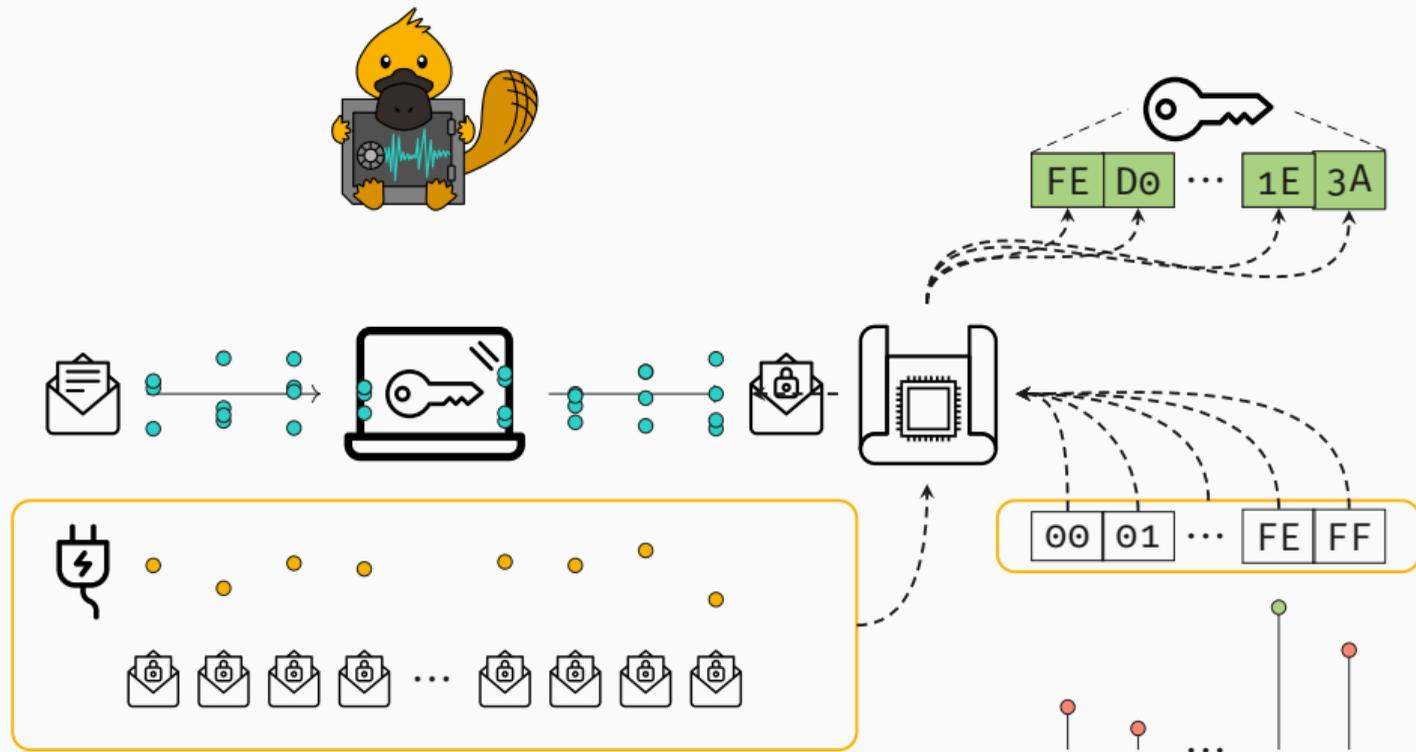
Correlation Power Analysis



- We **control** the plain text
- We **observe** the cipher text
- We **measure** the energy consumption over many operations
- We **guess** the key

- With our **model** and all **possible values**, **where** is the **correlation** the **highest**?

CPA Attack





- AMD **affected** as well
- Never heard back after disclosure
- Similar **Linux patch** as Intel



Countermeasures



- Remove the **unprivileged** access to the RAPL MSRs
- **1 Line Patch** for the Linux Kernel



- Threat model of SGX allows a **compromised operating system**
 - Operating system patch does not help
- **Microcode updates** are **necessary**
 - Fallback to a **model** of the energy consumption
 - Does **not allow** to distinguish data/operands any more
 - **Constant-time implementations** are **necessary**



- **Power side-channel attacks** can be exploited **from software** on modern CPUs
- Threat model of Intel SGX requires more **complex mitigations**

Remove Interface = The End?

- Home
- Shorts
- Subscriptions
- Library
- History
- Your videos
- Watch later
- Liked videos

Subscriptions

- Music
- Sports
- Gaming
- Movies

Explore

- Trending
- Music
- Movies
- Gaming
- News
- Sports

More from YouTube



0:19



3:06

Why MacBooks Get So Hot

290K views • 1 year ago

Apple Explained

If you've been using Apple notebooks for a while, you may've noticed how hot they get when sitting on your lap or playing games.

4K CC



15:01

How To Keep Your Macbook From Overheating (Top 10 Tips)

252K views • 2 years ago

Tom Scryleus

All of my gear → <https://kit.co/TomScryleus> Support this project ...

4K CC

Intro | What is Overheating | Tip 1 Understand Your Limitations | Tip 2 Consider Your Surface | Tip 3... 11 chapters



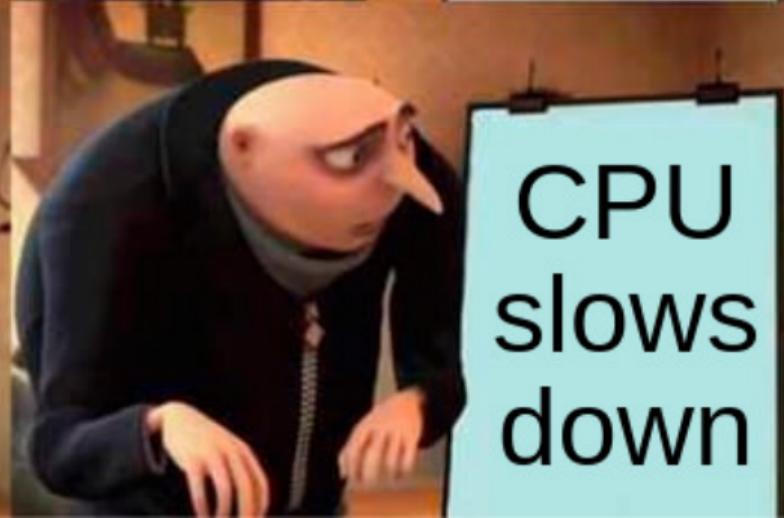
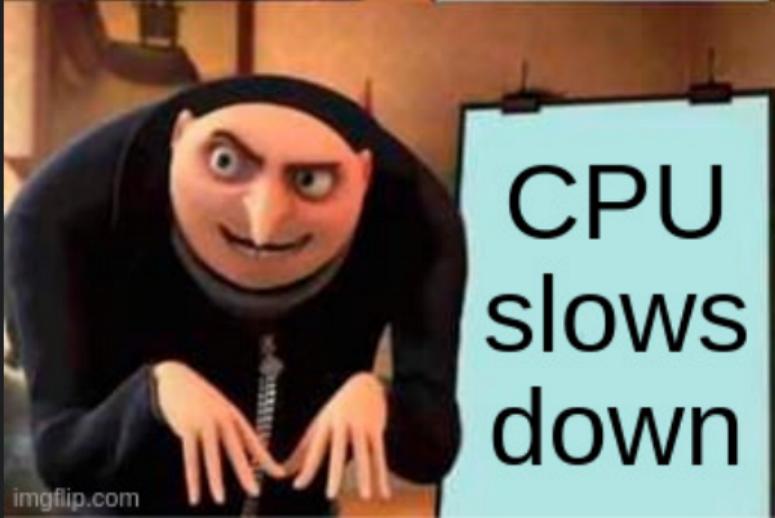
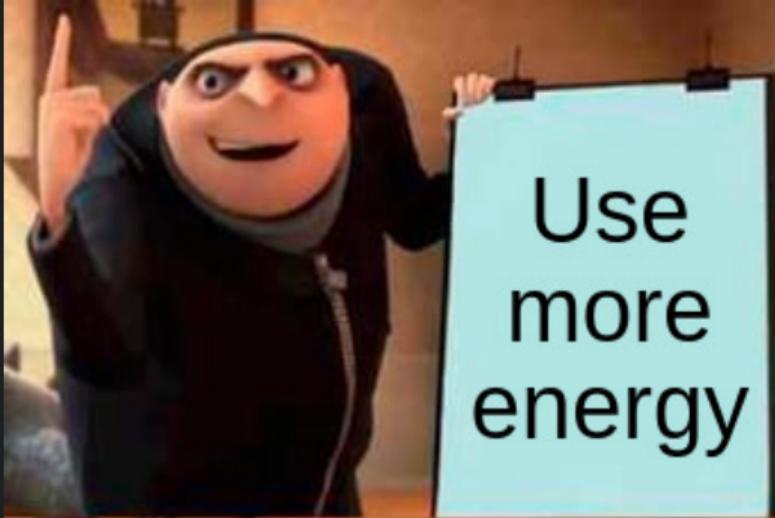
OVERHEATING MacBook Pro! Can We Fix It??

293K views • 5 years ago

Hardware Canucks

Macbook Pro is overheating... lets fix it :) Buy items in this video from Amazon at the links below: BUY The Phanteks HALOS RGB ...

4K



Remember?

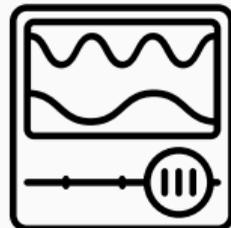
- CPU power management is **complex**
- In order to **save power**, you can ...



Shut down resources



Reduce **voltage**



Reduce **frequency**

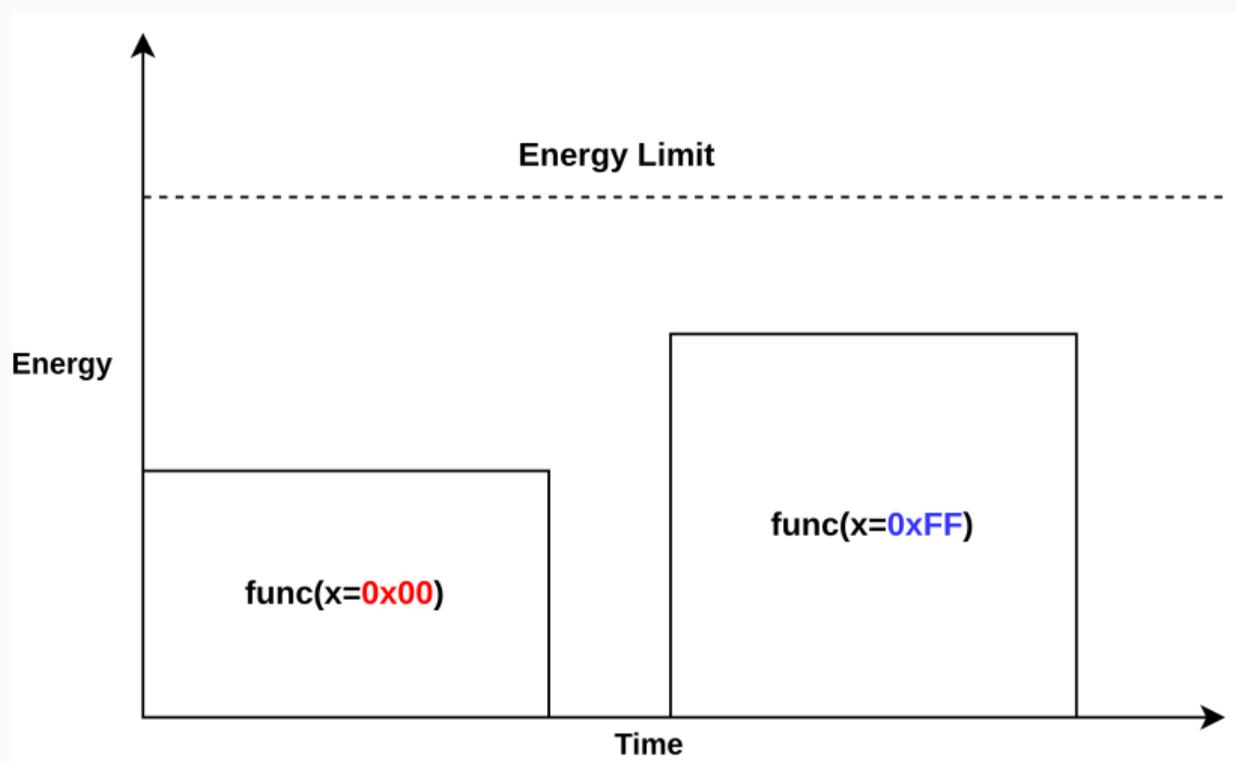


- The Hertzbleed attack from Wang et al. shows:
 - If more **energy** is used
 - The CPU gets **hotter**
 - Until the frequency is no longer sustainable
- The runtime of the executed code **slows down**
- Measure with **fixed** clock, e.g., `rdtsc`

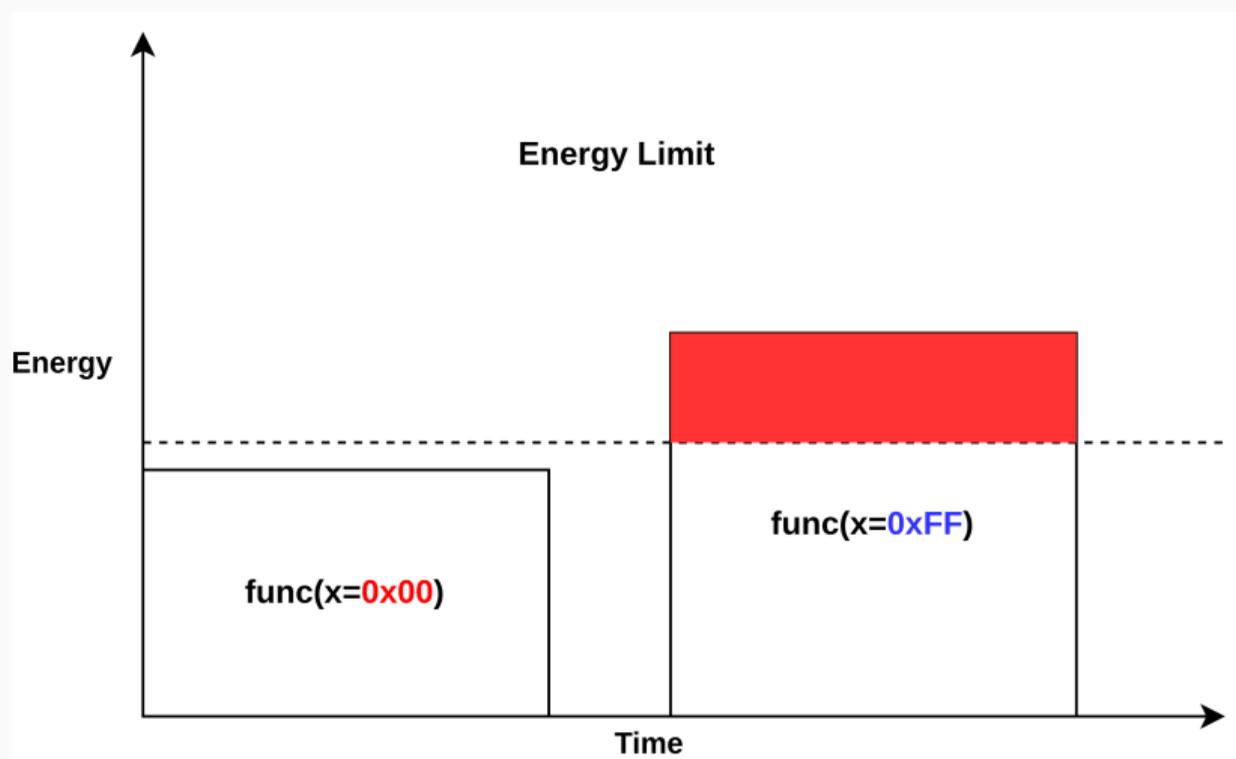


- RAPL provides energy **limits**
 - If exhausted CPU throttles the frequency
- Run **Stress** on the system
 - CPUs start throttling when using many threads

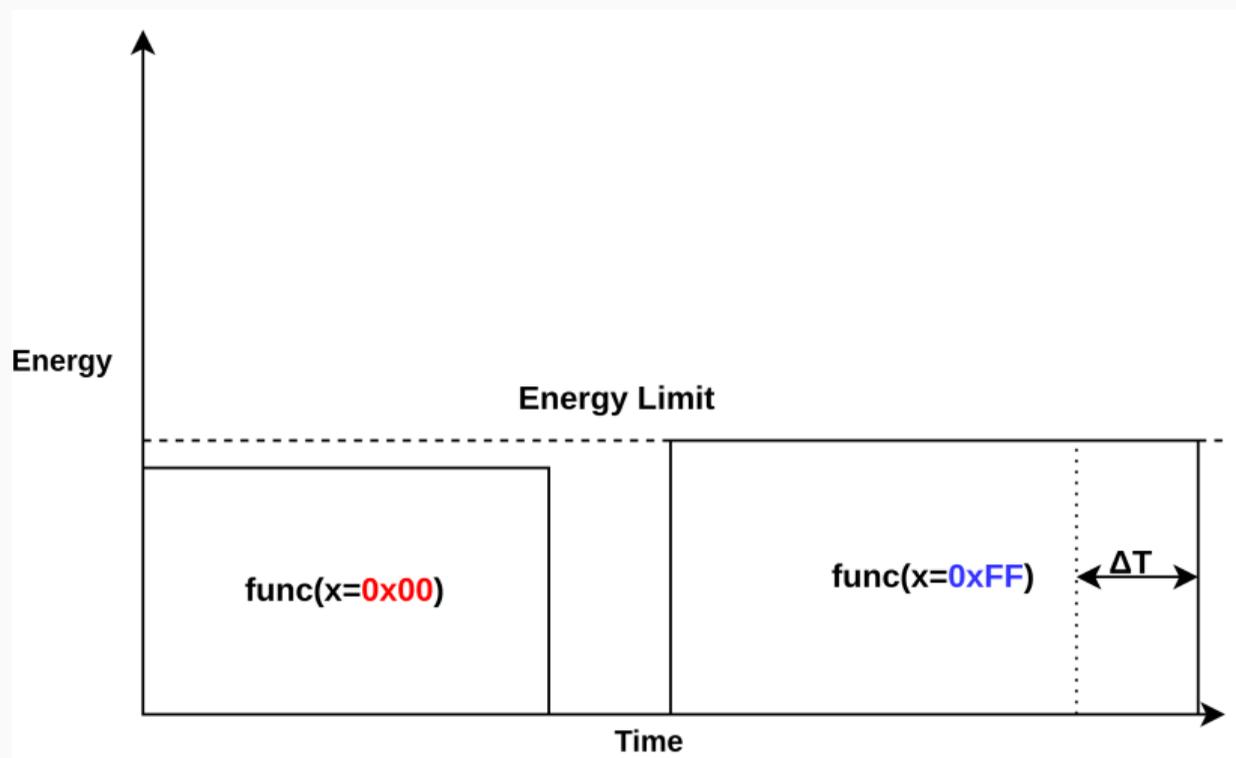
Converting Energy Differences



Convert Energy Differences



Convert Energy Differences







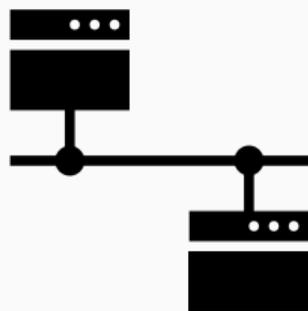
What can we do with this?



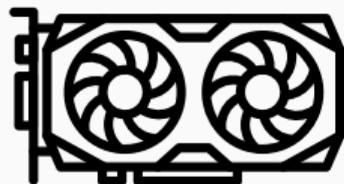
- **Hidden** communication channel
- **No** power interface required
- **Time/Frequency** measurements proxy power interface



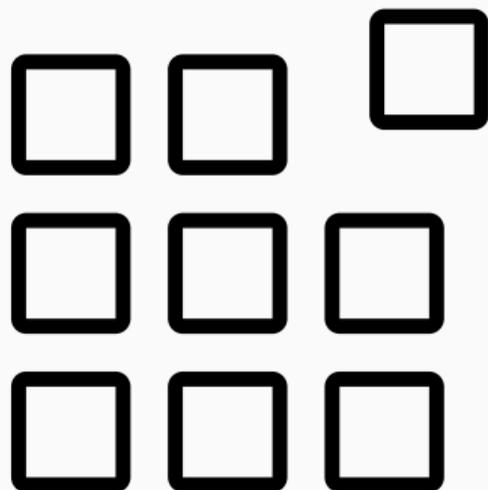
- **AES Correlation Power Analysis**
 - Measure **execution time** of AES encryptions
 - Apply CPA technique to recover key



- **Remote attacker** requests service from server
 - Cryptographic operation, i.e. encryption, signature
- Server computes response using secret
- **Hertzbleed-effect** influences **response times**
 - **Calculations using secret** influences server CPU frequency
- Attacker **recovers secret** using collected timings



- **Integrated** GPUs **share** power limits with the CPU
 - **CPU throttling** indicates high GPU consumption
- **Dedicated** GPUs have power limits too
 - **Observable** by **timing** a GPU workload



- What **secrets** are “*inside*” a GPU?
 - GPU renders windows and screen
→ **Privacy** related information
- **Pixel** color **represents** the information



- **Post-processing** **without** revealing the pixels
 - Pixel value is the **data operand**
 - Distinguishable power consumption
 - **Bright** pixel → **less** power
 - **Dark** pixel → **more** power
- **Measure timing and infer pixel value**

The End?

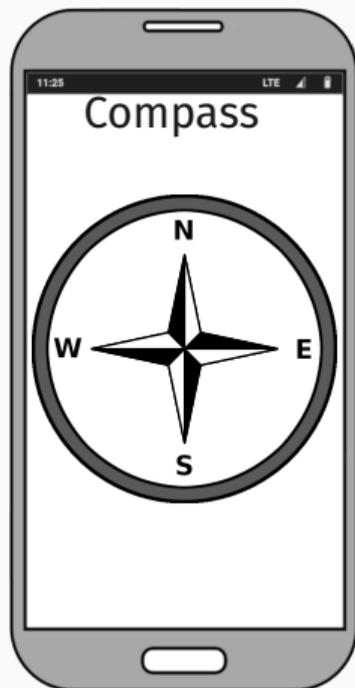


Are there other exploitable **power-related** signals?

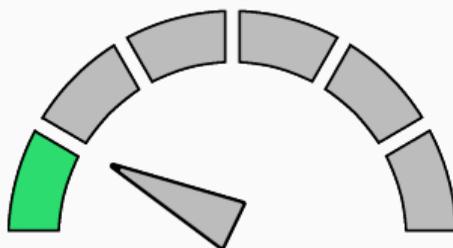
Android **power-related** side channel

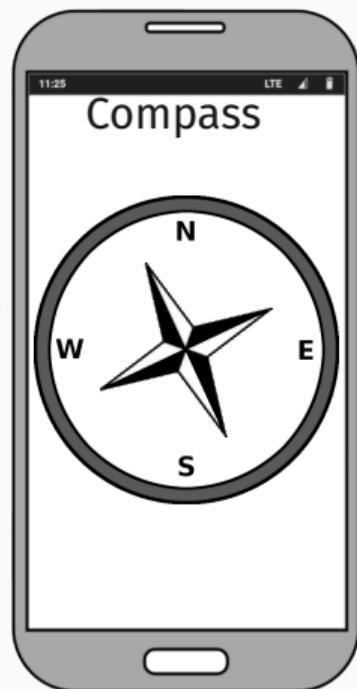
- Android sensor interface as a **proxy for power measurements** purely from software
- Systematic analysis of 9 Android smartphones:
 - ☛ Recovering leakage properties: **Integration interval, rotation-dependent leakage**
- Local attack:
 - ☛ Malicious app leaking processed AES key bytes
- Remote web-based JavaScript attack:
 - ☛ JavaScript **sensor-based pixel-stealing attack** leaking cross-origin pixels up to 5 s/pixel



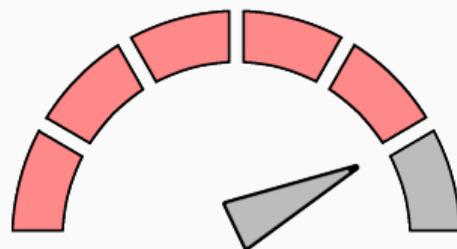


CPU utilization

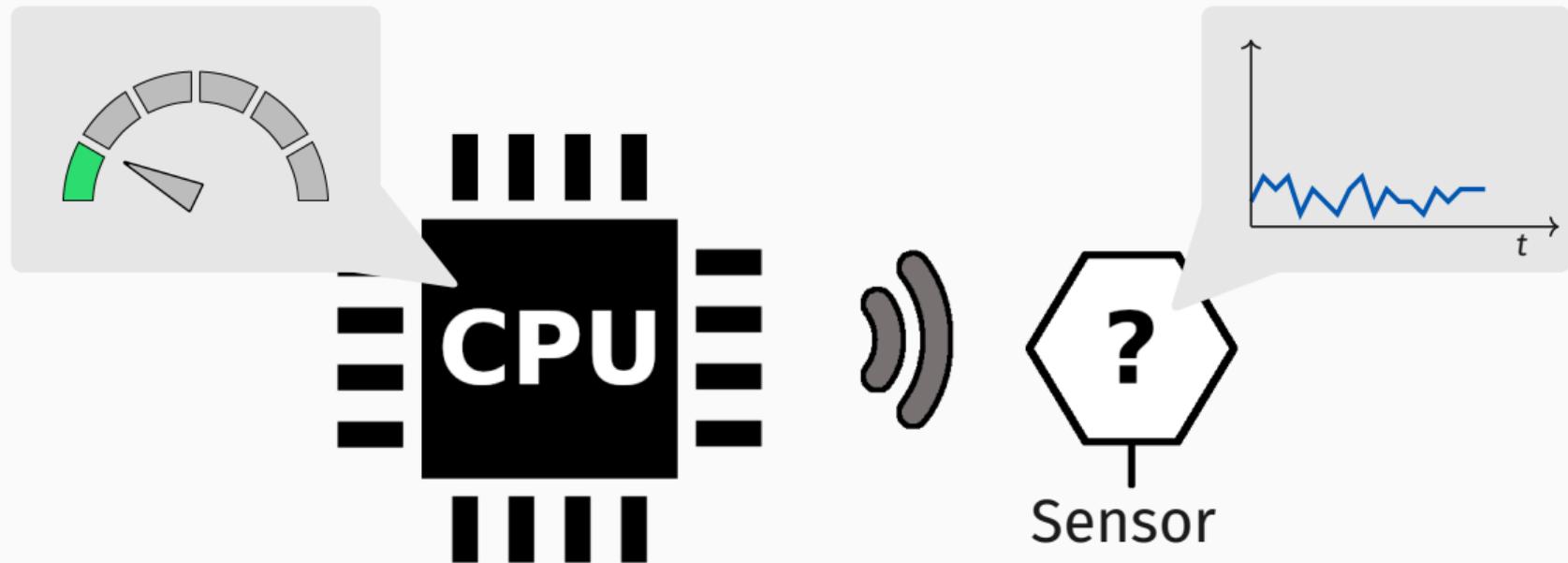




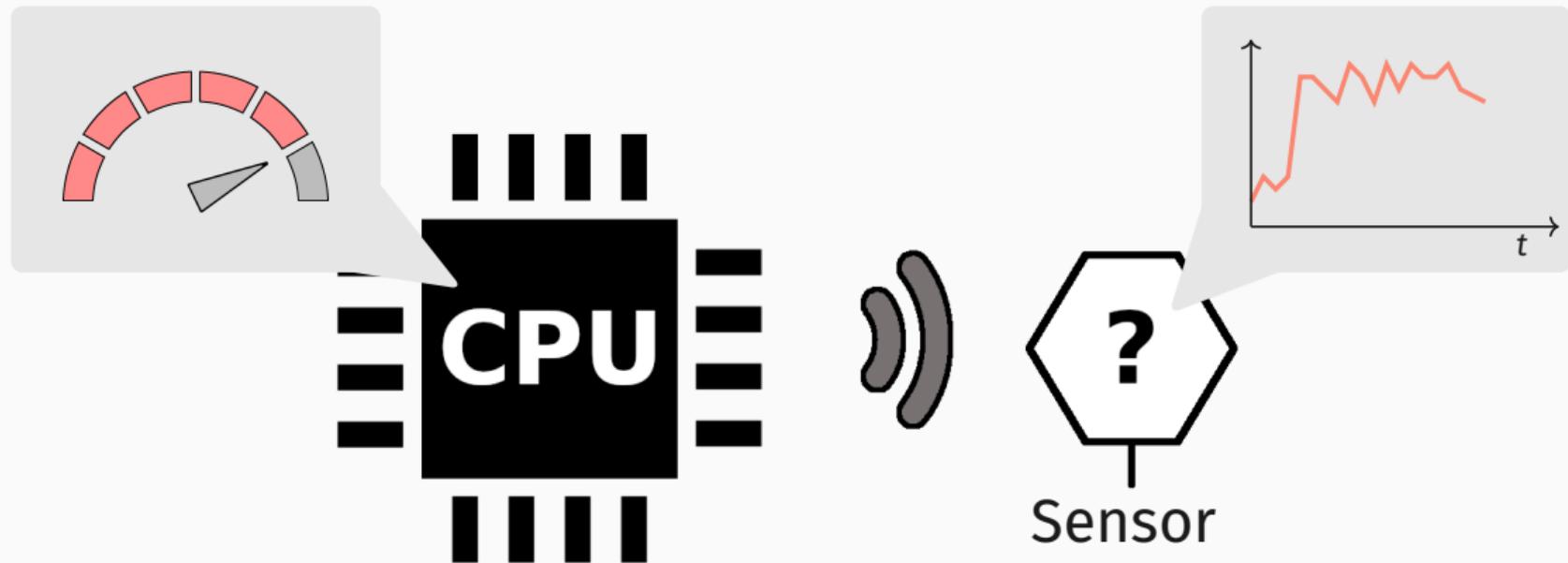
CPU utilization



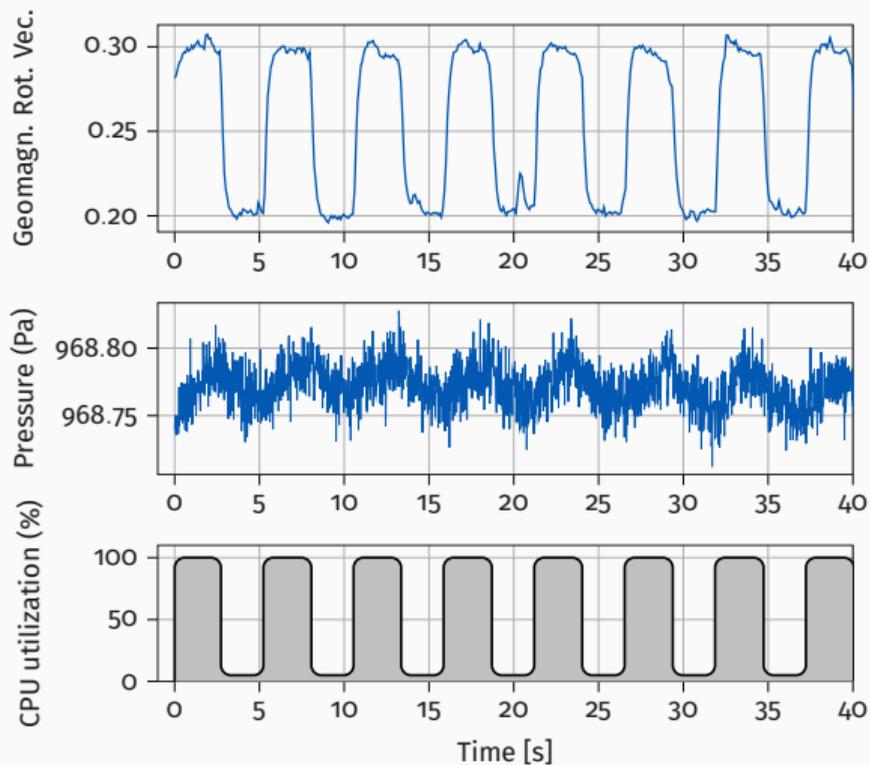
Motivation



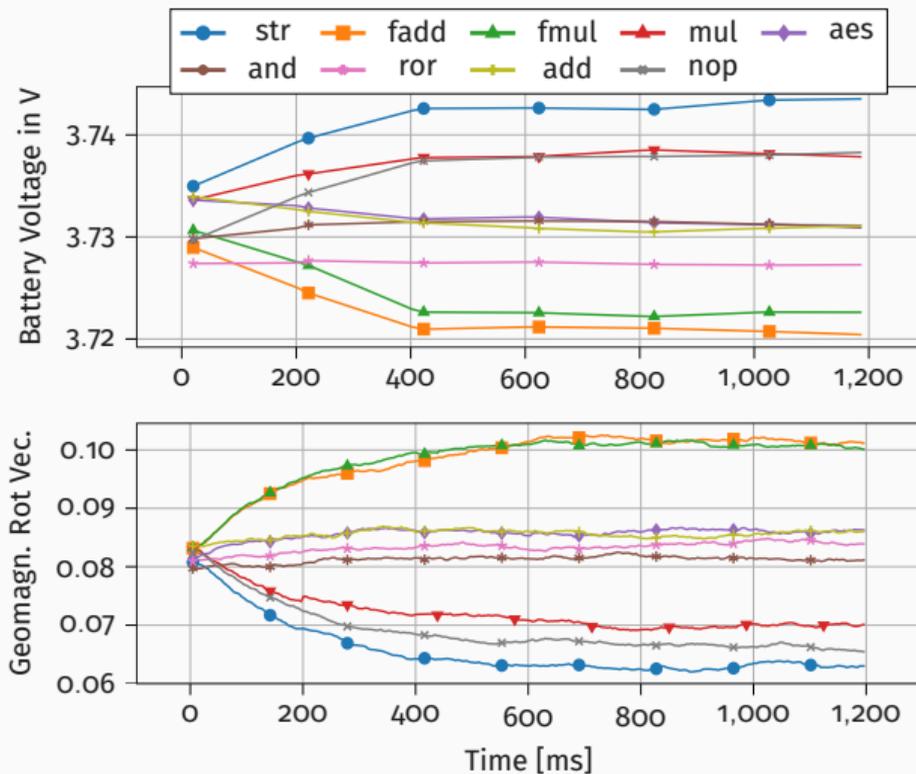
Motivation



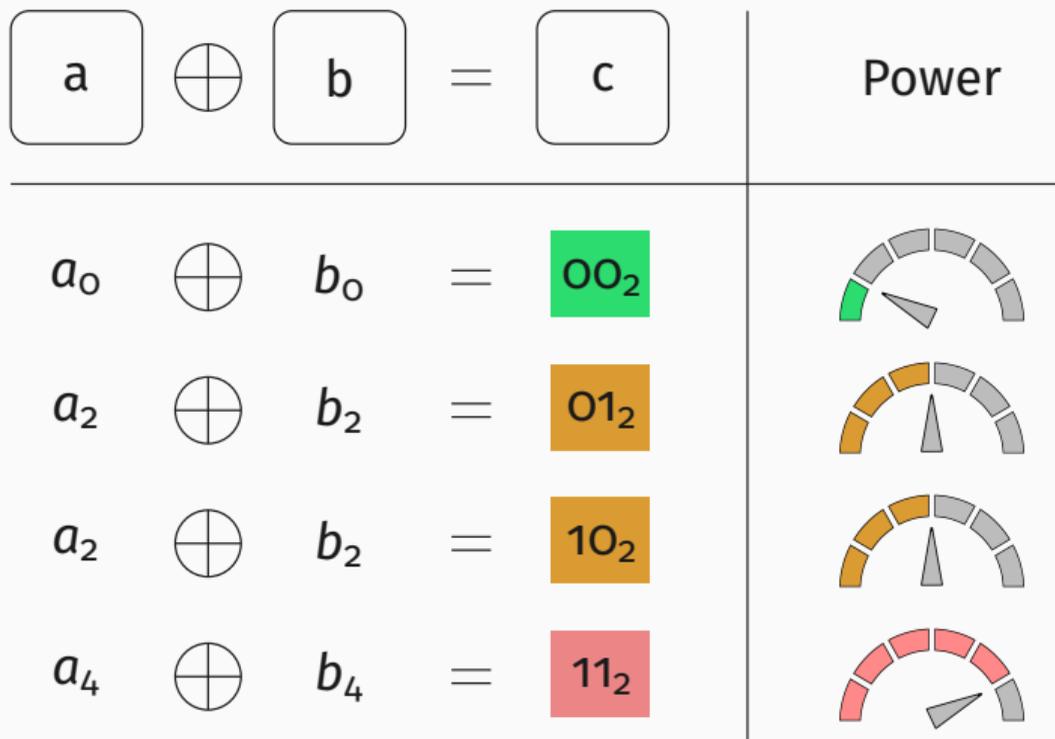
Systematic Evaluation: Varying CPU Utilization

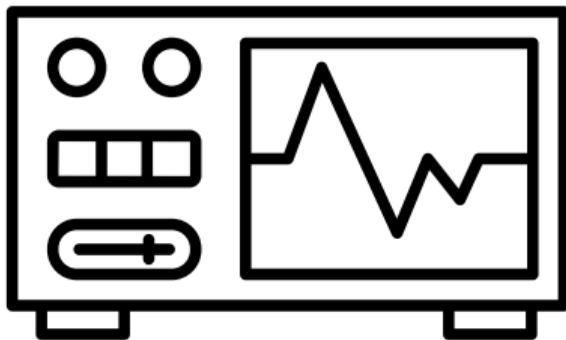


Systematic Sensor Analysis: Executed Instructions



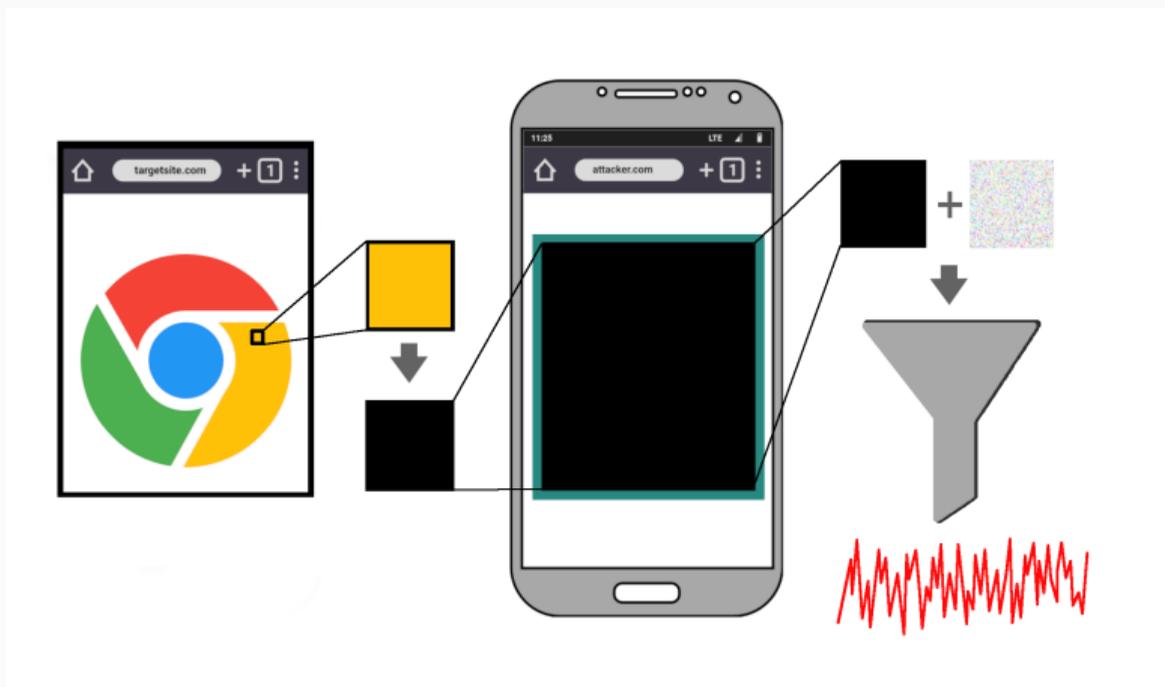
Systematic Sensor Analysis: Varying Data Operands





What can we do with this?

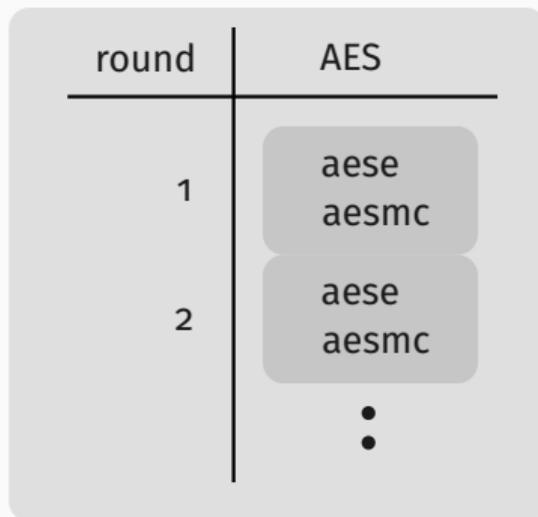
JavaScript Pixel-Stealing Attack



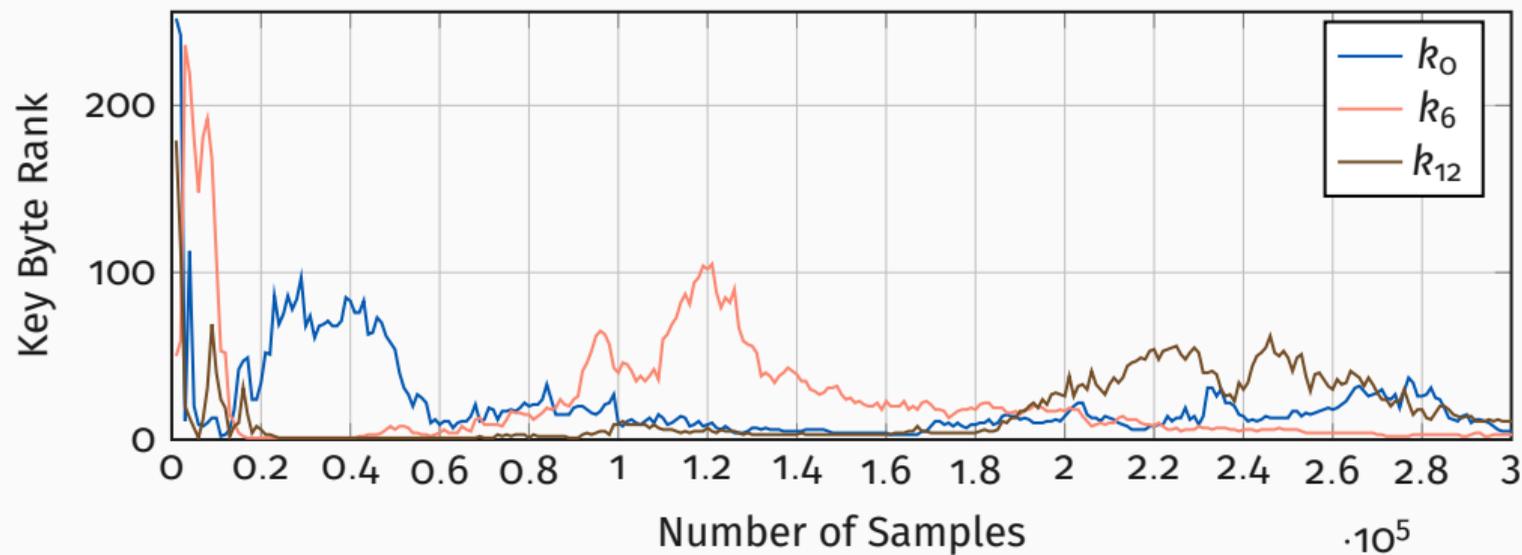
JavaScript Pixel Stealing: Evaluation

			
Image:	Original	Magnetometer	Abs. Orientation
Time/Pixel (s):		5	10
Accuracy (%):		90.2	70

AES Attack Case Study

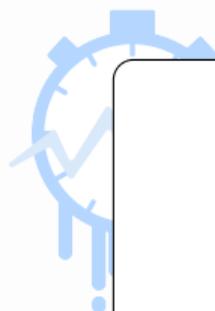
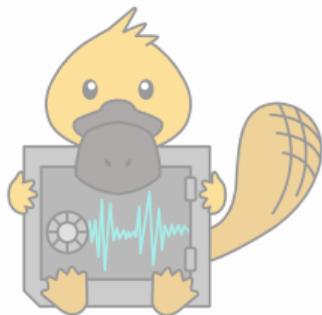


AES Correlation Power Analysis





How can we **transform** power side channels towards a broader scope?



Software-based Power Side Channel Attacks

- **Specific** targets: Algorithms
- Leak edge cases
- **Limited** to a side channel

Execution Attacks

- **Generic** targets: CPU components
- Leak arbitrary data
- **Agnostic** to side channels

Collide+Power



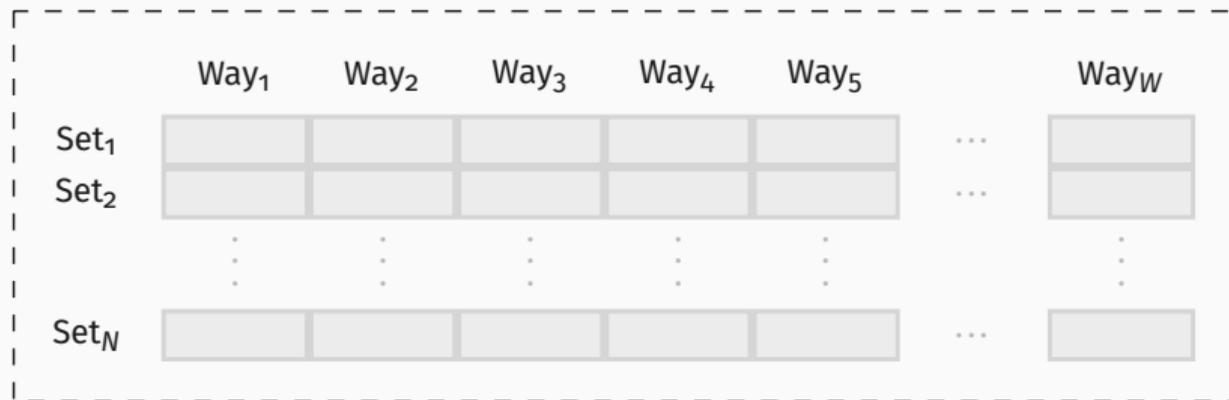


- **Collide+Power** exploits leakage between:
 - **Guess** \mathcal{G} : Attacker-controlled data
 - **Value** \mathcal{V} : Victim secret data

💡 Hamming distance: $\text{hd}(\mathcal{G}, \mathcal{V})$

→ **How to exploit this limited information?**

Collide+Power - Memory Subsystem



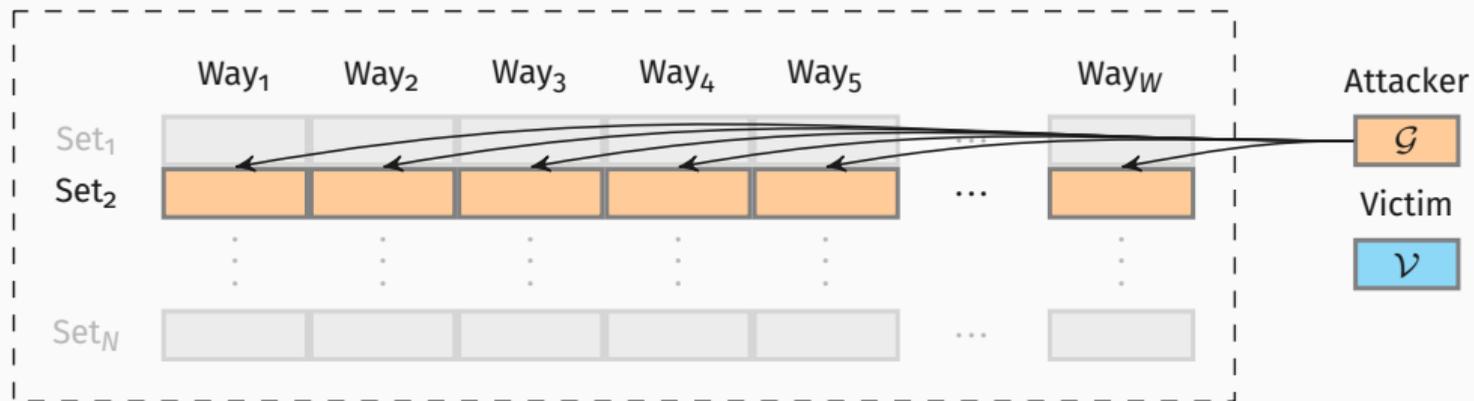
Attacker



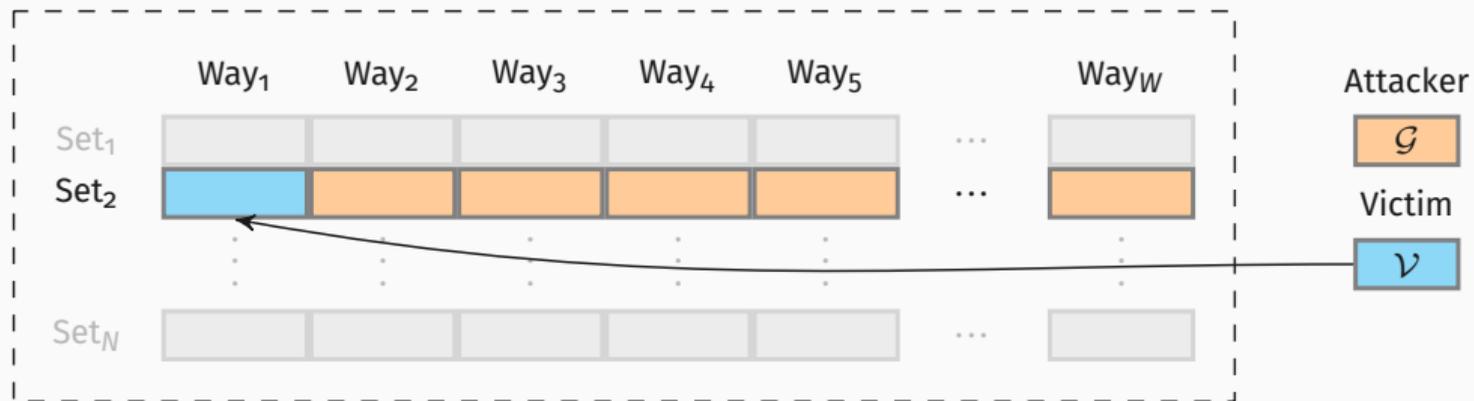
Victim



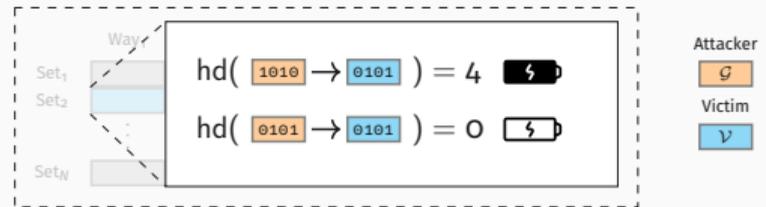
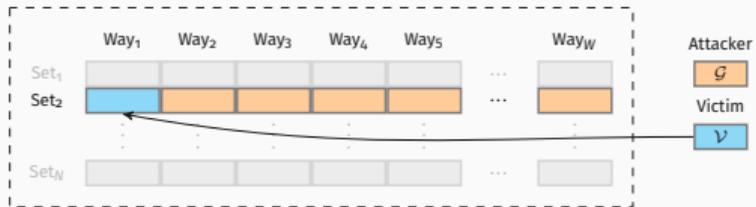
Collide+Power - Memory Subsystem



Collide+Power - Memory Subsystem

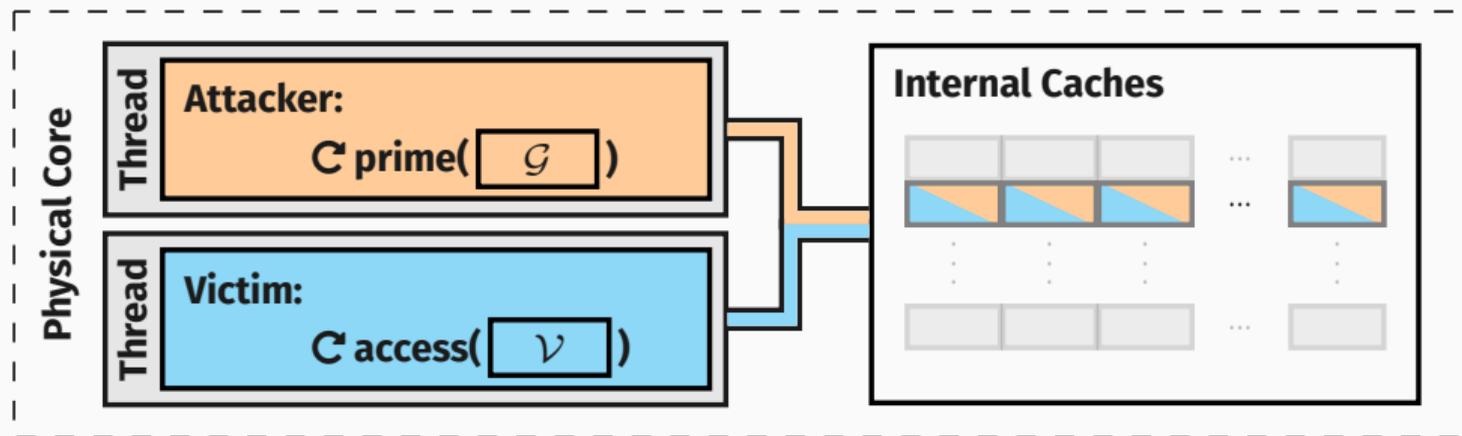


Collide+Power - Memory Subsystem

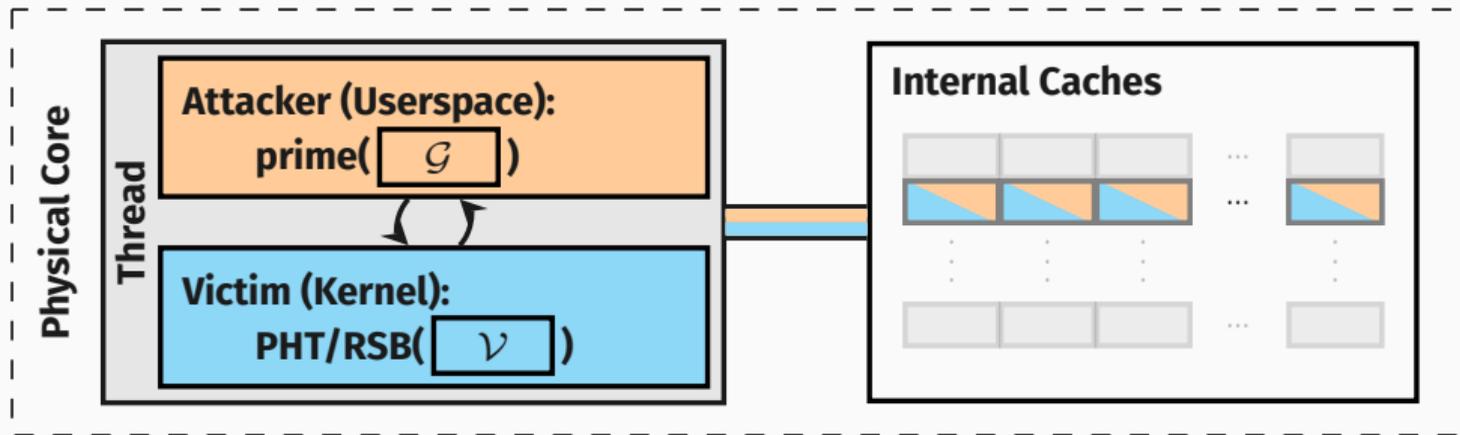


Generic Attacks

MDS-style Attack



Meltdown-style Attack

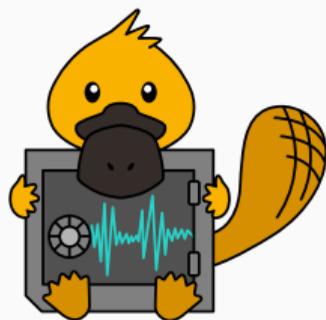


This must be slow?

NO!

It is EXTREMELY slow!¹

¹With the current state-of-the-art.



- **MDS-style:**
4.82 bit/h
- **Meltdown-style (RSB):**
0.84 bit/h



- **MDS-style:**
0.065 to 0.68 bit/h
- **Meltdown-style estimate (PHT):**
99.95 days/bit to 2.86 years/bit

Mitigations

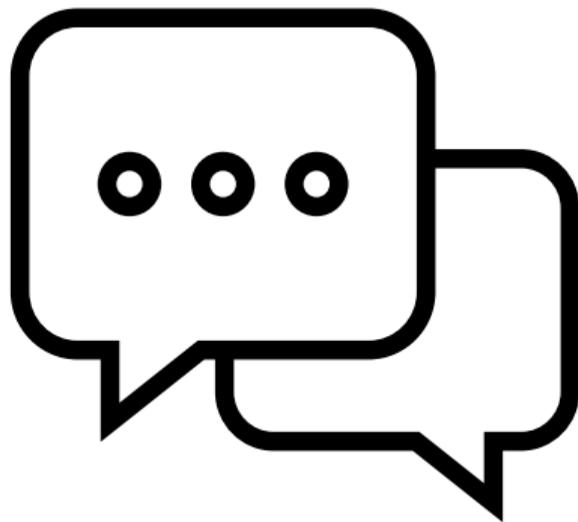




- **Preventing data collisions:**
 - **Redesign** of the **complete** shared data path
 - **Costly** to deploy
 - **Missed** components re-enable Collide+Power



- **Preventing observable power consumption:**
 - **Restricting** all direct power interfaces
 - **Mitigating** Hertzbleed is **challenging**
 - Thermal and power management is required
- **Collide+Power** is slow but **unmitigated** on modern CPUs!



Questions?