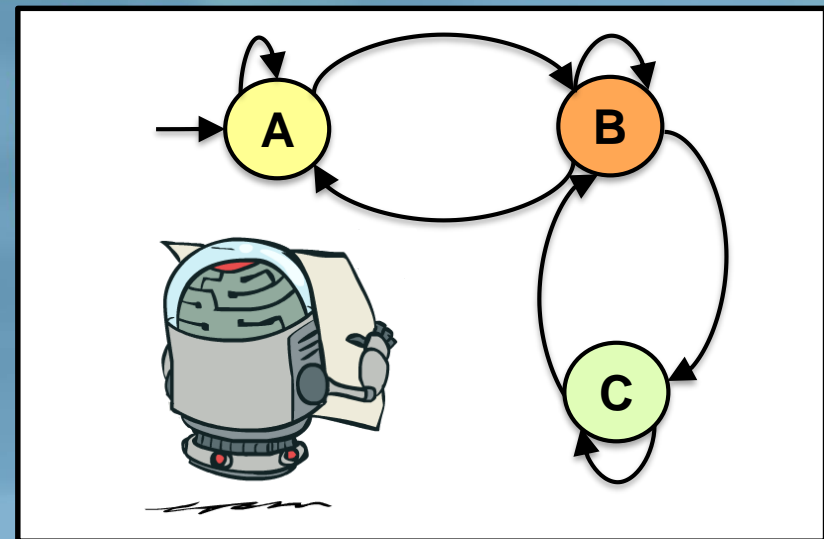
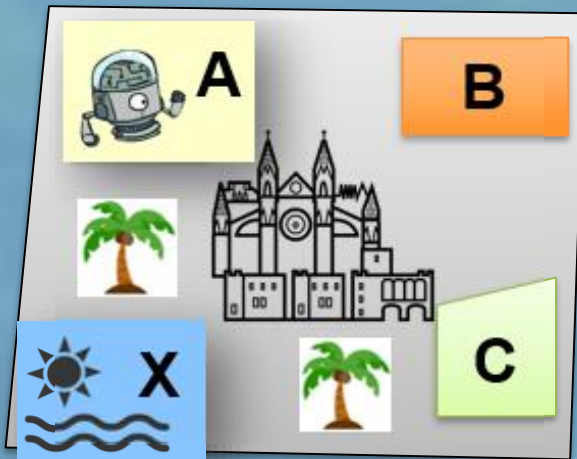
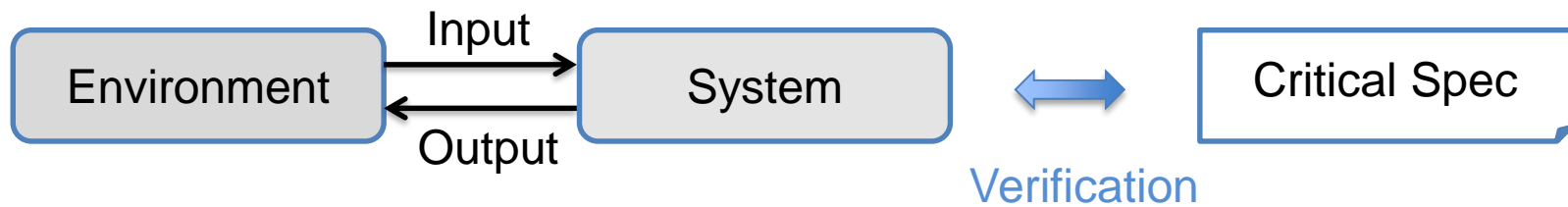


# Reactive Synthesis

Bettina Könighofer



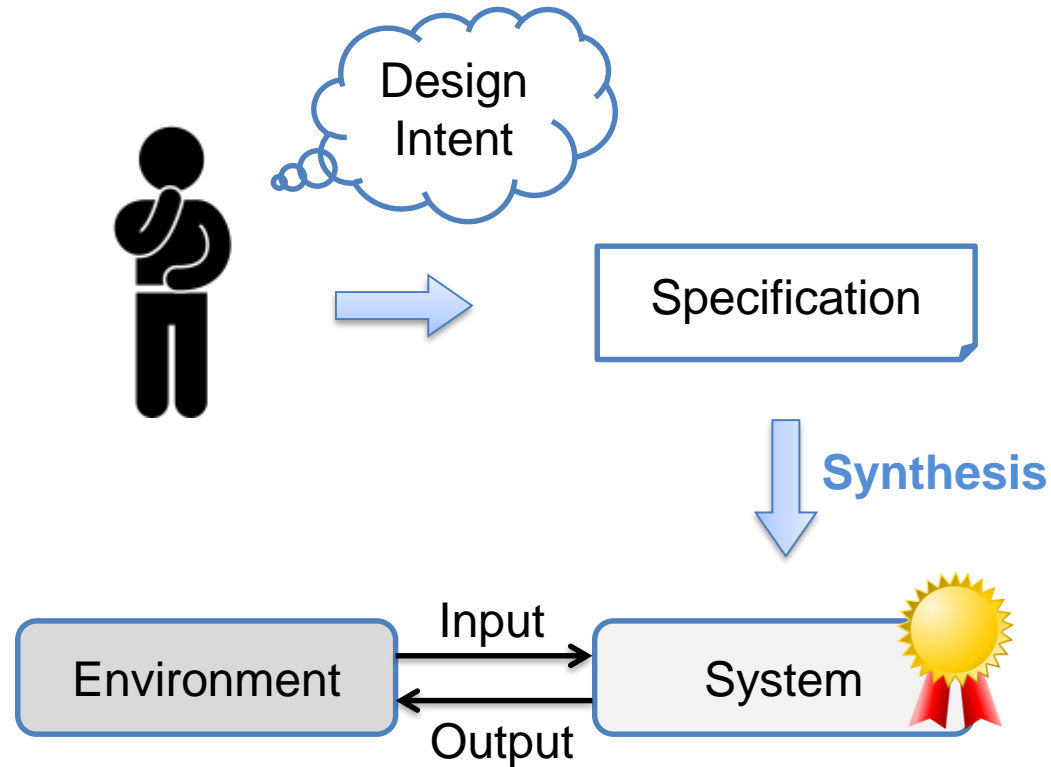
# Verification



**One needs to do a lot of work:**

- Need to write the system + specification

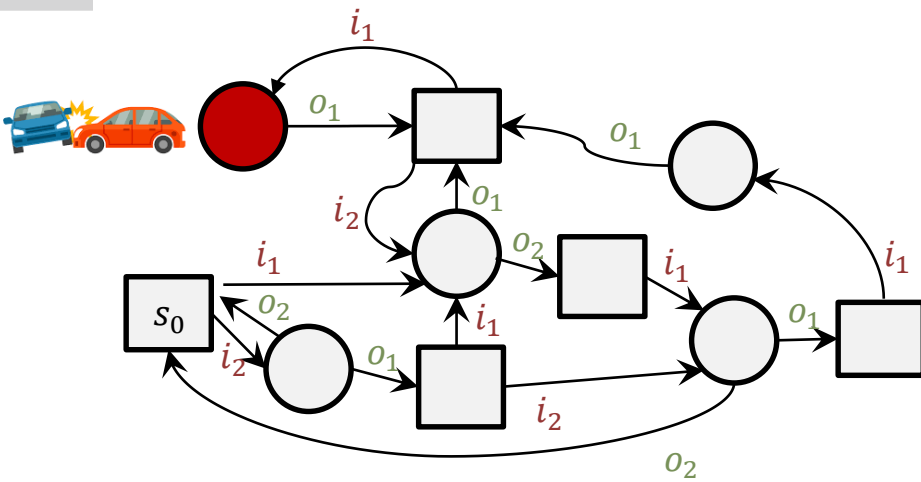
# Synthesize it!



## How can we compute the Sytem?

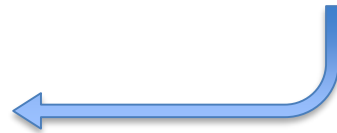
- By solving a **Game**
- Played between the **Environment-Player** and the **System-Player**

# Synthesis is a Game

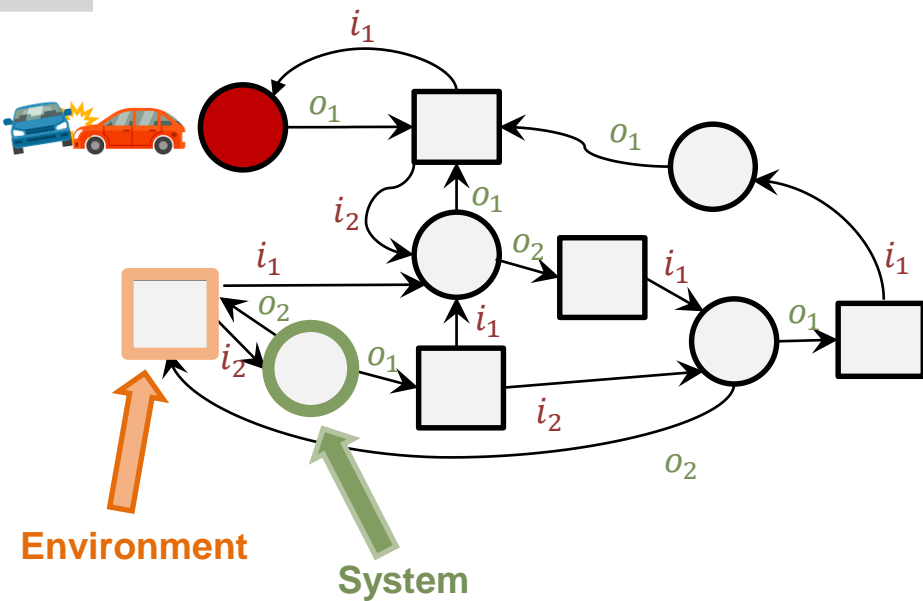
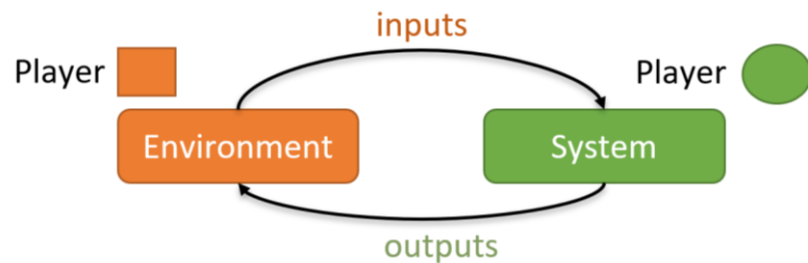


Formal safety specification

Model of environment

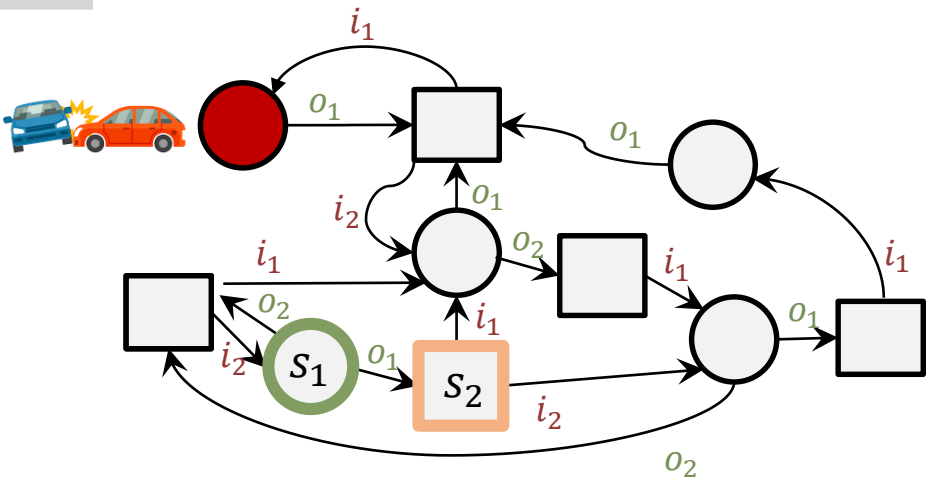
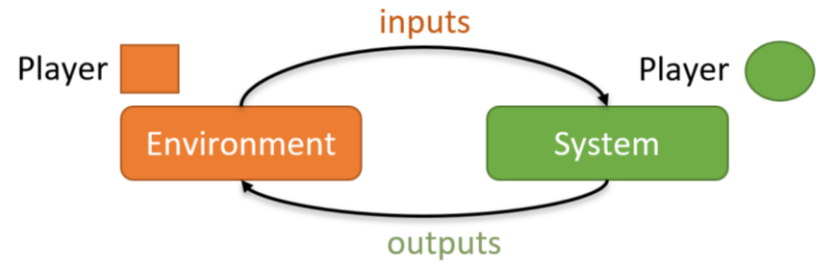


# Synthesis is a Game

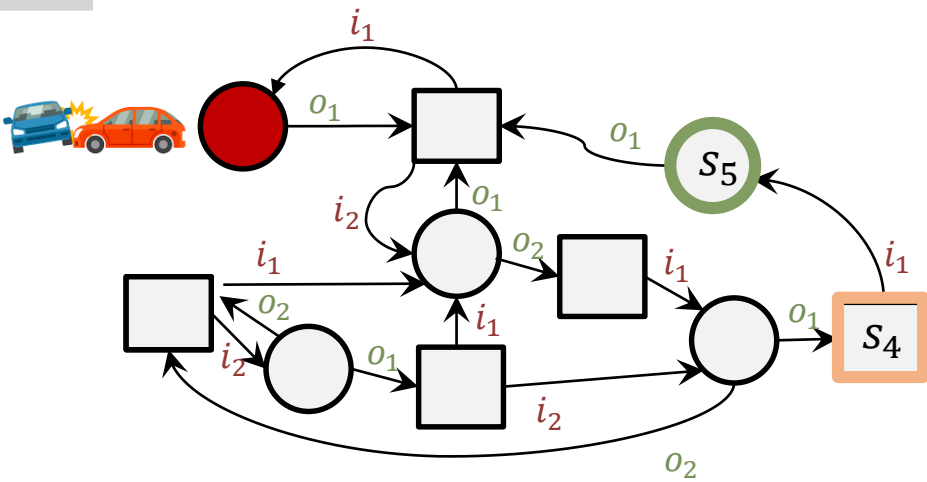
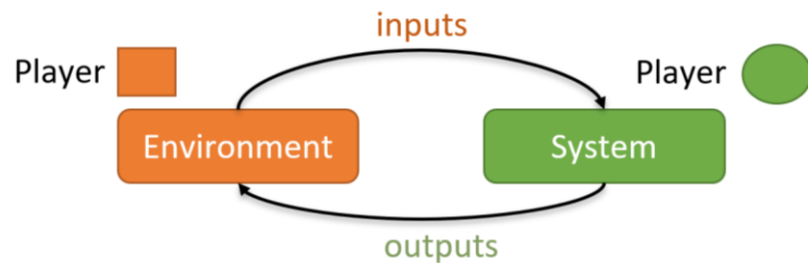




# Synthesis is a Game



# Synthesis is a Game

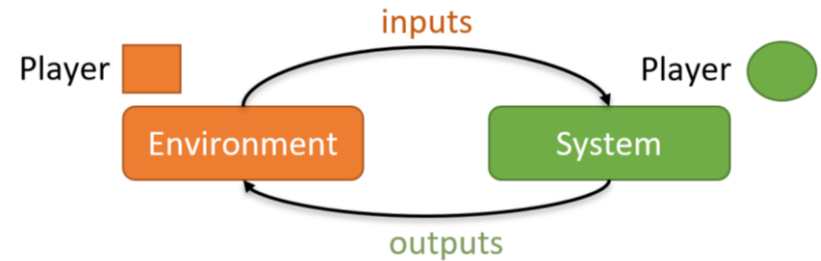


**System Player** wins, if is **never** visited

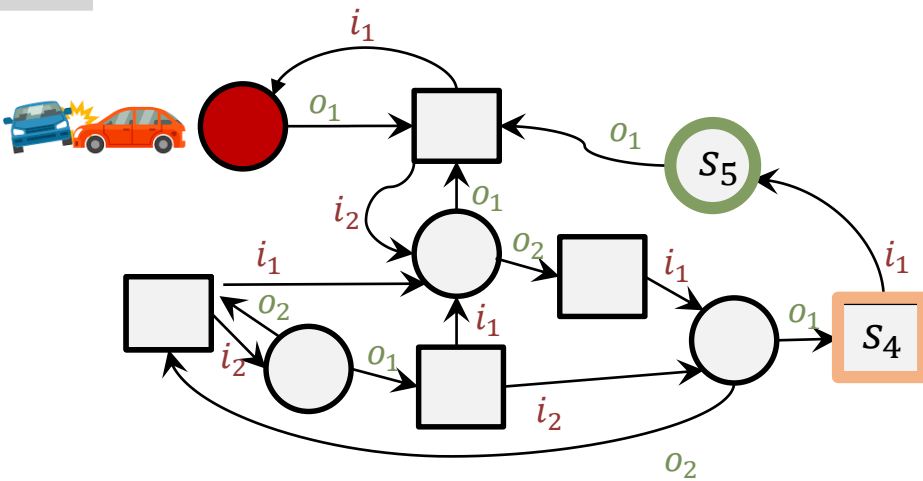
**Winning Region:** States from which the system can enforce that is **never** visited



# Synthesis is a Game



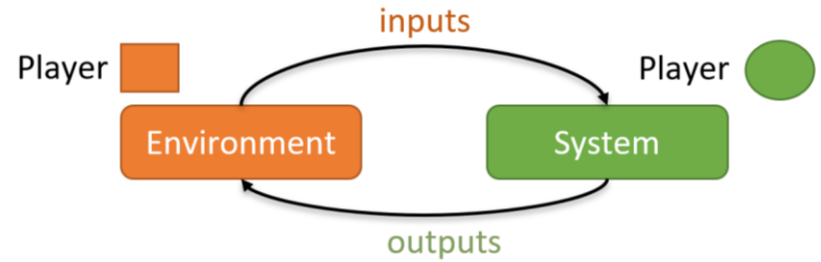
**What is the winning region for this example?**



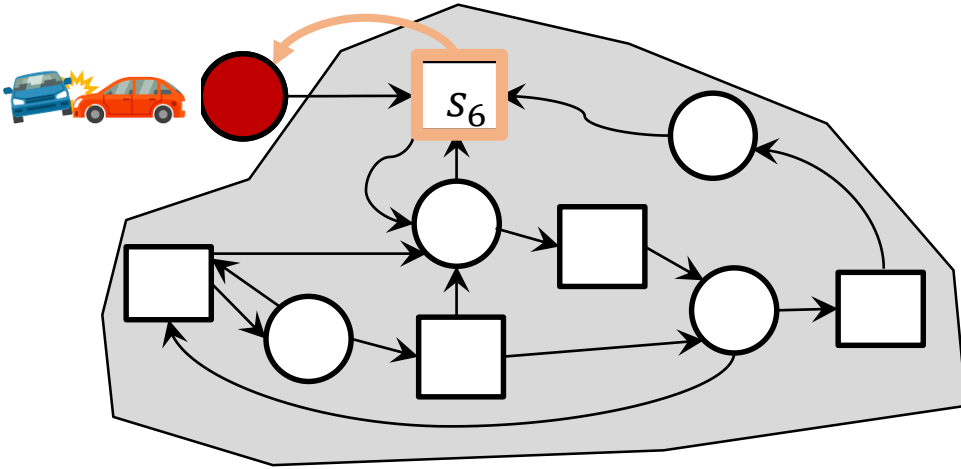
**System Player** wins, if is **never** visited

**Winning Region:** States from which the system can enforce that is **never** visited


# Synthesis is a Game



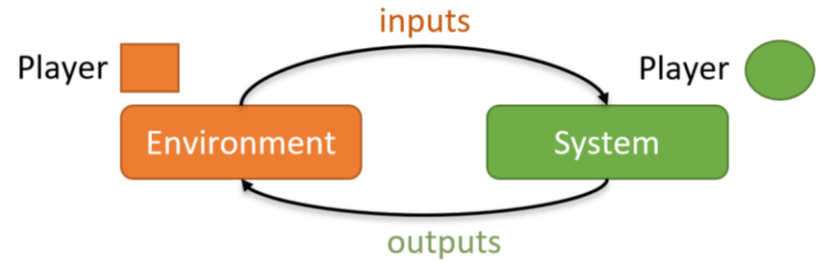
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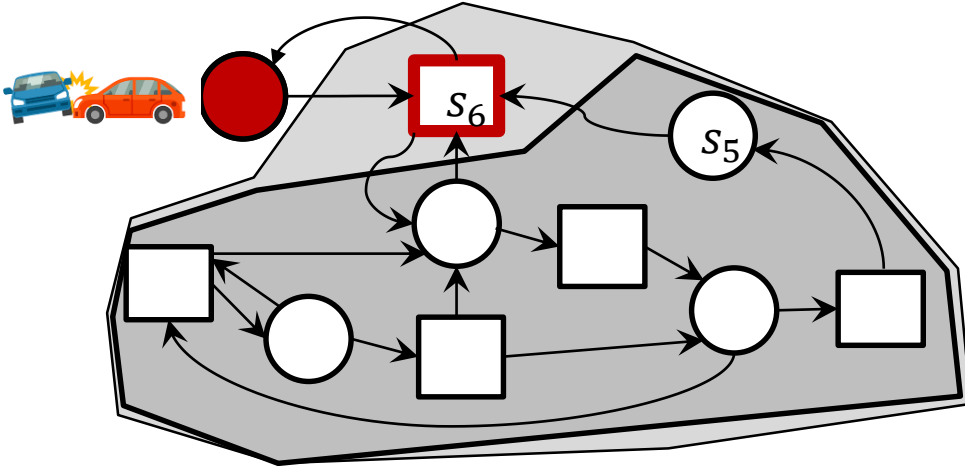
**System Player** wins, if  is **never** visited

**Winning Region:** States from which the system can enforce that  is **never** visited

# Synthesis is a Game



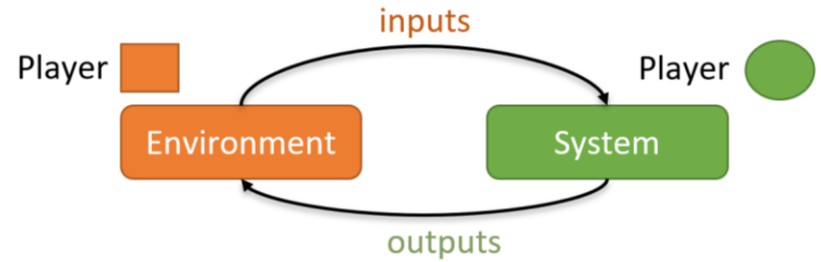
What is the winning region for this example?



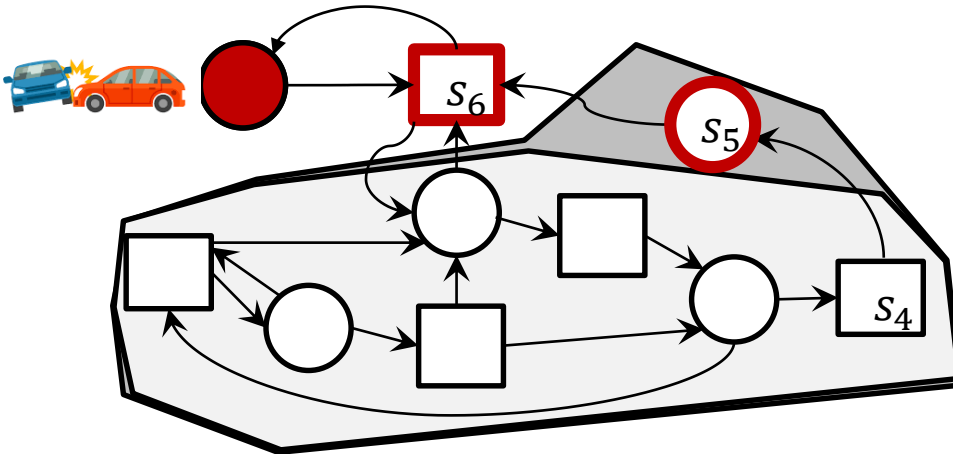
**System Player** wins, if is **never** visited

**Winning Region:** States from which the system can enforce that is **never** visited


# Synthesis is a Game



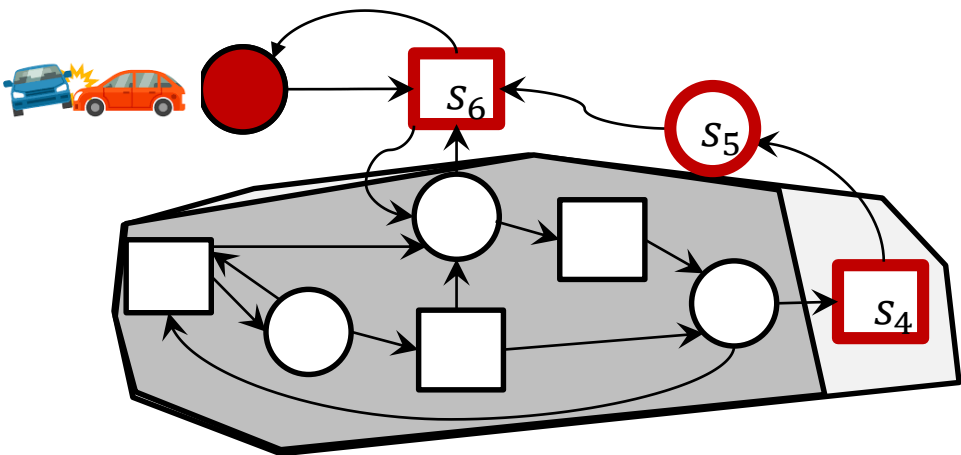
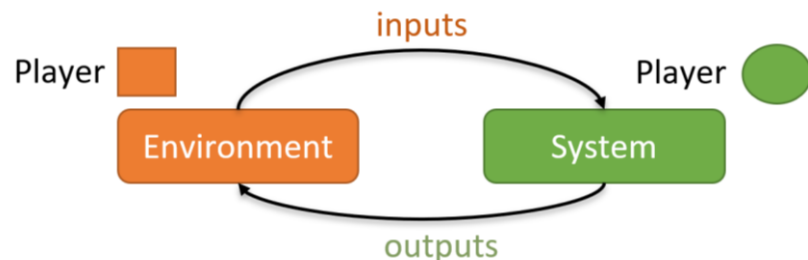
What is the winning region for this example?



**System Player** wins, if  is **never** visited

**Winning Region:** States from which the system can enforce that  is **never** visited

# Synthesis is a Game

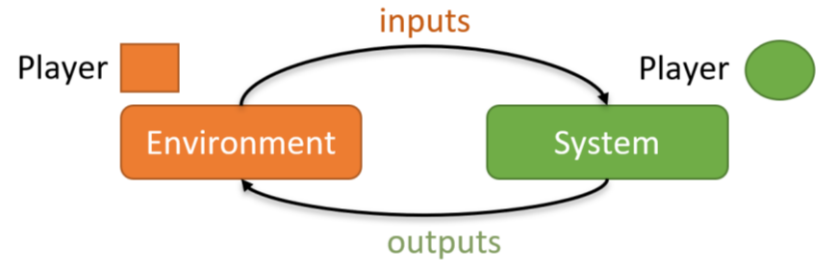


What is the winning region for this example?

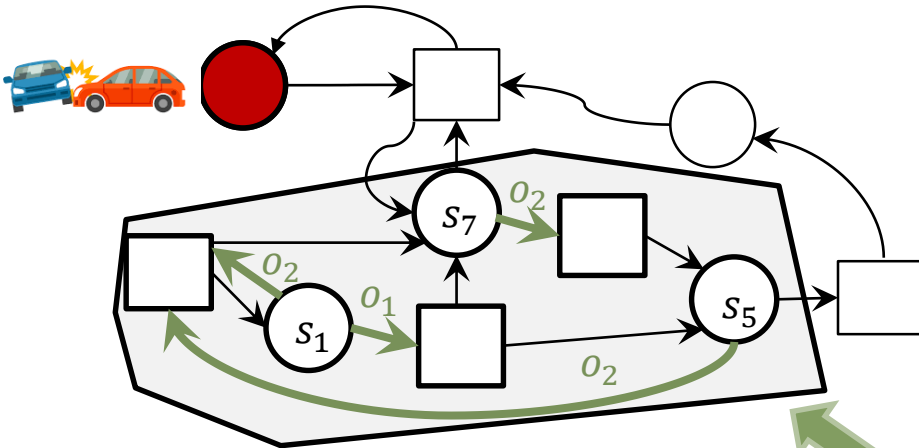
**System Player** wins, if is **never** visited

**Winning Region:** States from which the system can enforce that is **never** visited

# Synthesis is a Game



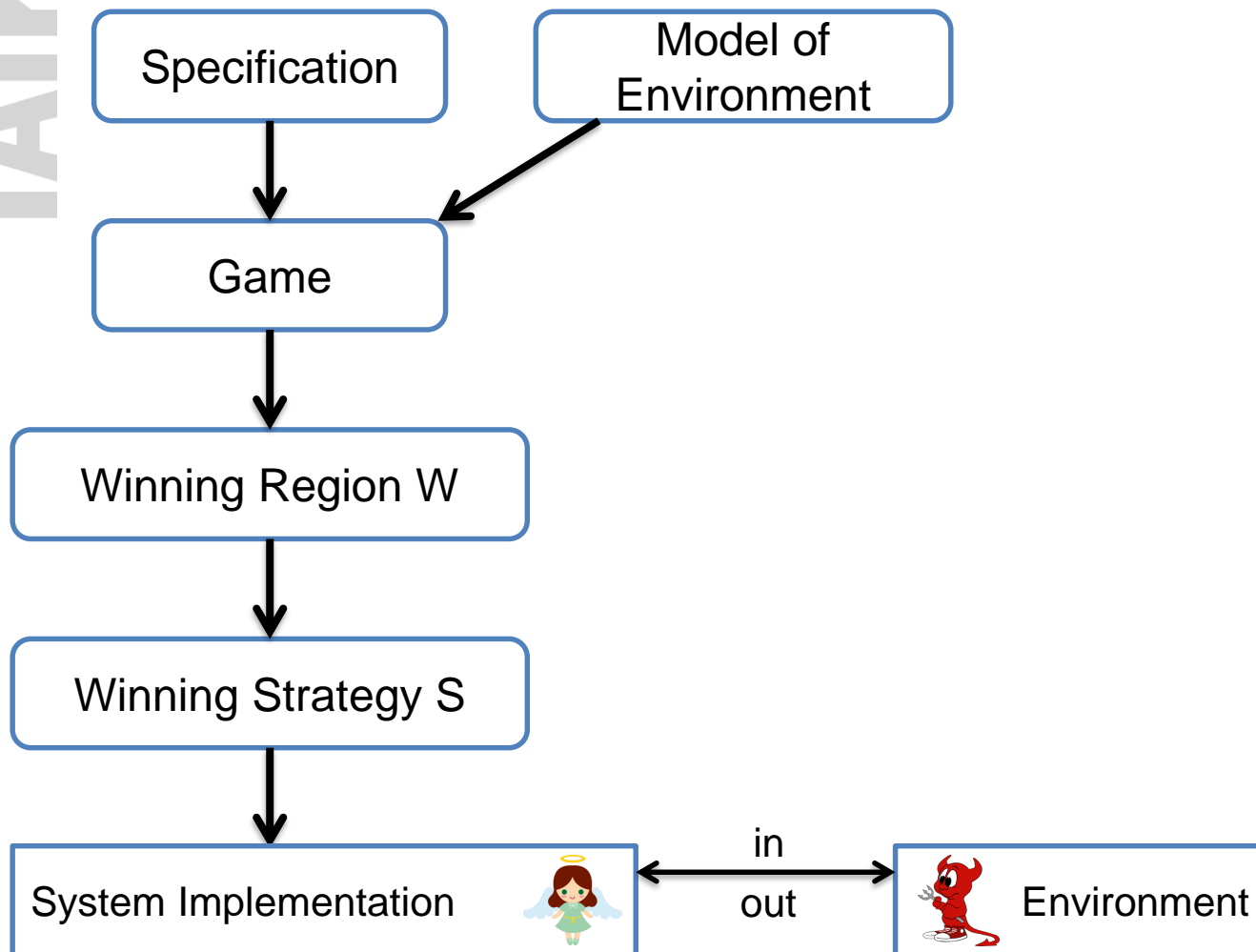
**What is the winning region for this example?**



**System Player** wins, if is **never** visited

**Winning Region:** States from which the system can enforce that is **never** visited

# Synthesis Flow



# Outline: Synthesis is a Game



- What is a game?
- Games on Graphs
- Solving Games



# What characterizes a game?

Games are **fun**

Several **players**

At least 2

## Games

Goal: **win** the game

Players can make **moves**

**Competitive:**

If one wins, the other loses

Moves must follow **rules**

Game has a **state**

You need a good **strategy** to win

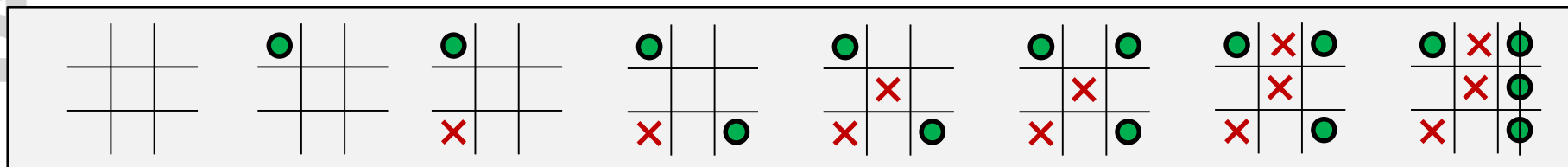
State is changed by moves

# Example: Tic-Tac-Toe

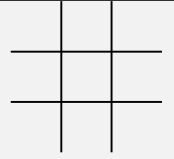
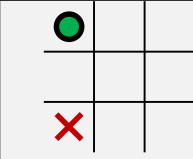
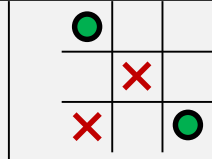
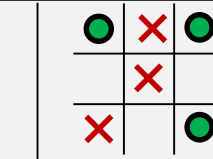
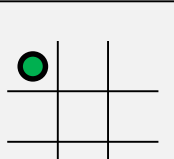
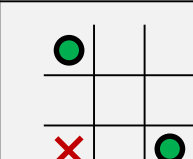
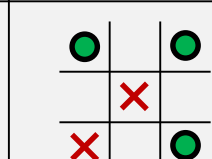
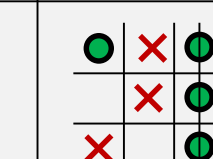
- 2 Players: ● and ×
- Players can make moves
  - ●-player: place ●
  - ×-player: place ×
- Rules for moves:
  - Players move in alternation
  - ... can only pick free slots
  - ...
- Winning condition: connect 3
- Players play **against** each other

# Terminology

- **Play:** Execution of the Game

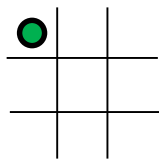


- **Strategy:** Defines what should be done when

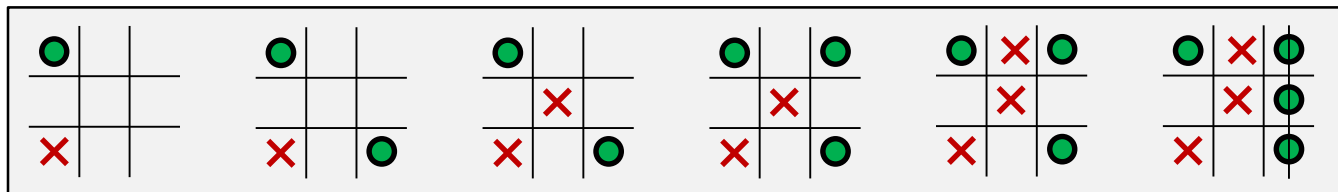
if:					and so on ...
then:					and so on ...

# Terminology

- Solving a game = finding a strategy that wins any game
  - No matter how the opponent plays
  
- Such a strategy is called a **Winning Strategy**
  - Not always possible
  - Computation needs to think ahead
  - Example: we are **x** and need to respond to

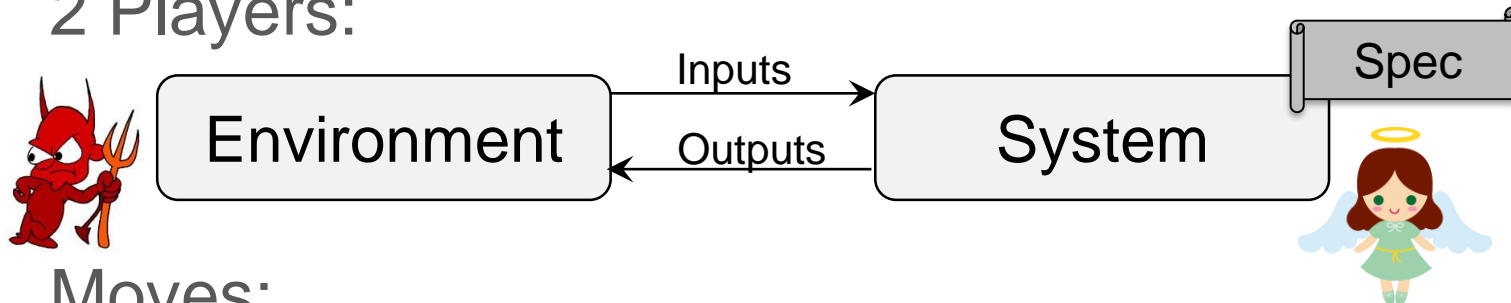


- Take the lower left corner  $\rightarrow$  dead in  $\leq 5$  moves:



# Synthesis is a Game

- 2 Players:



- Moves:

- Provide input values

Provide output values

- Winning Condition:

- Violate the specification

Satisfy the specification

- Strategy of System:

if:	state $s=q_7$ input $i=i_4$	$s=q_7$ $i=i_5$	$s=q_9$ $i=i_1$	and so on
then:	output $o=o_3$ next state $s=q_9$	$o=o_4$ $s=q_7$	$o=o_1$ $s=q_1$	and so on

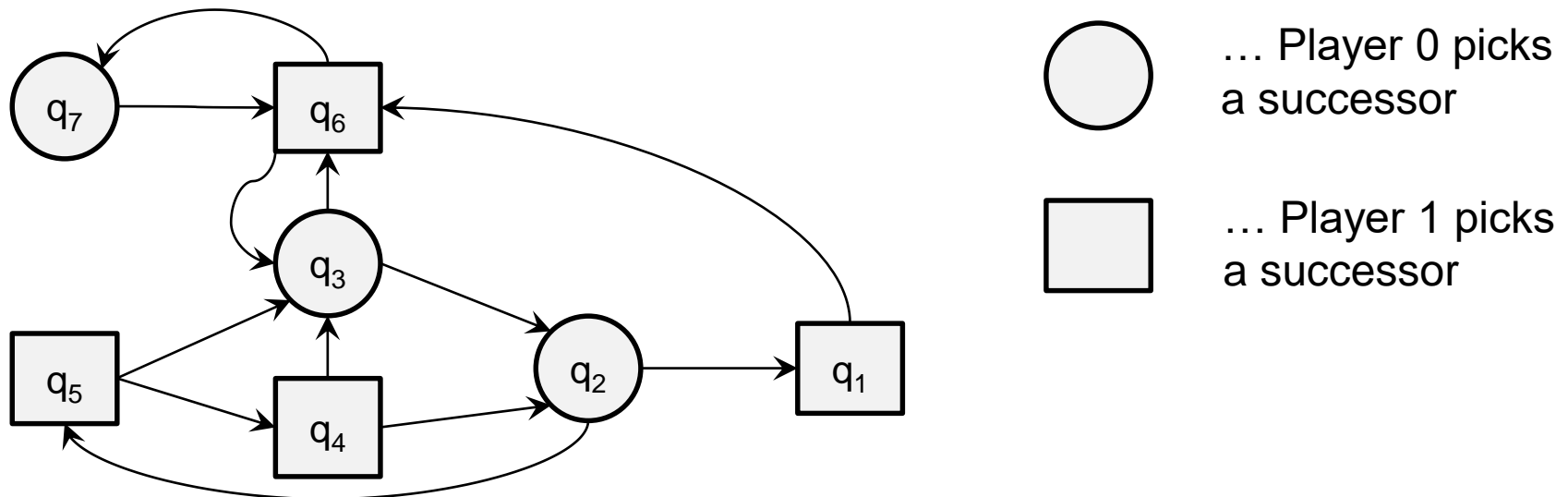
# Outline: Synthesis is a Game



- What is a game?
- Games on Graphs
- Game solving requires some math...
- Solving Games

# Game = Graph + Winning condition

- Game = Graph + Winning condition
- Game graph:  $G = (Q_0 \cup Q_1, E)$



## Play $\rho$ :

- Infinite sequence of states:  $\rho = q_0 q_1 q_2 \dots \in Q^\omega$

# Games on Graphs

- Game = Graph + Winning condition
- Game graph  $G = (Q, E)$ :
  - Every state has an outgoing edge:
    - $\forall q \in Q: \exists q' \in Q: (q, q') \in E$
  - $Q$  is partitioned into  $Q_0$  and  $Q_1$ :
    - $Q = Q_0 \cup Q_1$  with  $Q_0 \cap Q_1 = \emptyset$
    - Player 0 picks a successor state in  $Q_0$
    - Player 1 picks a successor state in  $Q_1$



# Game = Game Graph + Winning Condition

- **Winning condition**  $\varphi: Q^\omega \rightarrow \mathbb{B}$ 
  - $\rho$  is won by the **Player 0** iff  $\varphi(\rho) = \top$
  - $\rho$  is won by the **Player 1** iff  $\varphi(\rho) = \perp$

- **Types of winning conditions**

- Let  $F \subseteq Q$  be a set of states



1. Reach  $F$  at least once
2. Stay in  $F$  forever
3. Reach  $F$  infinitely often

- Safety Games
- Büchi Games
- Reachability Game

# Game = Game Graph + Winning Condition

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- **Types of winning conditions**
  - Let  $F \subseteq Q$  be a set of states
 

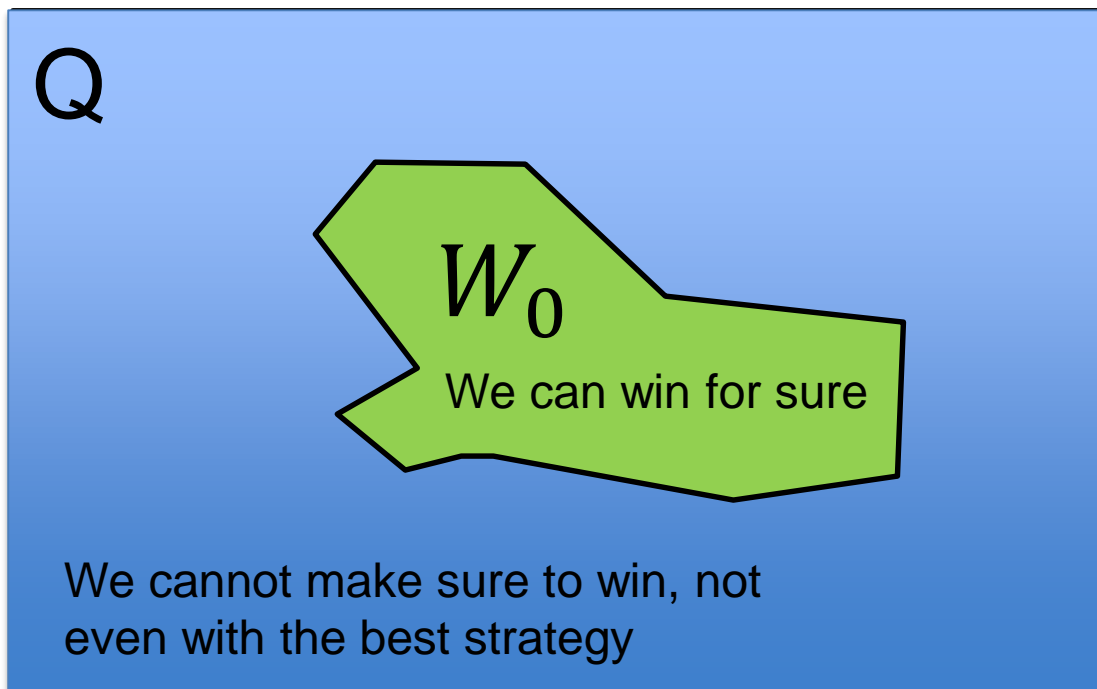
<ol style="list-style-type: none"> <li>1. Reach <math>F</math> at least once</li> <li>2. Stay in <math>F</math> forever</li> <li>3. Reach <math>F</math> infinitely often</li> </ol>	<ol style="list-style-type: none"> <li><span style="border: 1px solid black; padding: 2px;">2</span> Safety Games</li> <li><span style="border: 1px solid black; padding: 2px;">3</span> Büchi Games</li> <li><span style="border: 1px solid black; padding: 2px;">1</span> Reachability Game</li> </ol>
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# Winning Strategy

- Positional Strategy:
  - $f^0: Q_0 \rightarrow Q$
  - A play  $\rho = q_0q_1 \dots$  **follows** positional strategy  $f^0: Q_0 \rightarrow Q$  iff  $q_{i+1} = f^0(q_i)$  for all  $q_i \in Q_0$
- Winning Strategy:
  - Makes sure that P0 **always** wins

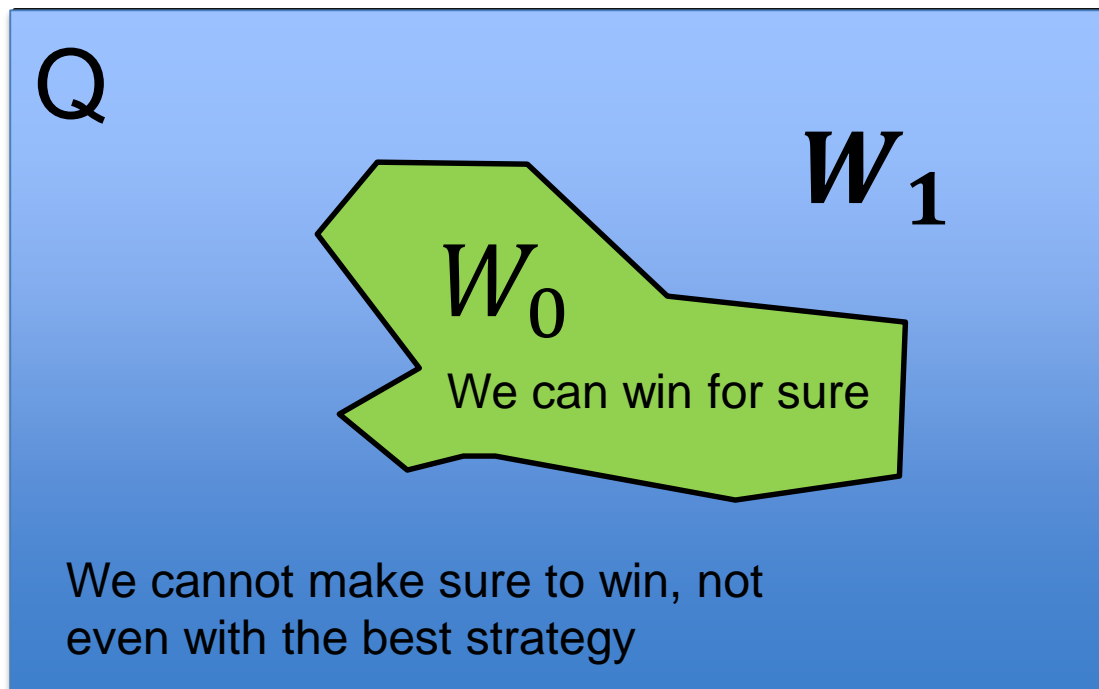
# Winning Region

- $W_0$  is the set of states from which a winning strategy exists
- ToDo What is the winning region  $W_1$  of P1?



# Winning Region

- $W_0$  is the set of states from which a winning strategy exists
- What is the winning region  $W_1$  of P1?



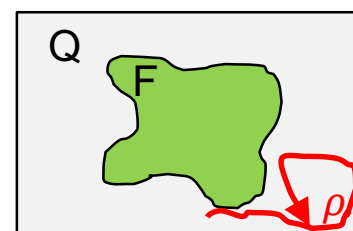
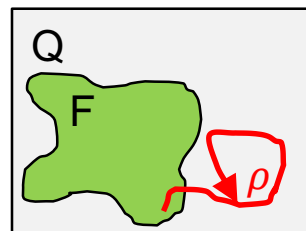
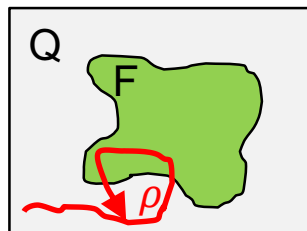
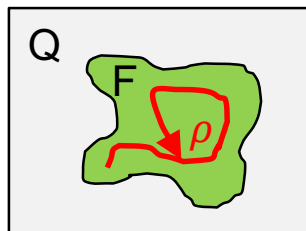
Winning Conditions:

# Reachability Games

- $\varphi$  is defined using a set **F** of “target states”
- Player 0 wins a play  $\rho = q_0q_1 \dots$  iff **F** is visited
- $\varphi(\rho) \Leftrightarrow \exists i: q_i \in F$



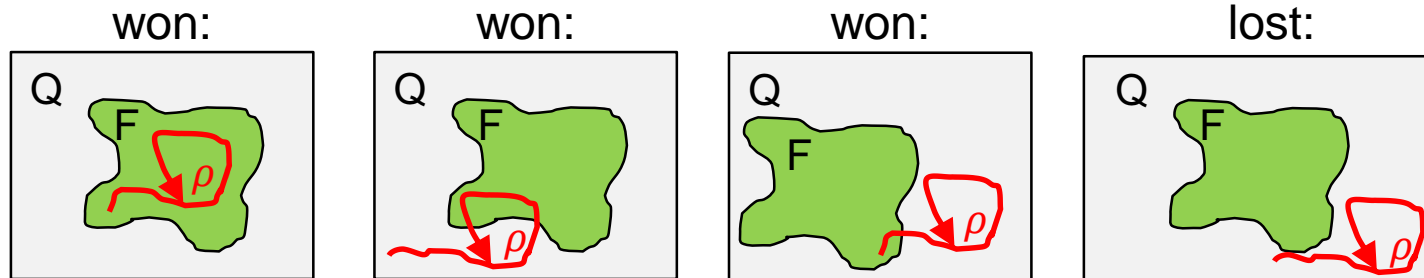
Which plays are winning plays for P0?



Winning Conditions:

# Reachability Games

- $\varphi$  is defined using a set **F** of “target states”
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Winning Conditions:

# Safety Games

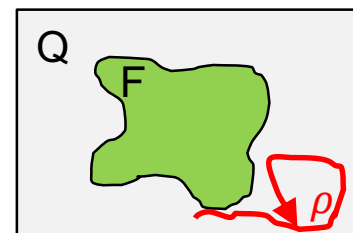
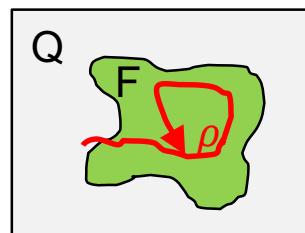
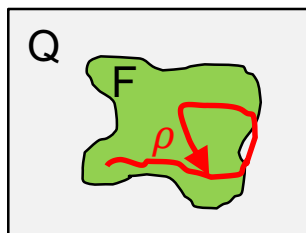
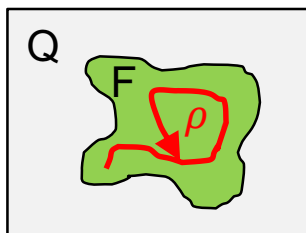
- $\varphi$  is defined using a set **F** of “safe states”
- Player 0 wins a play  $\rho = q_0 q_1 \dots$  iff **it stays in F**
- $\varphi(\rho) \Leftrightarrow \forall i: q_i \in F$



# Winning Conditions: Safety Games

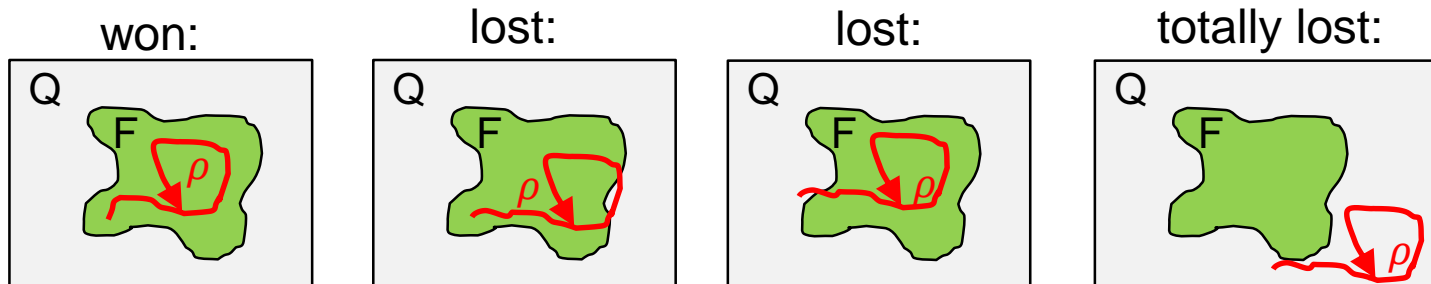
- $\varphi$  is defined using a set **F** of “safe states”
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 Which plays are winning plays for P0?



# Winning Conditions: Safety Games

- $\varphi$  is defined using a set **F** of “safe states”
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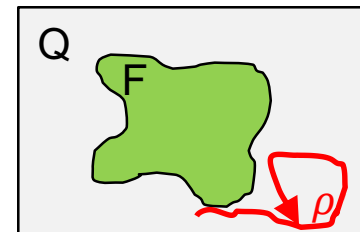
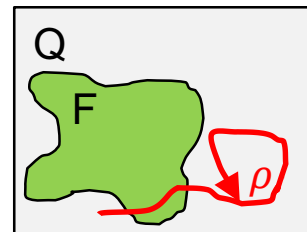
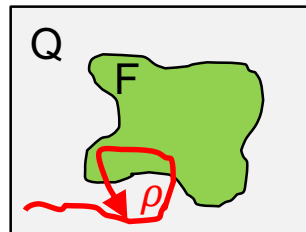
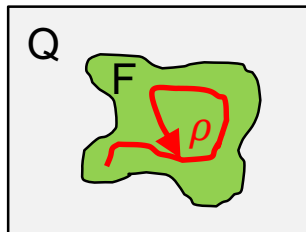


Winning Conditions:  
**Büchi Games**

- $\varphi$  is defined using a set **F** of “accepting states”
- Player 0 wins a play  $\rho$  iff **F** is visited infinitely often

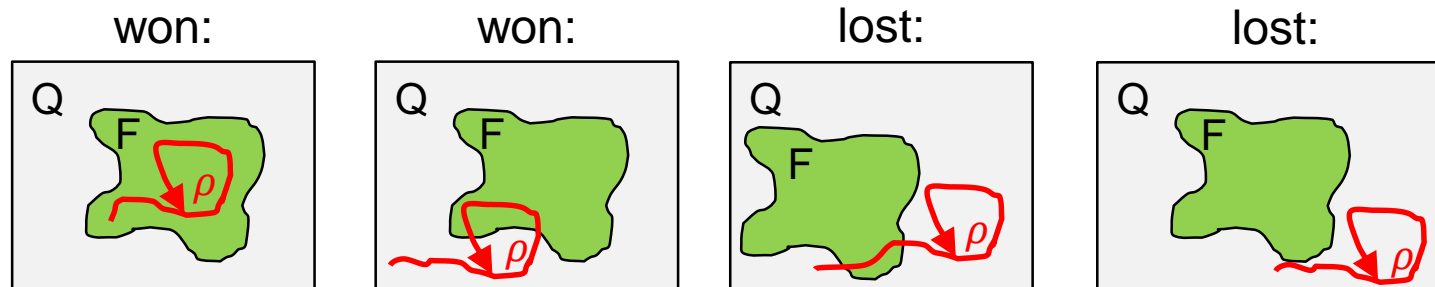
# Winning Conditions: Büchi Games

- $\varphi$  is defined using a set **F** of “accepting states”
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# Winning Conditions: Büchi Games

- $\varphi$  is defined using a set **F** of “accepting states”
- Player 0 wins a play  $\rho$  iff **F** is visited infinitely often



- $\text{Inf}(\rho)$ : the states occurring infinitely often in  $\rho$
- $\varphi(\rho) \Leftrightarrow \text{Inf}(\rho) \cap F \neq \emptyset$

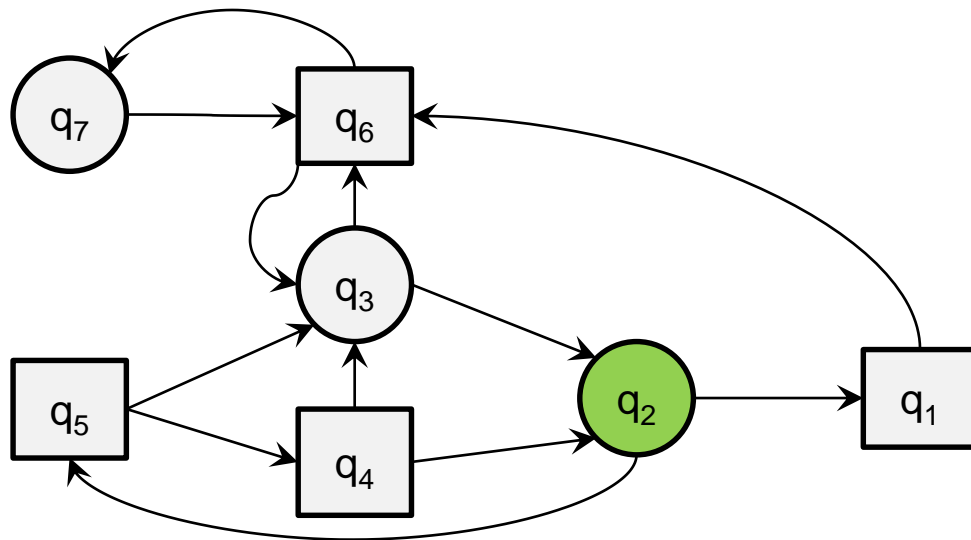
# Outline: Synthesis is a Game



- What is a game?
- Games on Graphs
- Solving Games
  - Reachability
  - Safety

# Reachability Game

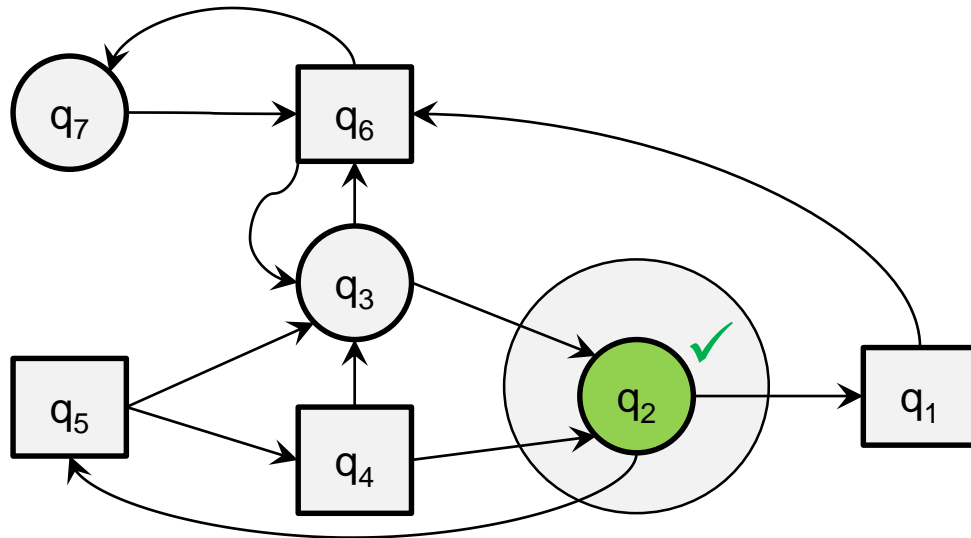
- Compute the winning region of Player 0 for:



- For the reachability game with  $F = \{q_2\}$

# Reachability Game

- Compute the winning region of Player 0 for:



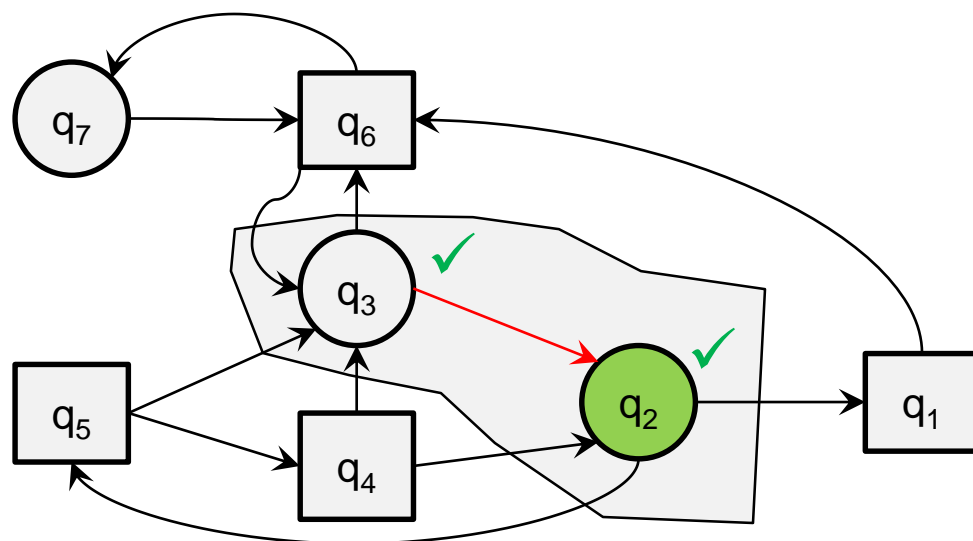
Every strategy is winning from  $q_2$   
 $\rightarrow$   
 $q_2$  is part of the winning region

- For the reachability game with  $F = \{q_2\}$



# Reachability Game

- Compute the winning region of Player 0 for:



If we start from  $q_3$ , we can get into the winning region (go from  $q_3$  to  $q_2$ , not  $q_6$ )

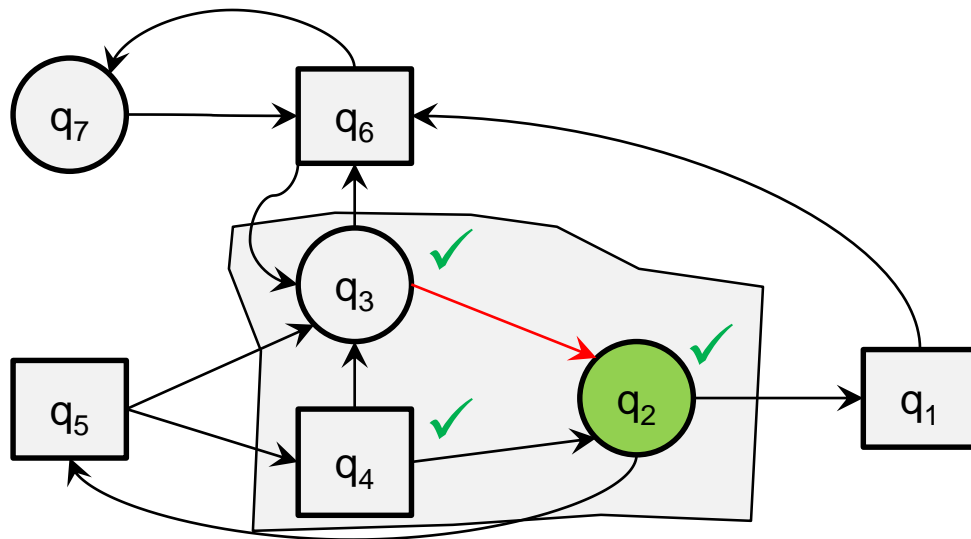


$q_3$  is part of the winning region

- For the reachability game with  $F = \{q_2\}$

# Reachability Game

- Compute the winning region of Player 0 for:

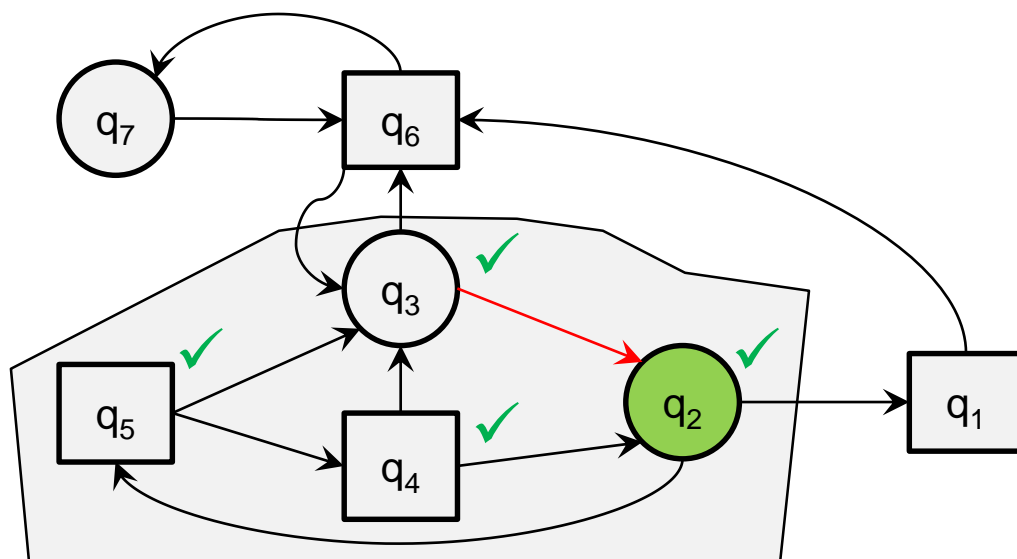


If we start from  $q_4$ , we must end up in the winning region  
 →  
 $q_4$  is part of the winning region

- For the reachability game with  $F = \{q_2\}$

# Reachability Game

- Compute the winning region of Player 0 for:



If we start from  $q_5$ , we must end up in the winning region

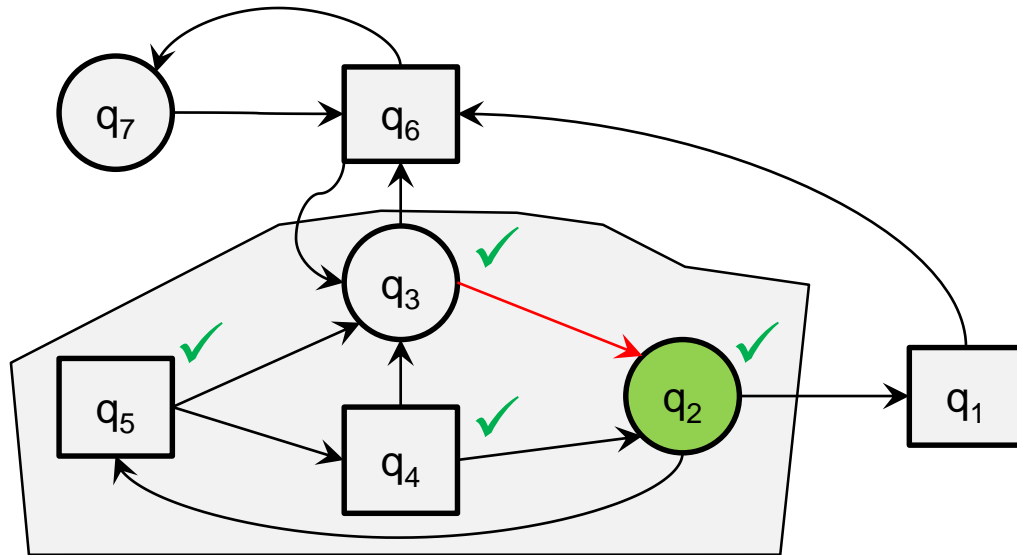


$q_5$  is part of the winning region

- For the reachability game with  $F = \{q_2\}$

# Reachability Game

- Compute the winning region of Player 0 for:



We cannot enforce going into the winning region from other states



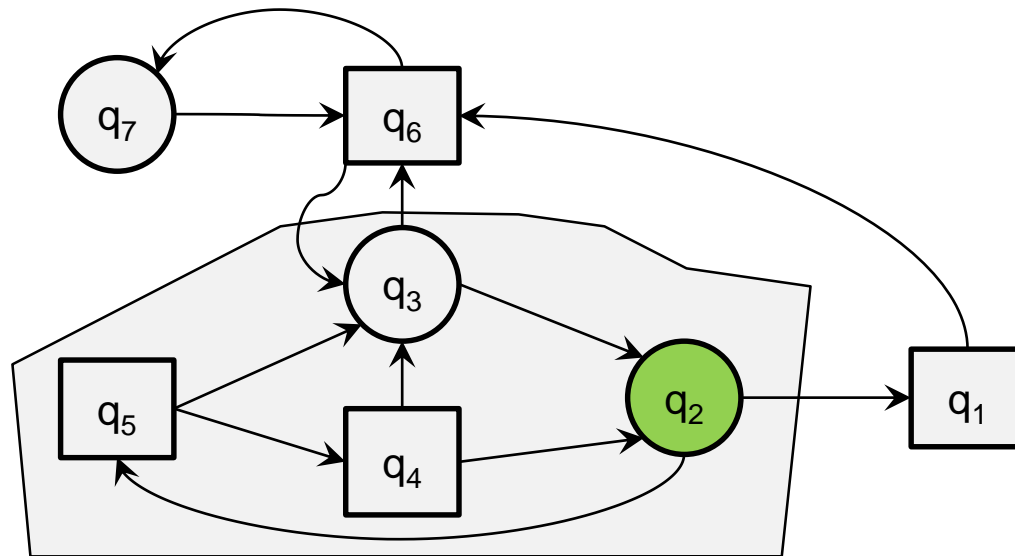
We are done! The winning region is  $\{q_2, q_3, q_4, q_5\}$

- For the reachability game with  $F = \{q_2\}$

# Reachability Game

## Winning Strategy

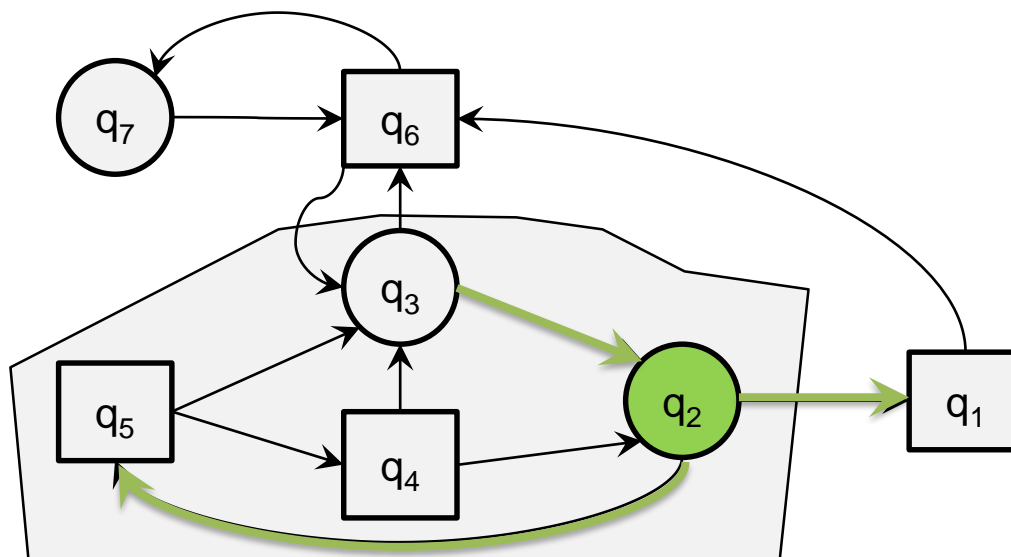
- Winning Strategy for Player 0?



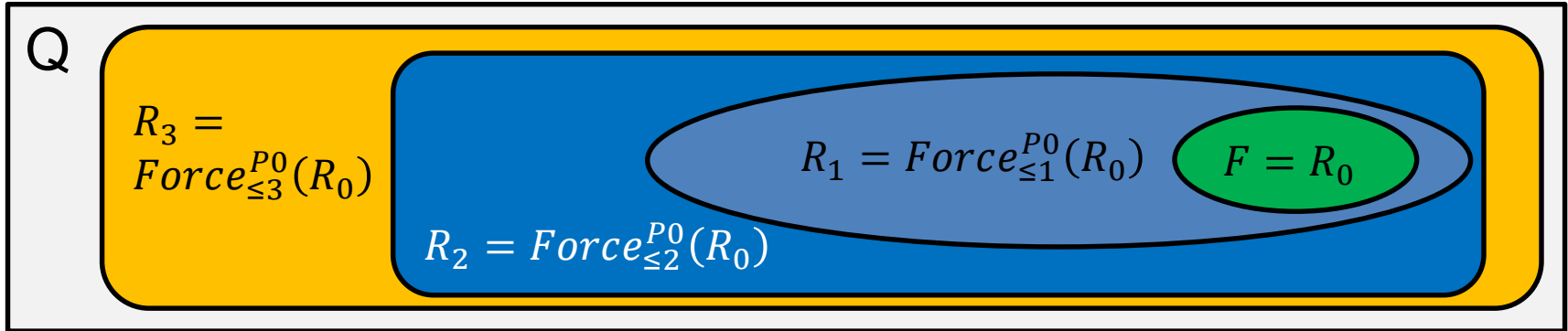
# Reachability Game

## Winning Strategy

- Winning Strategy for Player 0



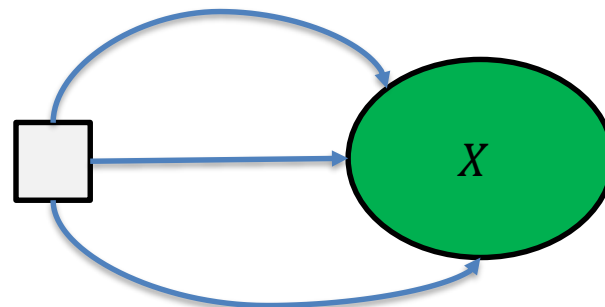
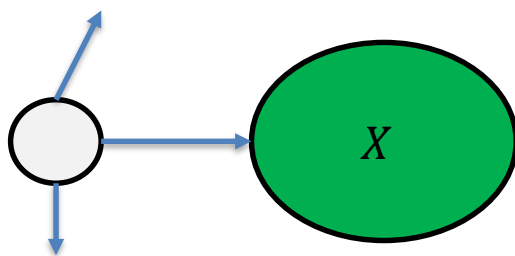
# Reachability Game: Winning Region



- $R_i(F) = \{q \in Q \mid \text{Player 0 can enforce to visit } F \text{ in } \leq i \text{ steps}\}$
- $R_\infty(F) = W = \{q \in Q \mid \text{Player 0 can enforce to visit } F\}$

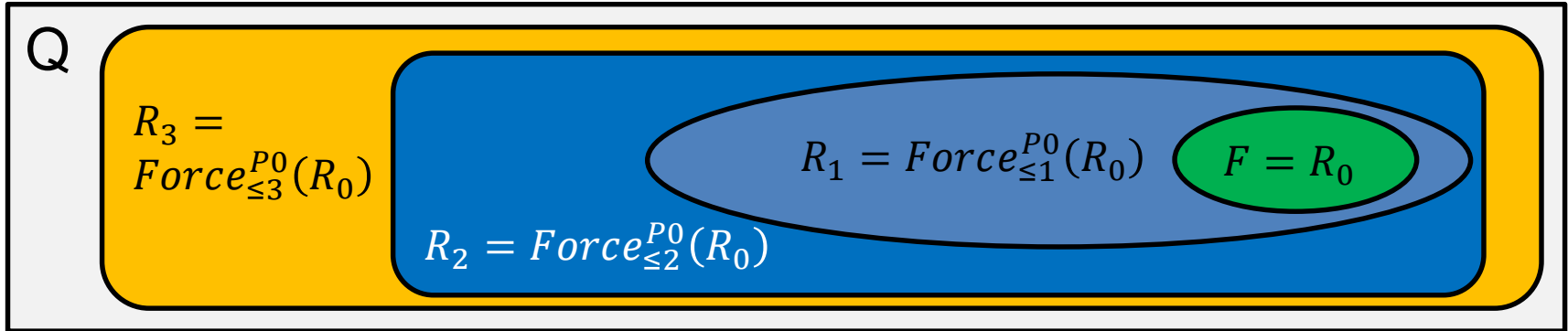
## Reachability Game: Winning Region

- $\text{Force}_1^{P0}(X) = \{q \in Q \mid \text{Player 0 can force to reach } X \text{ in exactly 1 step}\}$
- $\text{Force}_1^{P0}(X) = \{q \in Q_0 \mid \exists (q, q') \in E: q' \in X\} \cup \{q \in Q_1 \mid \forall (q, q') \in E: q' \in X\}$



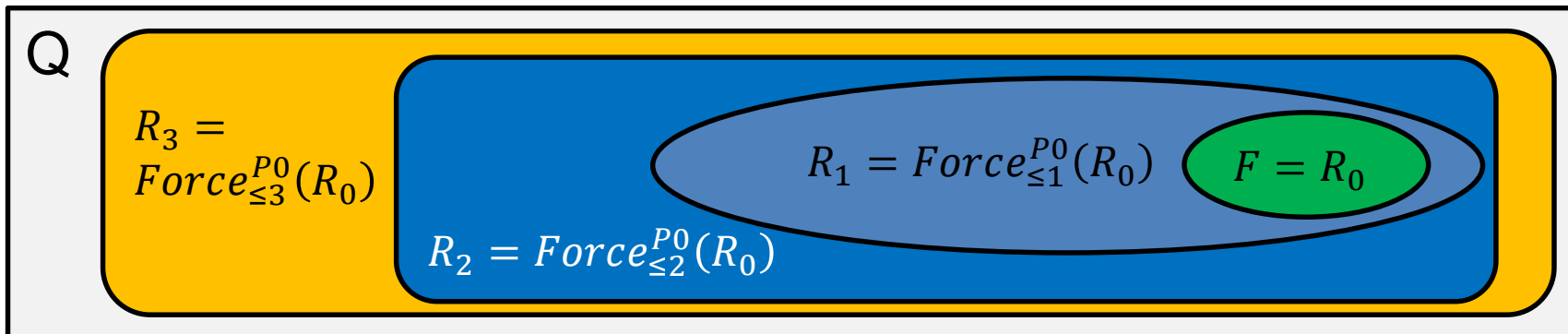


# Reachability Game: Winning Region



- $R_i(F) = \{q \in Q \mid \text{Player 0 can enforce to visit } F \text{ in } \leq i \text{ steps}\}$
- $R_\infty(F) = W = \{q \in Q \mid \text{Player 0 can enforce to visit } F\}$
- Algorithm to compute  $W$

# Reachability Game: Winning Region



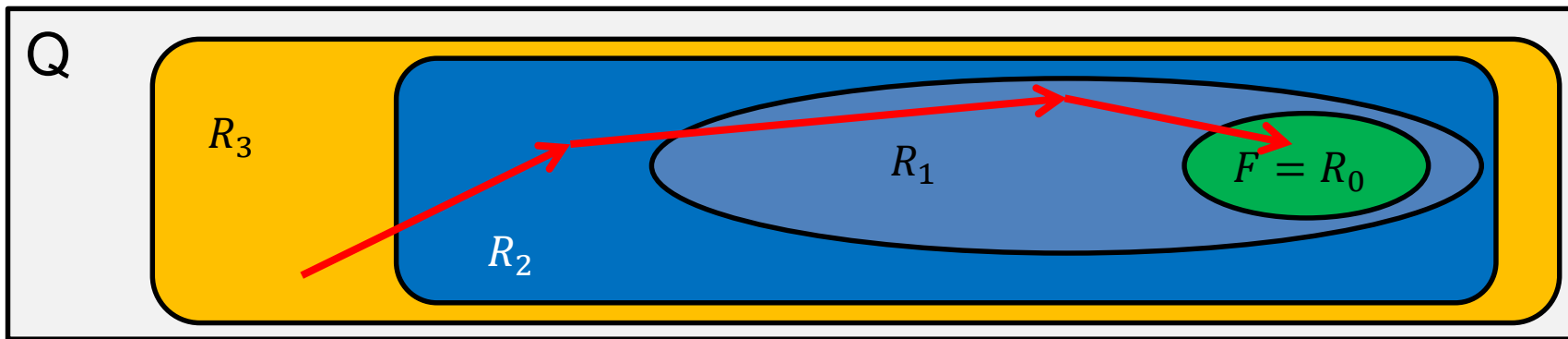
- $R_i(F) = \{q \in Q \mid \text{Player 0 can enforce to visit } F \text{ in } \leq i \text{ steps}\}$
- $R_\infty(F) = W = \{q \in Q \mid \text{Player 0 can enforce to visit } F\}$
- Algorithm to compute  $W$

```

W {
  R = {F}
  while (R changes)
    R = R ∪ Force1P0(R)
  return R
}

```

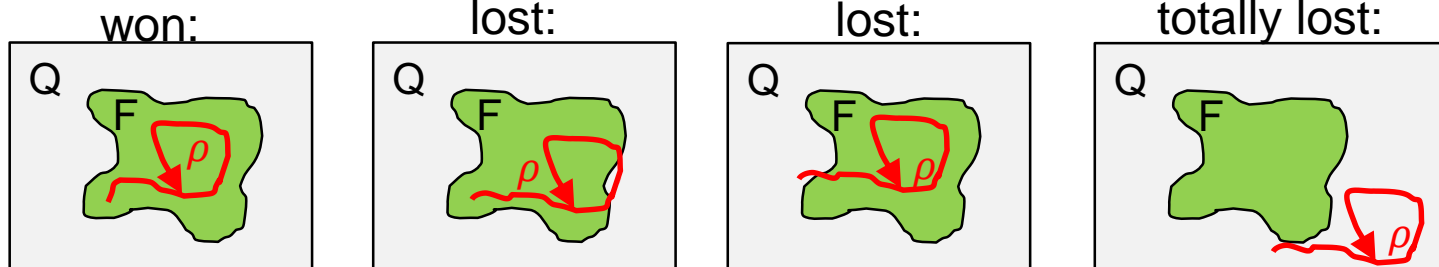
# Reachability Game: Winning Strategy



From  $R_i \setminus R_{i-1}$  go to  $R_{i-1}$

# Winning Conditions: Safety Games

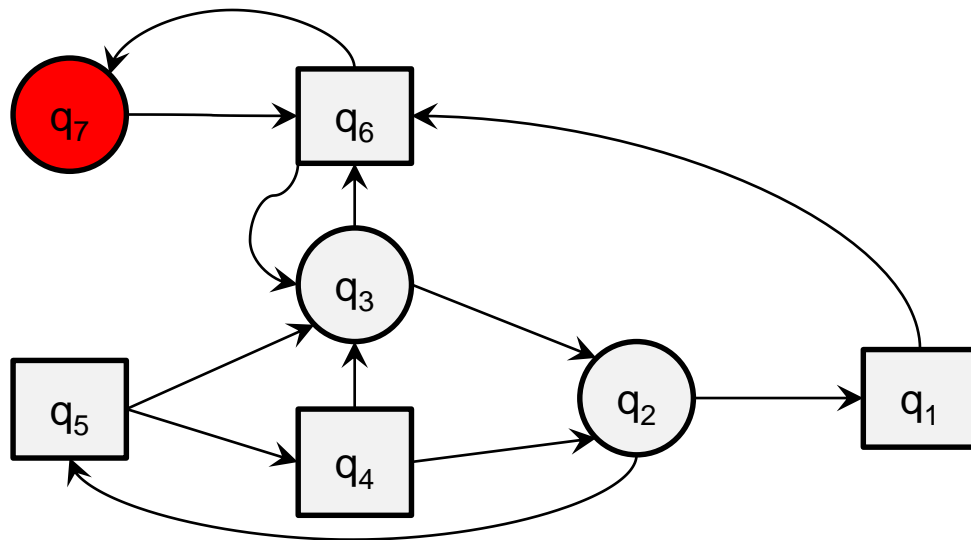
- set **F** of “safe states”
- Player 0 wins a play iff it **stays in F**



# Safety Game



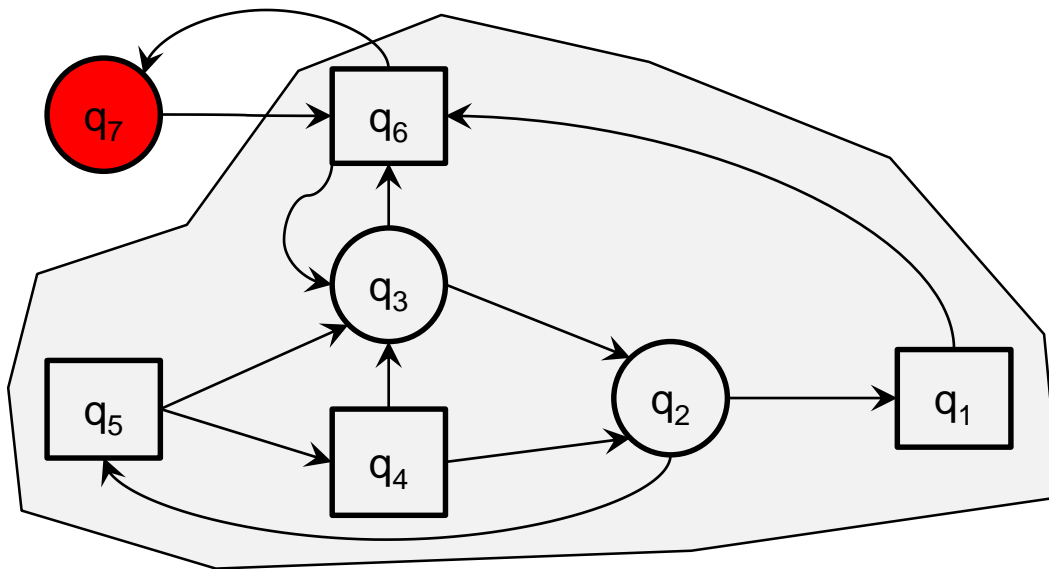
- Compute the winning region of Player 0



- For the safety game with  $F = \{q_1, q_2, q_3, q_4, q_5, q_6\}$

# Safety Game

- Compute the winning region of Player 0

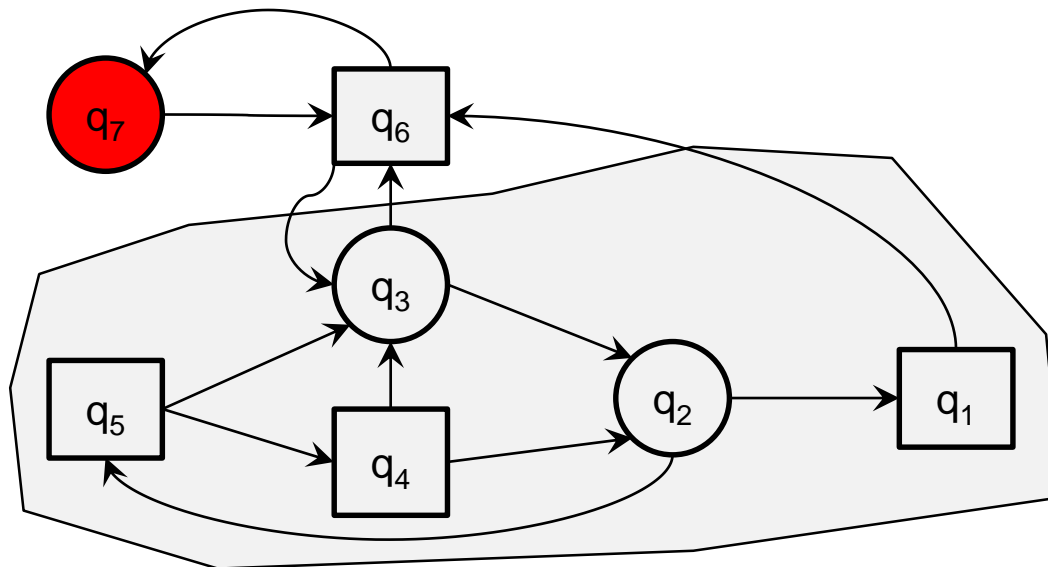


Let's start with all safe states and see from which states we can fall out of the safe region.

- For the safety game with  $F = \{q_1, q_2, q_3, q_4, q_5, q_6\}$

# Exercise: Safety Game

- Compute the winning region of Player 0

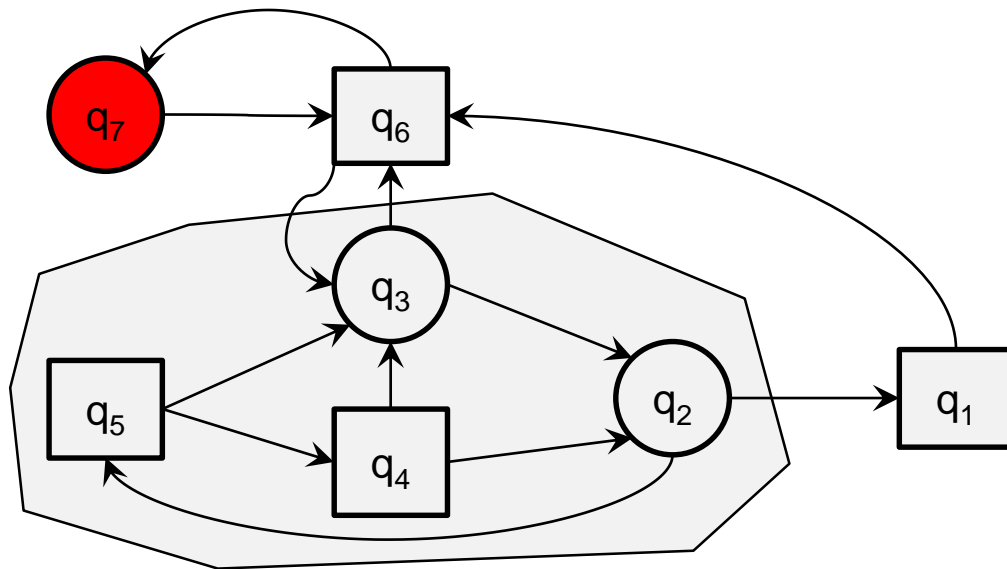


If we ever get to  $q_6$  we are in trouble  
 →  
 $q_6$  is not in the winning region

- For the safety game with  $F = \{q_1, q_2, q_3, q_4, q_5, q_6\}$

# Exercise: Safety Game

- Compute the winning region of Player 0



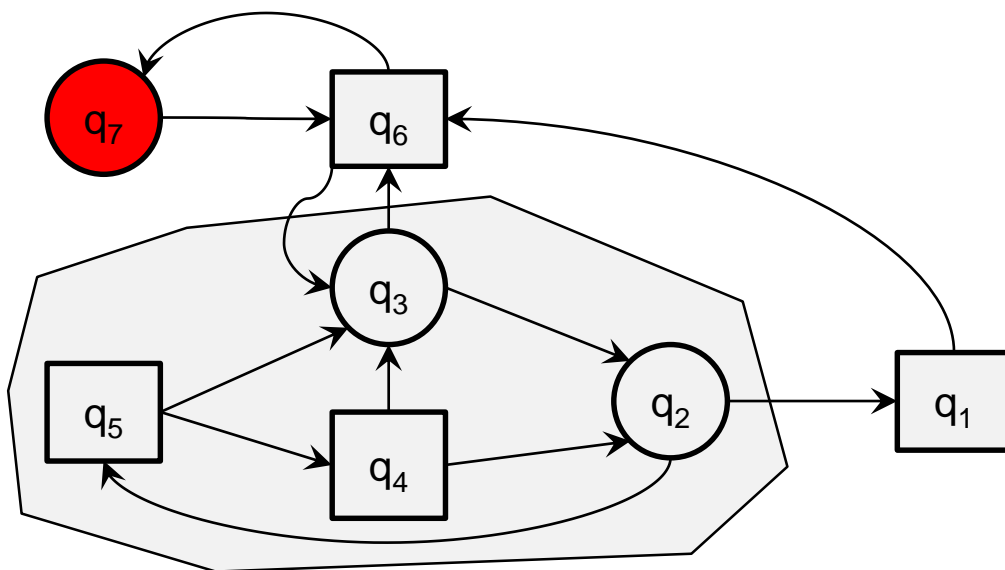
If we ever get to  $q_1$  we are in trouble  
 →  
 $q_1$  is not in the winning region

- For the safety game with  $F = \{q_1, q_2, q_3, q_4, q_5, q_6\}$



# Exercise: Safety Game

- Compute the winning region of Player 0



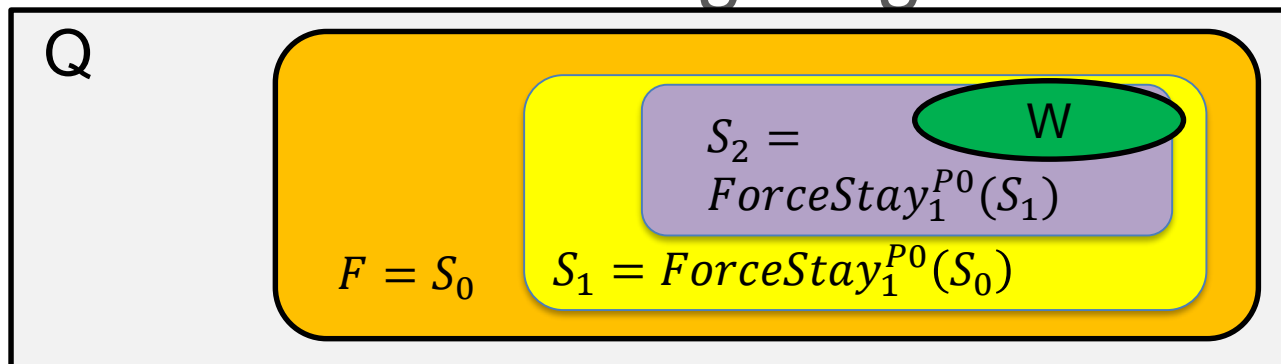
Now we can make sure that the grey area is not left

➔

We are done! The winning region is  $\{q_2, q_3, q_4, q_5\}$

- For the safety game with  $F = \{q_1, q_2, q_3, q_4, q_5, q_6\}$

## Safety Game: Winning Region



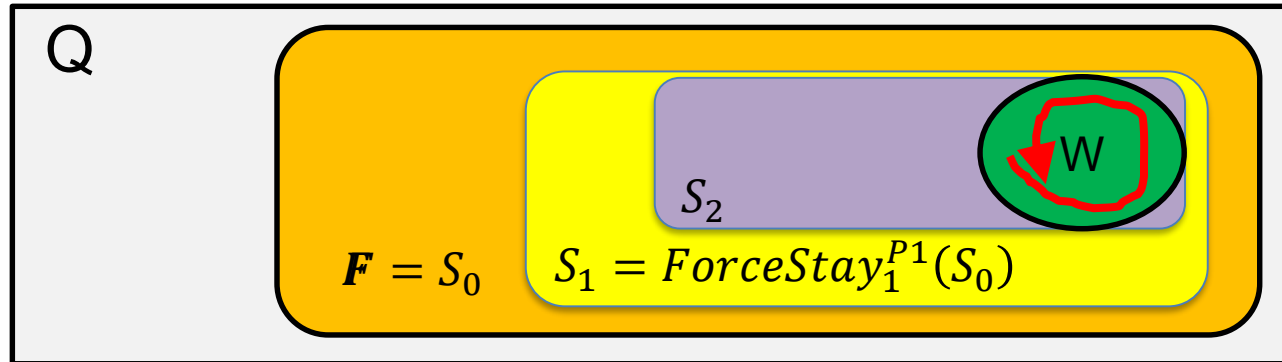
- $S_i(X) = \{q \in Q \mid \text{Player 0 can enforce to stay in } F \text{ for } \geq i \text{ steps}\}$
- $S_\infty(X) = W = \{q \in Q \mid \text{Player 0 can enforce to stay in } F \text{ forever}\}$
- Algorithm to compute  $W$

```

W {
  S = F
  while (S changes)
    S = F  $\cap$  Force $_1^{P_0}$ (S)
  return S
}
    
```

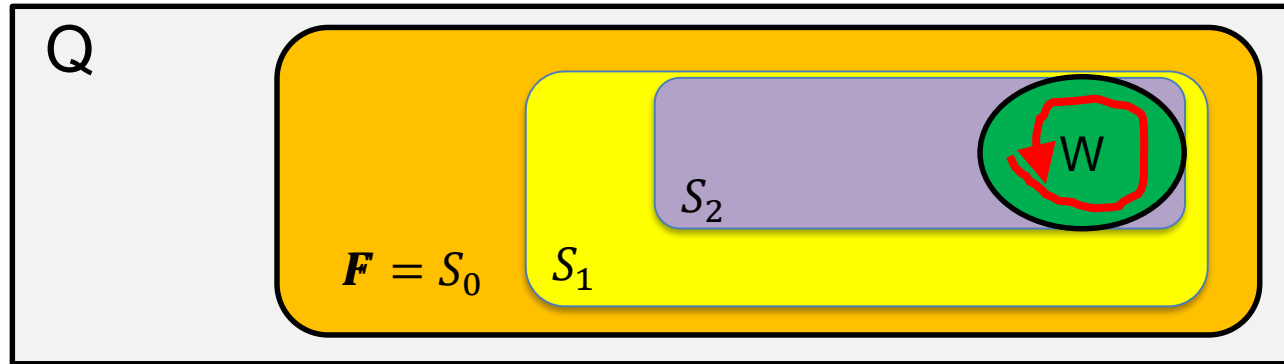
Safety Game:

## Winning Strategy for Player 0

From  $W$  go to  $W$ 


- In  $Q_0$  states: pick one such edge
- In  $Q_1$  states: Possible because  $W$  is constructed in such a way

# Safety Game: Winning Strategy For Player 0

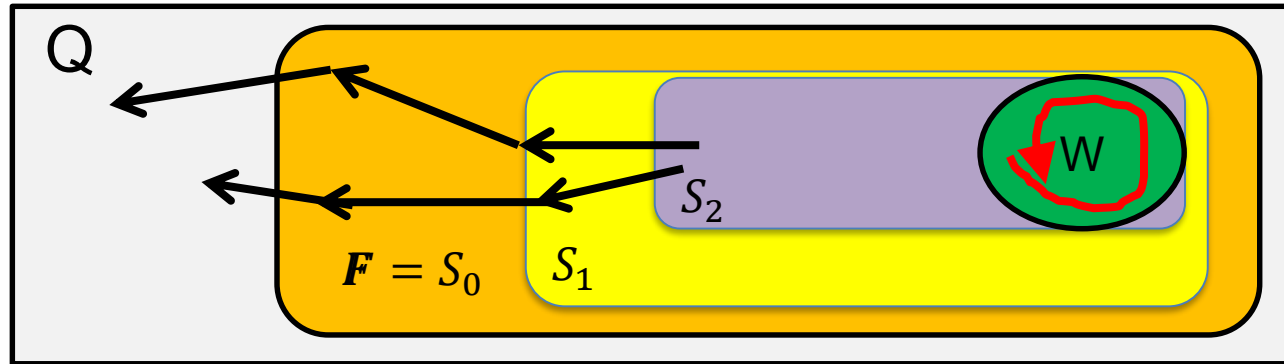




 Player 0 Strategy

 Player 1 Strategy

- 
  - In which states exists a winning strategy for player 1?
  - Which game is P1 playing?
  - What is the strategy of P1?

# Safety Game: Winning Strategy For Player 0



 Player 0 Strategy  
 Player 1 Strategy

- - In which states exists a winning strategy for player 1?
    - $W_1 = Q \setminus W_0$
  - Which game is P1 playing?
    - Reachability Game
  - What is the strategy of P1?

