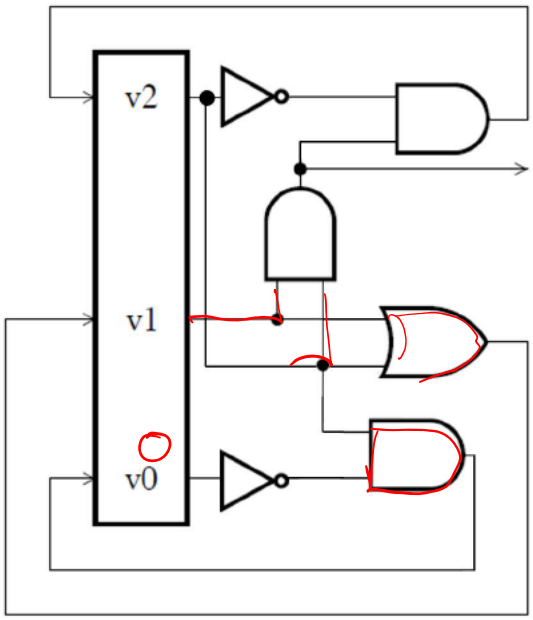


$V = \{v_0, v_1, v_2\}$

$V' = \{v_0', v_1', v_2'\}$

The initial value of the state variable v_0 of the circuit is $false$. The initial values of v_1 and v_2 are unknown.

Task 1a. [5 Points]

- State the formula S_0 that represents the set of initial states and the formula R that represents the transition relation of C.

Task 1b. [5 Points]

- Draw the Kripke Structure $M = (S, S_0, R, AP, L)$ that represents C.

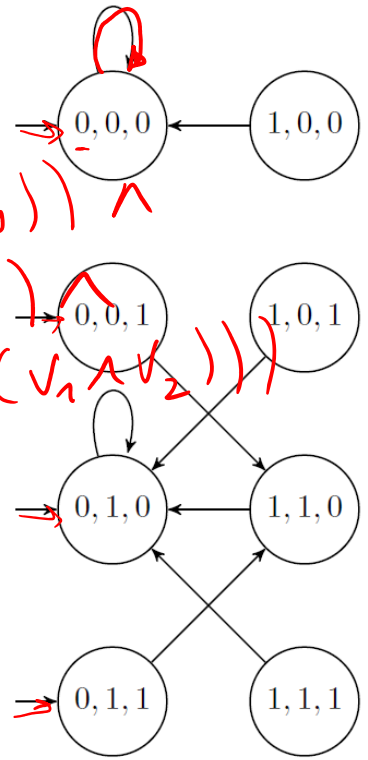
$S_0(v_0, v_1, v_2) = \neg v_0$

$R(V, V') = (v_0' \leftrightarrow (v_2 \wedge \neg v_0)) \wedge$

$(v_1' \leftrightarrow (v_1 \vee v_2)) \wedge$

$(v_2' \leftrightarrow (\neg v_2 \wedge (v_1 \wedge v_2)))$

v0, v1, v2



taken from an anonymous student

AGp
p is always True

EFp
there is a path
to p.

SAT-Based Model Checking

Chapter 10

Outline

- Bounded Model Checking
- Verifying Reachability Properties with k-induction
- Model Checking with Inductive Invariants
- Model Checking with Craig Interpolants
- Property-Directed Reachability

Overview

- SAT solvers can solve propositional formulas. (See Chapter 9.)
- SAT solvers have become very fast

Techniques

- Bounded Model Checking: Is there a trace of length k that violates the property?
- k-induction: Can we prove the property inductively for *any* trace?
- Create an inductive invariant that is stronger than the property
  - Craig Interpolants
  - Property-Directed Reachability

In this chapter, we will only consider LTL formulas invariants.

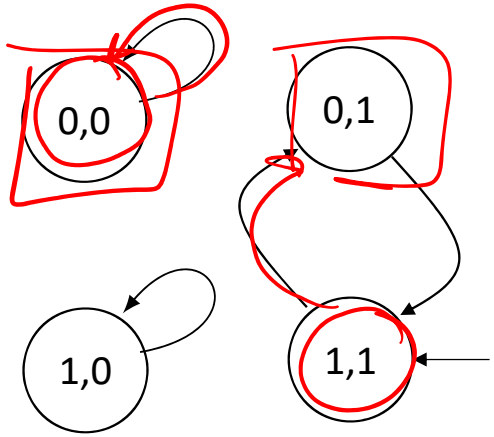
Preliminaries

Postimage

postimg(S) = states reachable from S in one step

postimg({ (0,1) }) = { (1,1) }

postimg({ (0,0), (1,1)}) = { (0,0), (0,1) }



Kripke structure

Paths

Initial states: $S_0(x, y) = (x = 1 \wedge y = 1)$

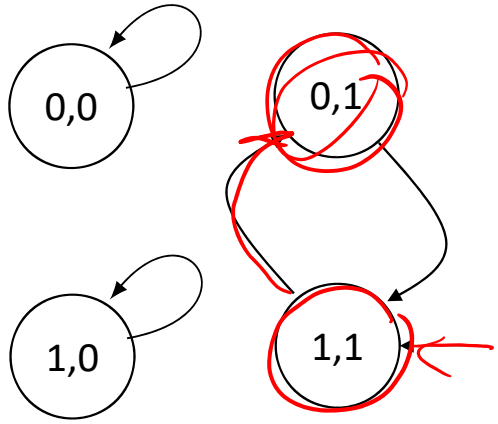
Transitions: $\mathcal{R}(x, y, x', y') = (x' = (x + y) \bmod 2) \wedge (y' = y)$

$Path_1(S(V)) = S \wedge R(V, V') =$

$\exists V. \neg x \wedge y \wedge (x' = (x + y) \bmod 2) \wedge (y' = y) =$

$x \wedge y \wedge (x' = (x - y \bmod 2)) \wedge (y' = y)$

$x = 1, y = 1 \quad x' = 0, y' = 1.$



Kripke structure

Paths

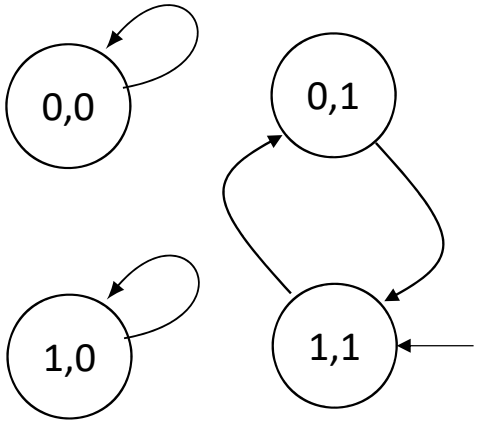
Initial states: $\mathcal{S}_0(x, y) = x = 1 \wedge y = 1$

Transitions: $\mathcal{R}\ (x, y, x', y') = (x' = (x + y)\ mod\ 2) \wedge (y' = y)$

$Path_1(S(V)) = S \wedge R(V, V') =.$

$\neg x \wedge y \wedge (x' = (x + y)\ mod\ 2) \wedge (y' = y) =.$

$\neg x \wedge y \wedge x' \wedge y'.$



Kripke structure

Bounded Model Checking

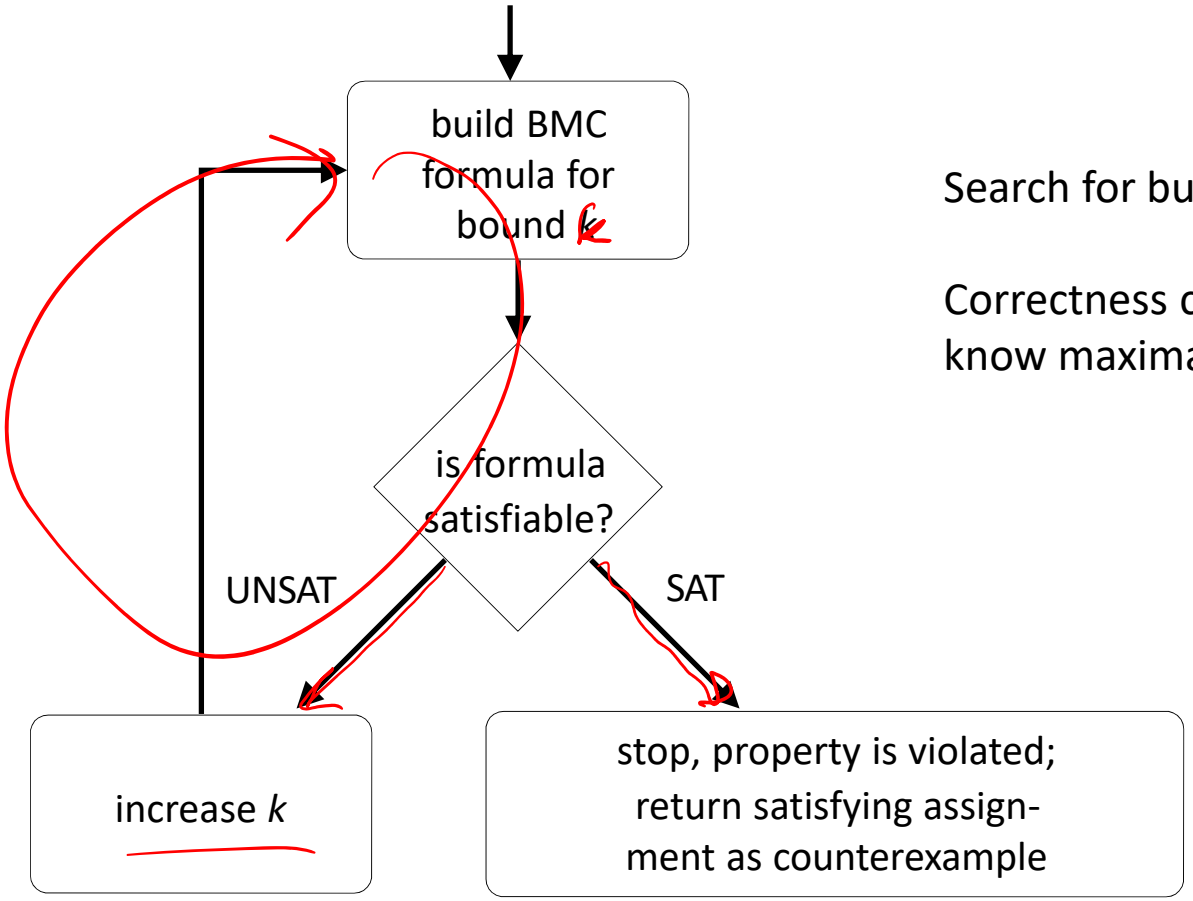
Computer Aided Verification Award 2018

Bounded Model Checking has revolutionized the way model checking is used and perceived. It has increased the capabilities of model checkers by orders of magnitude, turning them into a standard tool for hardware verification and a very important component of the toolkit available for software verification…

(1999)

Flavio Lerda, Daniel Kroening, Armin Biere, Alessandro Cimatti,
Not pictured: Edmund M. Clarke, Yunshan Zhu

Bounded Model Checking



Search for bugs within k steps.

Correctness can only be proven if we know maximal value for *k*

Reachability Properties

Property

p always holds iff we cannot reach a state with $\neg p$

Kripke Structure

$M = (S, S_0, R, AP, L)$ – represented symbolically (See Chapter 3)

State variables $\mathrm{V} = \{v_1, \ldots, v_n\}$

Reachability – Paths

$$path_0(s_0) = S_0(s_0)$$

$$path_1(s_0, s_1) = S_0(s_0) \wedge R(s_0, s_1)$$

$$path_2(s_0, \ldots, s_2) = S_0(s_0) \wedge R(s_0, s_1) \wedge R(s_1, s_2)$$

$$\vdots$$

$$path_k(s_0, \ldots, s_k) = S_0(s_0) \wedge \bigwedge_{i=0}^{k-1} R(s_i, s_{i+1})$$

Reachability – Paths

$$path_0(s_0) = S_0(s_0)$$
$$path_1(s_0, s_1) = S_0(s_0) \wedge R(s_0, s_1)$$
$$path_2(s_0, \dots, s_2) = S_0(s_0) \wedge R(s_0, s_1) \wedge R(s_1, s_2)$$
$$path_3(s_0, \dots, s_3) = S_0(s_0) \wedge R(s_0, s_1) \wedge R(s_1, s_2) \wedge R(s_1, s_3)$$

$$path_k(s_0, \dots, s_k) = S_0(s_0) \wedge \bigwedge_{i=0}^{k-1} R(s_i, s_{i+1})$$

Reachability – Building the Formula

Reachability – Building the Formula

$$path_k(s_0, \ldots, s_k) = S_0(s_0) \wedge \bigwedge_{i=0}^{k-1} R(s_i, s_{i+1})$$

Path starts in initial state

Exists transition from s_i to s_{i+1}

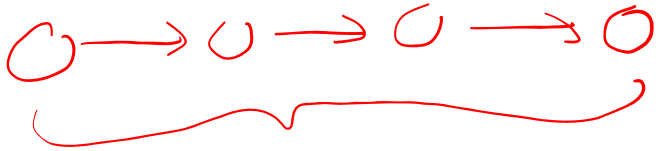
spec: p hold in all states

System is **incorrect** within k steps if

$$path_k(s_0, \ldots, s_k) \wedge \bigvee_{i=0}^{k} \neg p(s_i)$$

There is a path to s_k

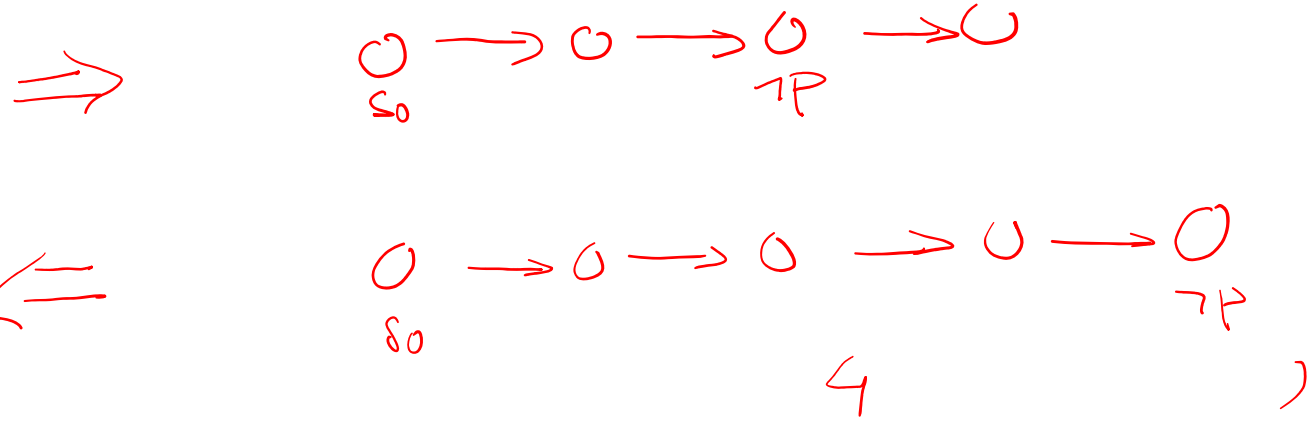
One of the state violates p



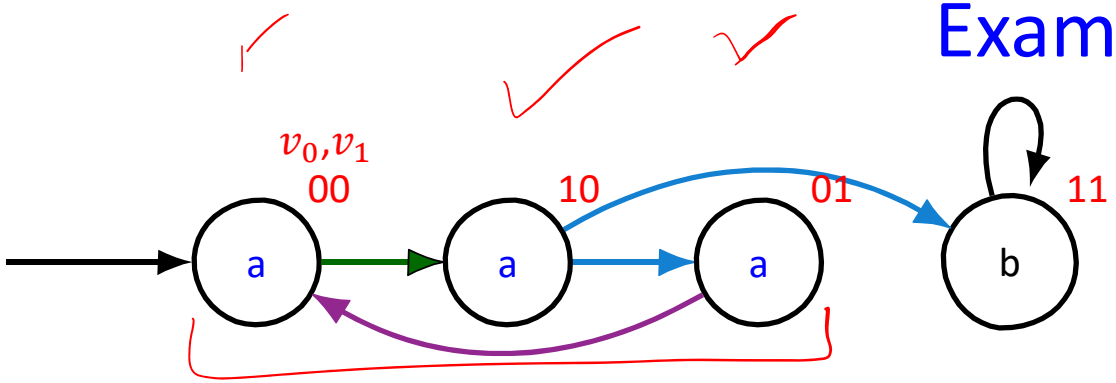
Reachability – Correctness

Theorem 10.1. $path_{k(s_0,\ldots,s_k)} \wedge \bigvee_{i=0}^{k} \neg p(s_i)$ is satisfiable iff there is a counterexample to $\mathrm{AG}\, p$ of length $\leq k$.

~~Proof uses fact that R is left total.~~



Example



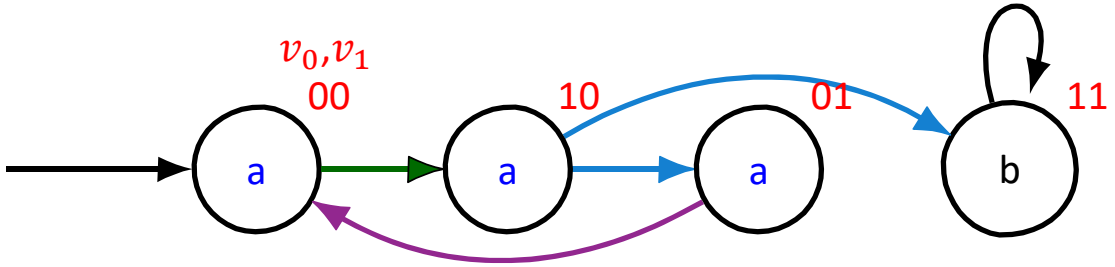
AG a a true in all reachable states.

$S_0(v_0, v_1) =.$

$R(v_0, v_1, v_0', v_1') =$

$a(v_0, v_1) =.$ $\neg v_0 \vee \neg v_1$

Example



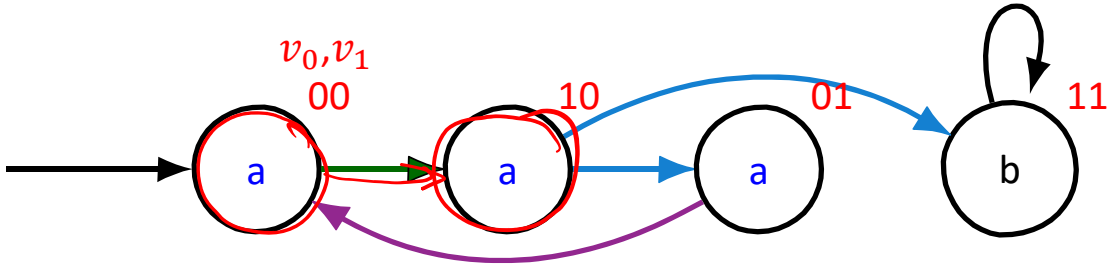
AG a

$$S_0(v_0, v_1) = \neg v_0 \wedge \neg v_1$$

$$\begin{aligned} R(v_0, v_1, v_0', v_1') = {} & \neg v_0 \wedge \neg v_1 \wedge v_0' \wedge \neg v_1' \\ & \vee v_0 \wedge \neg v_1 \wedge v_1' \\ & \vee \neg v_0 \wedge v_1 \wedge \neg v_0' \wedge \neg v_1' \\ & \vee v_0 \wedge v_1 \wedge v_0' \wedge v_1' \end{aligned}$$

$$a(v_0, v_1) = \neg v_0 \vee \neg v_1$$

Example



$$S_0(v_0, v_1) = \neg v_0 \wedge \neg v_1$$

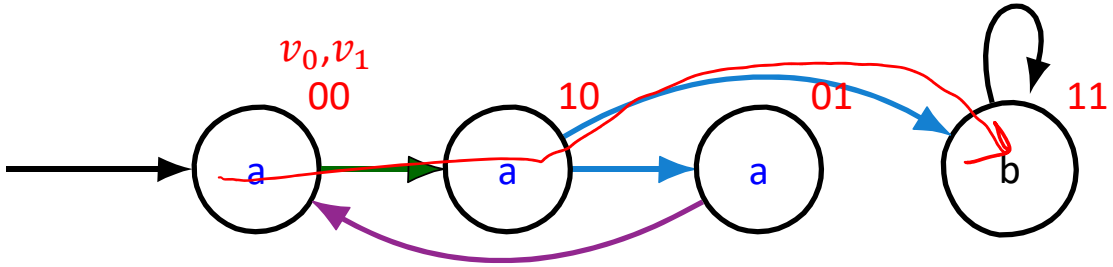
$$\begin{aligned} R(v_0, v_1, v_0', v_1') = {} & \neg v_0 \wedge \neg v_1 \wedge v_0' \wedge \neg v_1' \\ & \vee v_0 \wedge \neg v_1 \wedge v_1' \\ & \vee \neg v_0 \wedge v_1 \wedge \neg v_0' \wedge \neg v_1' \\ & \vee v_0 \wedge v_1 \wedge v_0' \wedge v_1' \end{aligned}$$

$$a(v_0, v_1) = \neg v_0 \vee \neg v_1$$

$$\begin{aligned} path_1(s_0, s_1, s_2) &= S_0(s_0) \wedge R(s_0, s_1) \wedge R(s_1, s_2) \\ &= \neg v_{0,0} \wedge \neg v_{1,0} \wedge \\ & \begin{pmatrix} \neg v_{00} \wedge \neg v_{10} \wedge v_{01} \wedge \neg v_{11} \\ \vee\, v_{00} \wedge \neg v_{10} \wedge v_{11} \\ \vee \neg v_{00} \wedge v_{10} \wedge \neg v_{01} \wedge \neg v_{11} \\ \vee\, v_{00} \wedge v_{10} \wedge v_{01} \wedge v_{11} \end{pmatrix} \end{aligned}$$

$$\bigvee_{i=0}^{k} \neg a(s_i) = \neg(\neg v_{00} \vee \neg v_{10}) \vee \neg(\neg v_{01} \vee \neg v_{11})$$

Example



$$S_0(v_0, v_1) = \neg v_0 \wedge \neg v_1$$

$$R(v_0, v_1, v_0', v_1') = \neg v_0 \wedge \neg v_1 \wedge v_0' \wedge \neg v_1'$$
$$\vee\ v_0 \wedge \neg v_1 \wedge v_1'$$
$$\vee\ \neg v_0 \wedge v_1 \wedge \neg v_0' \wedge \neg v_1'$$
$$\vee\ v_0 \wedge v_1 \wedge v_0' \wedge v_1'$$

$$a(v_0, v_1) = \neg v_0 \vee \neg v_1$$

$$path_2(s_0, s_1, s_2) = S_0(s_0) \wedge R(s_0, s_1) \wedge R(s_1, s_2)$$

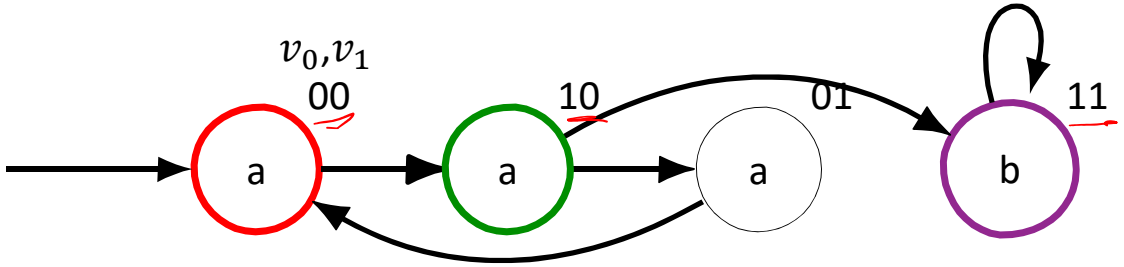
$$= \neg v_{0,0} \wedge \neg v_{1,0} \wedge$$

$$\begin{pmatrix} \neg v_{00} \wedge \neg v_{10} \wedge v_{01} \wedge \neg v_{11} \\ \vee\ v_{00} \wedge \neg v_{10} \wedge v_{11} \\ \vee\ \neg v_{00} \wedge v_{10} \wedge \neg v_{01} \wedge \neg v_{11} \\ \vee\ v_{00} \wedge v_{10} \wedge v_{01} \wedge v_{11} \end{pmatrix} \wedge$$

$$\begin{pmatrix} \neg v_{01} \wedge \neg v_{11} \wedge v_{02} \wedge \neg v_{12} \\ \vee\ v_{01} \wedge \neg v_{11} \wedge v_{12} \\ \vee\ \neg v_{01} \wedge v_{11} \wedge \neg v_{02} \wedge \neg v_{12} \\ \vee\ v_{01} \wedge v_{11} \wedge v_{02} \wedge v_{12} \end{pmatrix}$$

$$\bigvee_{i=0}^{k} \neg a(s_i) = \neg(\neg v_{00} \vee \neg v_{10}) \vee \neg(\neg v_{01} \vee \neg v_{11}) \vee \neg(\neg v_{02} \vee \neg v_{12})$$

Example



AG a

$$path_2(s_0, s_1, s_2) = S_0(s_0) \wedge R(s_0, s_1) \wedge R(s_1, s_2)$$

$$= \neg v_{0,0} \wedge \neg v_{1,0} \wedge \begin{pmatrix} \neg v_{00} \wedge \neg v_{10} \wedge v_{01} \wedge \neg v_{11} \\ \vee\, v_{00} \wedge \neg v_{10} \wedge v_{11} \\ \vee\, \neg v_{00} \wedge v_{10} \wedge \neg v_{01} \wedge \neg v_{11} \\ \vee\, v_{00} \wedge v_{10} \wedge v_{01} \wedge v_{11} \end{pmatrix} \wedge \begin{pmatrix} \neg v_{01} \wedge \neg v_{11} \wedge v_{02} \wedge \neg v_{12} \\ \vee\, v_{01} \wedge \neg v_{11} \wedge v_{12} \\ \vee\, \neg v_{01} \wedge v_{11} \wedge \neg v_{02} \wedge \neg v_{12} \\ \vee\, v_{01} \wedge v_{11} \wedge v_{02} \wedge v_{12} \end{pmatrix}$$

$$\bigvee_{i=0}^{k} \neg a(s_i) = \neg(\neg v_{00} \vee \neg v_{10}) \vee \neg(\neg v_{01} \vee \neg v_{11}) \vee \neg(\neg v_{02} \vee \neg v_{12})$$

Satisfying assignment

i	0	1	2
v_{0i}	0	1	1
v_{1i}	0	0	1

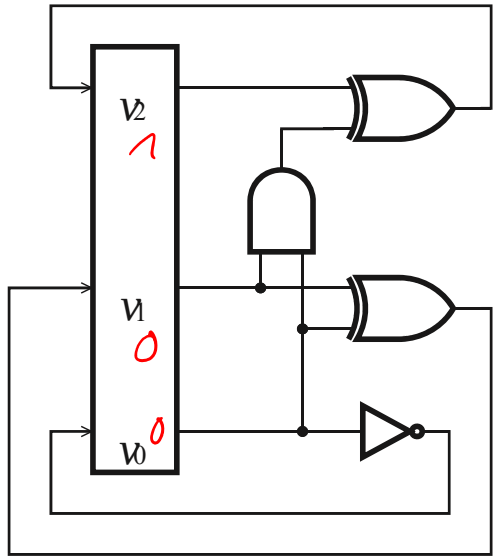
Try it with Z3!

```
(declare-const v00 Bool)
(declare-const v10 Bool)
(declare-const v01 Bool)
(declare-const v11 Bool)
(declare-const v02 Bool)
(declare-const v12 Bool)

(define-fun S0 ((v0 Bool) (v1 Bool)) Bool
  (and (not v0) (not v1)))

(define-fun R ((v0 Bool) (v1 Bool) (w0 Bool) (w1 Bool))
Bool
  (or
      (and (not v0) (not v1) w0 (not w1))
      (and v0 (not v1) w1)
      (and (not v0) v1 (not w0) (not w1))
      (and v0 v1 w0 w1))
)

(define-fun a ((v0 Bool) (v1 Bool)) Bool
   (or (not v0) (not v1))
)

(define-fun path1 ((v00 Bool) (v10 Bool) (v01 Bool) (v11
Bool)) Bool
  (and (S0 v00 v10)
    (R v00 v10 v01 v11)
  )
)
```

```
(define-fun path2 ((v00 Bool) (v10 Bool) (v01 Bool) (v11
Bool) (v02 Bool) (v12 Bool)) Bool
  (and (S0 v00 v10)
    (R v00 v10 v01 v11)
    (R v01 v11 v02 v12)
  )
)

; k = 1
;(assert
;  (and (path1 v00 v10 v01 v11)
;    (or (not (a v00 v10))
;      (not (a v01 v11))
;    )
;  )
;)

; k =2
(assert
  (and (path2 v00 v10 v01 v11 v02 v12)
    (or (not (a v00 v10))
      (not (a v01 v11))
      (not (a v02 v12))
    )
  )
)

(check-sat)
(get-model)
```

https://rise4fun.com/Z3

Another Example



Does modulo-8 counter of Chapter 3.5 ever reach 4?

$\neg p = \neg v_0 \wedge \neg v_1 \wedge v_2.$

$S_0(V) = \neg v_0 \wedge \neg v_1 \wedge \neg v_2.$

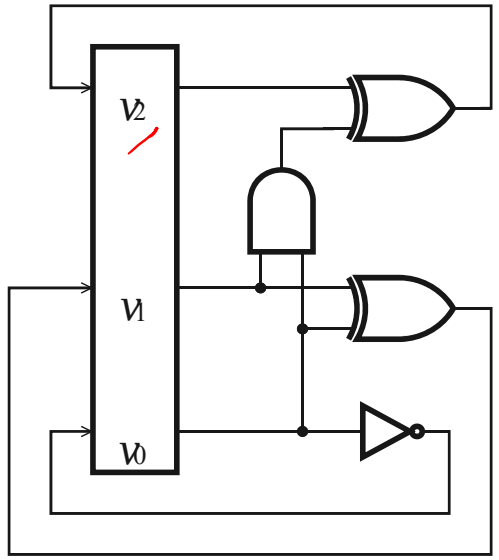
$R(V, V') = (v_0' \leftrightarrow \neg v_0) \wedge (v_1' \leftrightarrow v_0 \oplus v_1) \wedge (v_2' \leftrightarrow (v_0 \wedge v_1) \oplus v_2)$.

k =0:

$S_0(v)$ $\neg v_0 \wedge \neg v_1 \wedge \neg v_2 \wedge$

\neg p(v) $\neg v_0 \wedge \neg v_1 \wedge v_2.$

Another Example



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$R(V, V') = (v'_0 \leftrightarrow \neg v_0) \wedge (v'_1 \leftrightarrow v_0 \oplus v_1) \wedge (v'_2 \leftrightarrow (v_0 \wedge v_1) \oplus v_2)$.

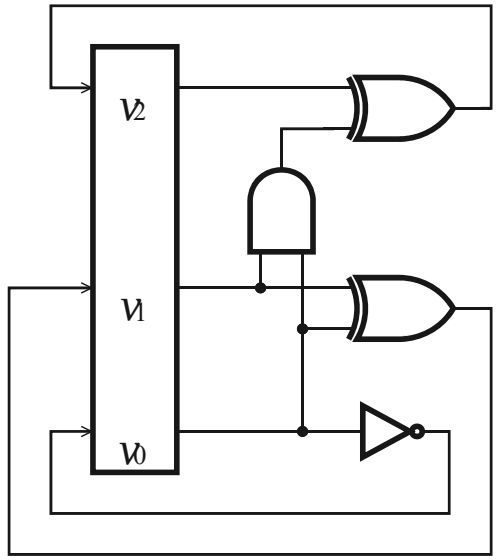
$k = 1$:

$S_0(v)$ $\neg v_0 \wedge \neg v_1 \wedge \neg v_2 \wedge$

$R(v, v')$ $(v'_0 \leftrightarrow \neg v_0) \wedge (v'_1 \leftrightarrow v_0 \oplus v_1) \wedge (v'_2 \leftrightarrow (v_0 \wedge v_1) \oplus v_2) \wedge$

$(\underbrace{\neg v_0 \wedge \neg v_1 \wedge v_2}_{\neg p(v)} \vee \underbrace{\neg v'_0 \wedge \neg v'_1 \wedge v'_2}_{\neg p(v')}).$

Another Example



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$R(V, V') = (v'_0 \leftrightarrow \neg v_0) \wedge (v'_1 \leftrightarrow v_0 \oplus v_1) \wedge (v'_2 \leftrightarrow (v_0 \wedge v_1) \oplus v_2)$.

$k = 2$:

$S_0(v)$ $\neg v_0 \wedge \neg v_1 \wedge \neg v_2 \wedge$

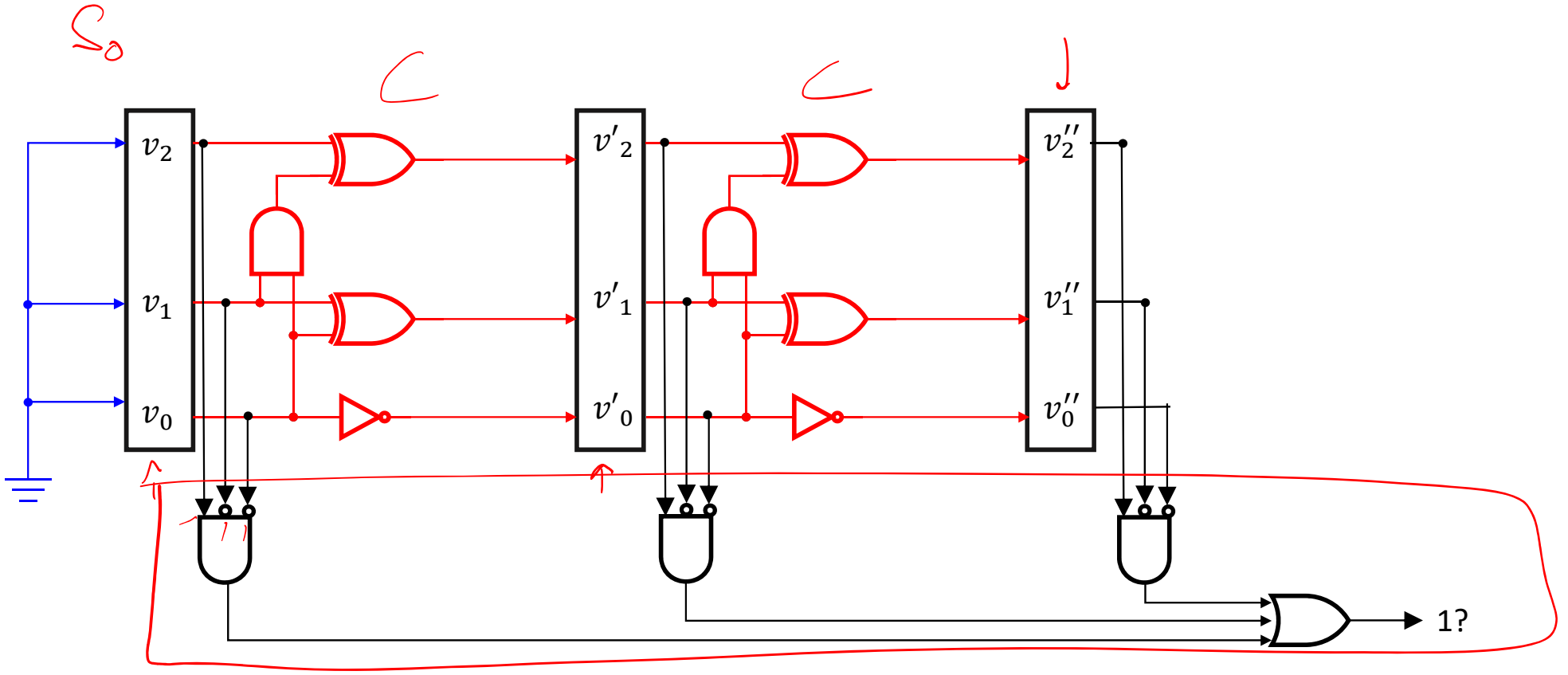
$R(v, v')$ $(v'_0 \leftrightarrow \neg v_0) \wedge (v'_1 \leftrightarrow v_0 \oplus v_1) \wedge (v'_2 \leftrightarrow (v_0 \wedge v_1) \oplus v_2) \wedge$

$R(v', v'')$ $(v''_0 \leftrightarrow \neg v'_0) \wedge (v''_1 \leftrightarrow v'_0 \oplus v'_1) \wedge (v''_2 \leftrightarrow (v'_0 \wedge v'_1) \oplus v'_2) \wedge$

$(\underbrace{\neg v_0 \wedge \neg v_1 \wedge v_2}_{\neg p(v)} \vee \underbrace{\neg v'_0 \wedge \neg v'_1 \wedge v'_2}_{\neg p(v')} \vee \underbrace{\neg v''_0 \wedge \neg v''_1 \wedge v''_2}_{\neg p(v'')}.)$

Model Checking

The Electrical Engineer's Viewpoint



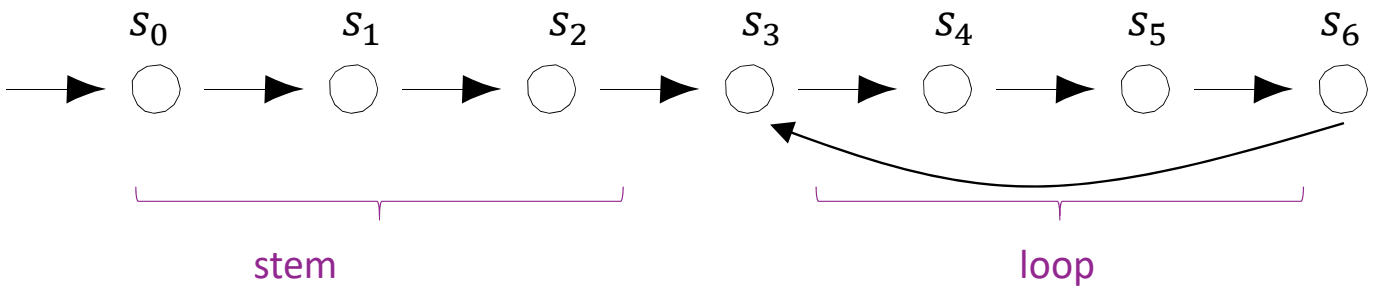
Eventuality Properties

Property

Suppose $\phi = \mathbf{AF}\ p$

Counterexamples fulfill $\mathbf{EG}\ \neg p$

Counterexamples have *Lasso Shape* (See Chapter 4.)



Eventuality Properties

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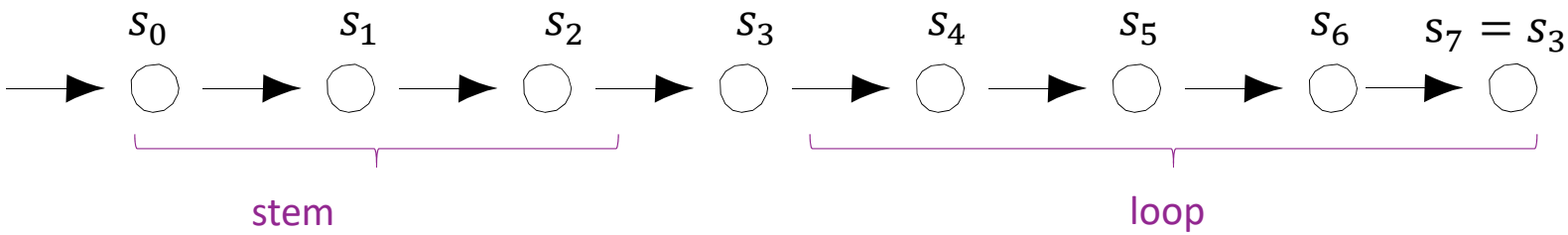


$$lasso_k(s_0, \ldots, s_k) = path_k(s_0, \ldots, s_k) \wedge \bigvee_{i=0}^{k-1} s_i = s_k$$

Eventuality Properties

Property

Suppose $\phi = \mathbf{AF}\, p$

Counterexamples fulfill $\mathbf{EG}\, \neg p$

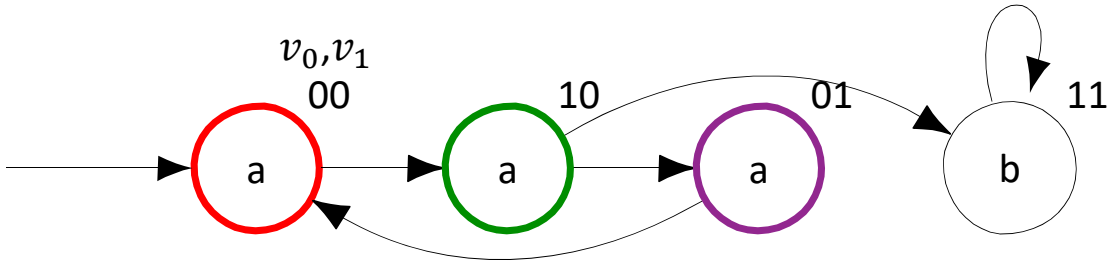
Counterexamples have *Lasso Shape* (See Chapter 4.)

$$lasso_k(s_0, \ldots, s_k) = path_k(s_0, \ldots, s_k) \wedge \bigvee_{i=0}^{k-1} s_i = s_k$$

Counterexample lasso:

$$lasso_k(s_0, \ldots, s_k) \wedge \bigwedge_{i=0}^{k-1} \neg p$$

Example



AF b

$$lasso_3(s_0, \ldots, s_3) =$$

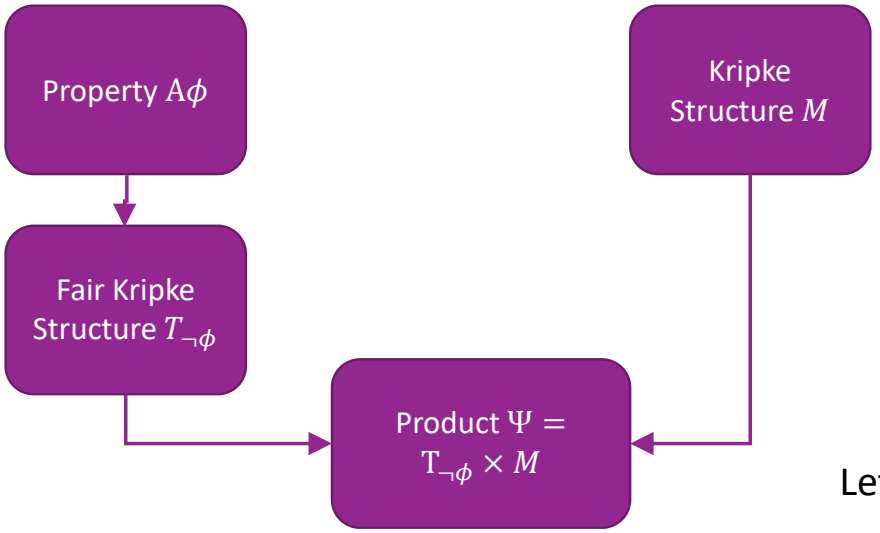
$$path_3(s_0, \ldots, s_3) \wedge \bigvee_{i=0}^{2} s_i = s_3 \wedge \bigwedge_{i=0}^{2} \neg b$$

Satisfying assignment

i	0	1	2	3
v_{0i}	0	1	1	0
v_{1i}	0	0	1	0

Full LTL

See Chapter 6



Let $\Psi = (S^{\Psi}, S_0^{\Psi}, R^{\Psi}, AP^{\Psi}, L^{\Psi}, F^{\Psi})$

$$S_0^{\Psi}(s_0) \wedge \bigwedge_{i=0}^{k-1} R^{\Psi}(s_i, s_{i+1}) \wedge \bigvee_{i=0}^{k-1} (s_i = s_k \wedge Fair_i^P)$$

Satisfiable iff there is a fair lasso of length k

Completeness

BMC finds bugs of length $\leq k$.

Longer counterexamples may exist!

Is there a k big enough to exclude any counterexample?

Def. $M \models_k \phi$: all computations of length k satisfy ϕ.

Completeness threshold: Number CT such that $M \models_{CT} \phi \Rightarrow M \models \phi$

If completeness threshold known, stop BMC when $k = CT$

Ideas for finding CT?

- number of states $\approx 2^n$ states for n flipflops

Completeness Threshold

Finding smallest CT is as difficult as model checking!

- Smallest CT is size of shortest counterexample, or 0 if the property is satisfied

Simple values for CT

- Number of state of M is bound on CT
- Diameter (longest simple path between two states) is bound on CT

These are typically really large numbers

Verifying Reachability Properties with k-induction



Mary Sheeran, Koen Claessen, Per Bjesse, 2000

Motivation

- Completeness thresholds usually very large
- Can we **prove** a property with fewer unrollings?
- Idea: Use induction.

Base: Prove $Q(0)$

Induction: Prove $Q(n-1) \Rightarrow Q(n)$

Conclusion: $\forall n. Q(n)$

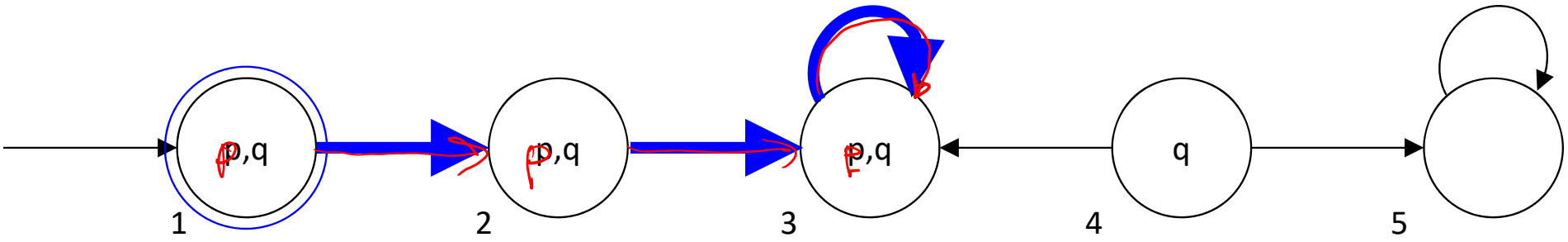
Caveat: Property may be true, but not inductive (see below)

Induction

Let's prove **AG** p on the following structure.

Take arbitrary path π starting in q_1.

- **Base case:** $\pi(0) \vDash p$ true: $q_1 \vDash p$ true: $q_1 \vDash p$
- **Induction:** if $\pi(n-1) \vDash p$ then $\pi(n) \vDash p$ true: any successor of a p-state is a p-state
- **Conclusion**: for any path π we have $\forall n.\, \pi(n) \vDash p$



Satisfiability

Let's prove $\mathrm{AG}\, p$ on the following structure. How can these properties be violated?

Take arbitrary path π

- **Base case:** $\pi(0) \vDash p$
- **Induction:** if $\pi(n-1) \vDash p$ then $\pi(n) \vDash p$
- **Conclusion**: for any path π we have $\forall n.\, \pi(n) \vDash p$

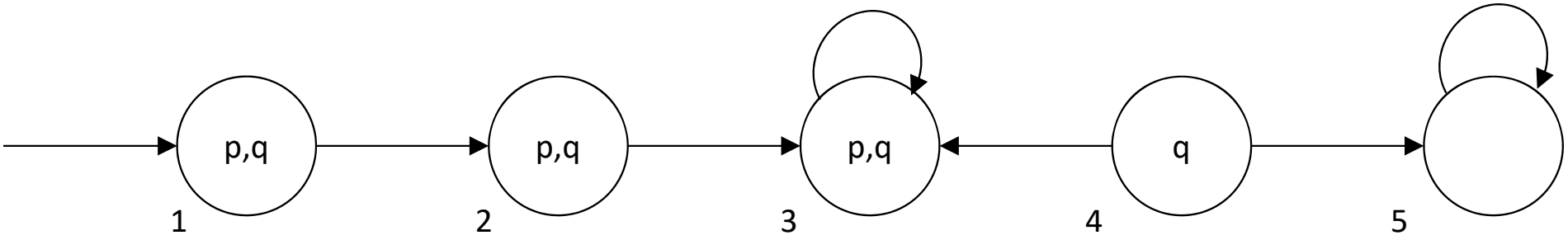
$S_0(s) \wedge \neg p(s)$ Unsatisfiable

$p(s) \wedge R(s,s') \wedge \neg p(s')$ Unsatisfiable

UNSAT

UNSAT

System correct!

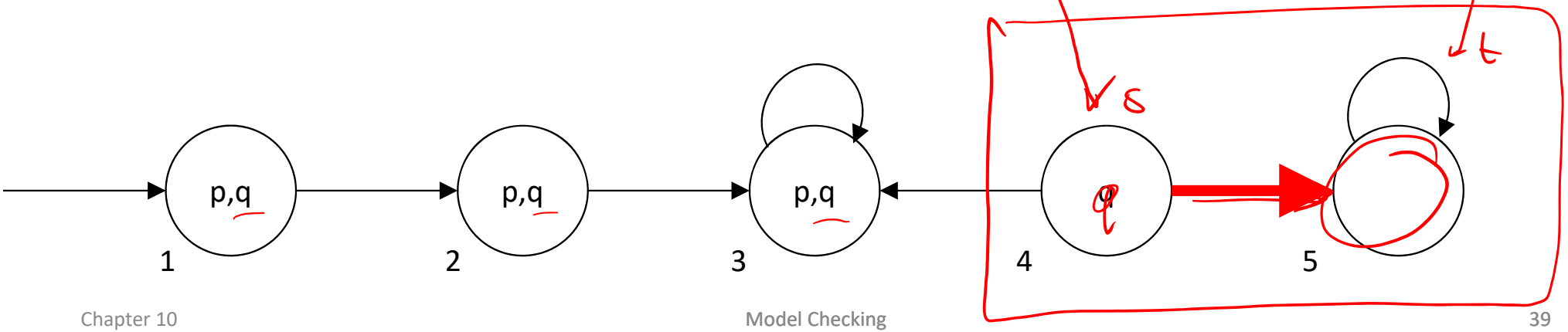


A Problem

Let's prove $\mathbf{AG}\ q$ on the following structure.

Take arbitrary path π

- **Base case:** $\pi(0) \vDash q$
- **Induction:** if $\pi(n-1) \vDash q$ then $\pi(n) \vDash q$ **not true!**
- ~~**Conclusion:** for any path π we have $\forall n.\ \pi(n) \vDash q$~~ **not all true properties are inductive**



***k*-induction**

$q(s_0) \wedge R(s_0, s_1) \wedge q(s_1)$

$q(n-2) \wedge q(n-1) \rightarrow q(n)$

Base:

Induction:

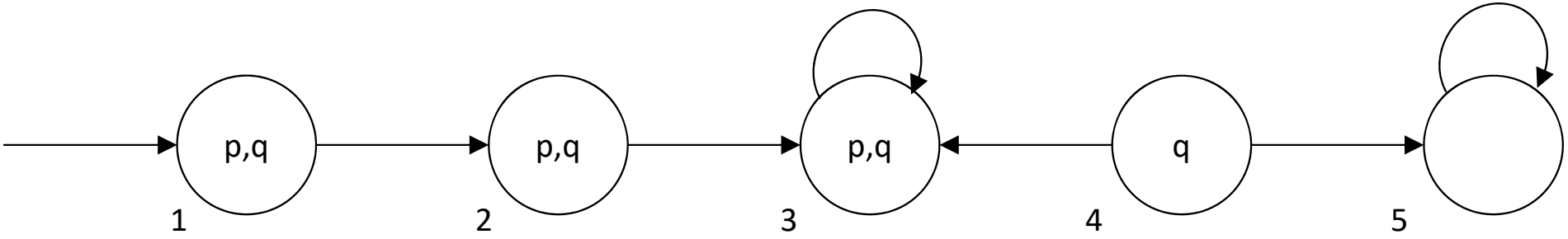
Conclusion: $\forall n. Q(n)$

In our setting:

Base. all paths of length *k* from S_0 are labeled *q*

Induction. all paths of length *k* labeled with all *q*s are followed by a *q*

Conclusion. All paths from S_0 are labeled *q*



k-induction

Base: Prove $Q(0) \wedge \cdots \wedge Q(k-1)$

Induction: Prove $Q(n-k) \wedge \cdots \wedge Q(n-1) \Rightarrow Q(n)$

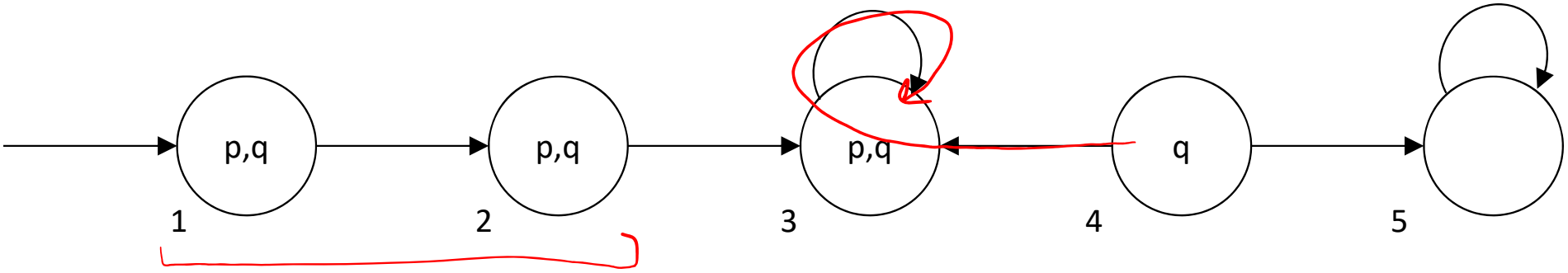
Conclusion: $\forall n.\, Q(n)$

In our setting:

Base. all paths of length *k* from S_0 are labeled *q*

Induction. all paths of length *k* labeled with all *q*s are followed by a *q*

Conclusion. All paths from S_0 are labeled *q*



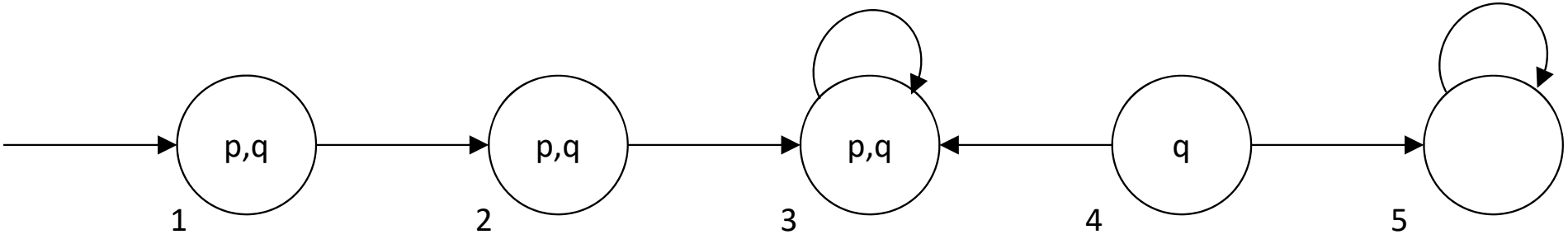
Prove AG q using 2-induction

Base: Consider all paths of length 2 from q_1: $q_1 \vDash q$ and $q_2 \vDash q$.

Induction: Do all successors of paths of length 2 labeled q, q fulfill q?

- (q_1, q_2)
- (q_2, q_3)
- (q_3, q_4)
- (q_4, q_4)

Conclusion: for any path π we have $\forall n. \pi(n) \vDash p$



k-induction as Satisfiability

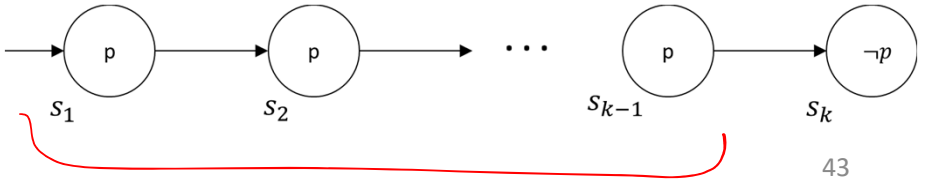
Base. all paths of length *k* from S_0 are labeled *p*

This is BMC! $S_0(s_0) \wedge \Lambda_{i=0}^{k-1} R(s_i, s_{i+1}) \wedge \vee_{i=0}^{k} \neg p(s_i)$

Induction. all paths of length *k* labeled with all *p*'s are followed by a *p*

$$\bigwedge_{i=0}^{k-1} R(s_i, s_{i+1}) \wedge \bigwedge_{i=0}^{k-1} p(s_i) \wedge \neg p(s_k)$$

Formula satisfiable iff there is a counterexample



k-induction

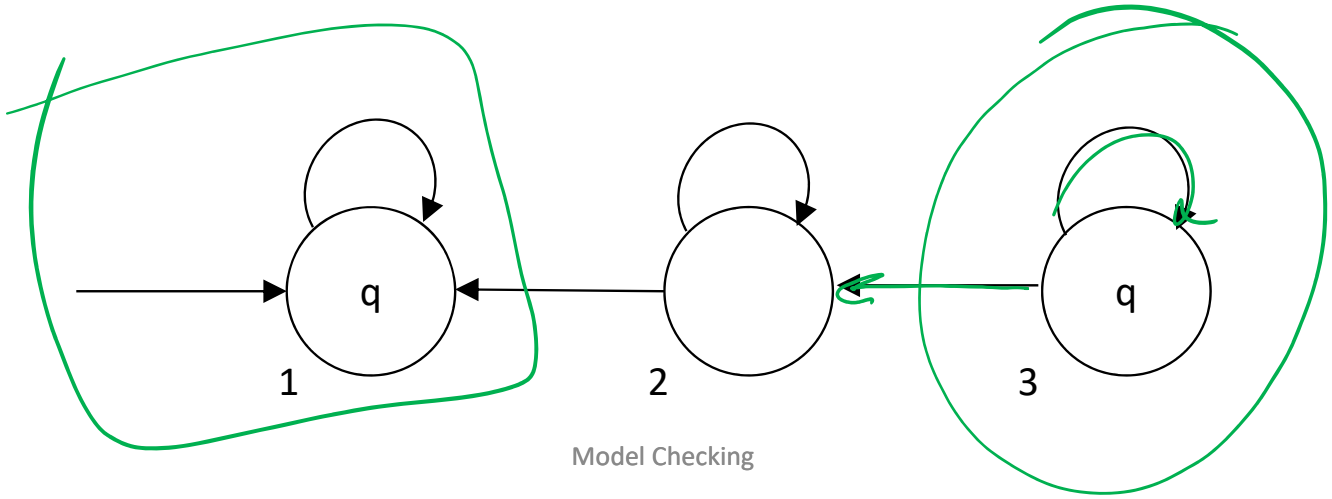
```
while(k=0; ; k++){
        build BMC formula φ
        if φ SAT return "bug!"

        build k-induction formula ψ
         if φ UNSAT return "correct!"
}
```

k-induction is not Complete

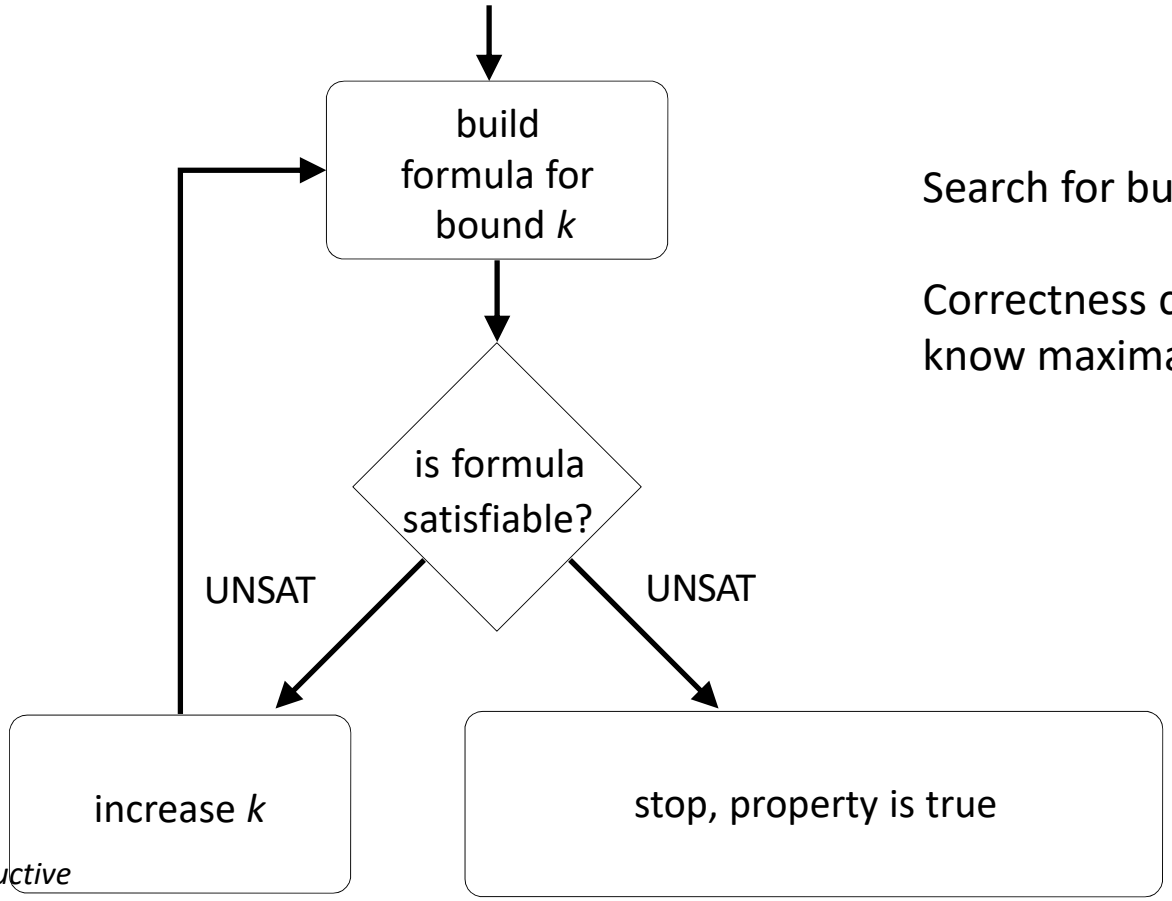
System satisfies **AG** q, but induction step fails for any *k*

Base. all paths of length *k* from S_0 are labeled *q*

Induction. all paths of length *k* labeled with all *q*s are followed by a *q*. *FALSE*!



k-induction



Search for bugs within k steps.

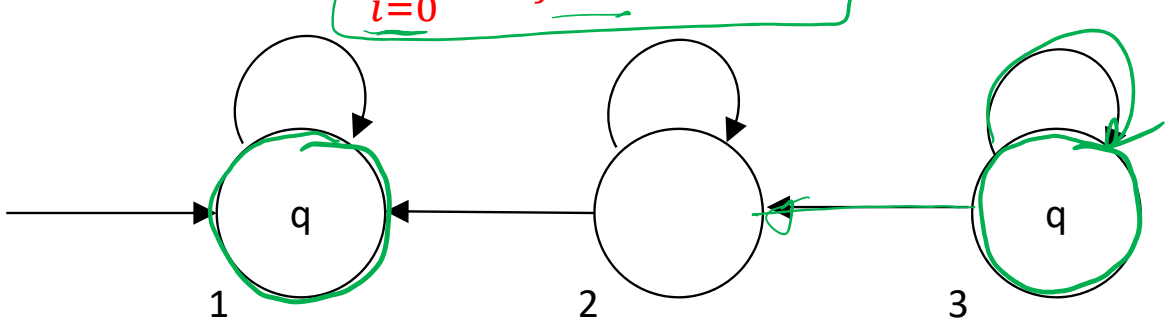
Correctness can only be proven if we know maximal value for *k*

Property not k-inductive
May or may not be true

Making k-induction Complete

Induction. all noncyclic paths of length *k* labeled with all *q*s are followed by a *q*

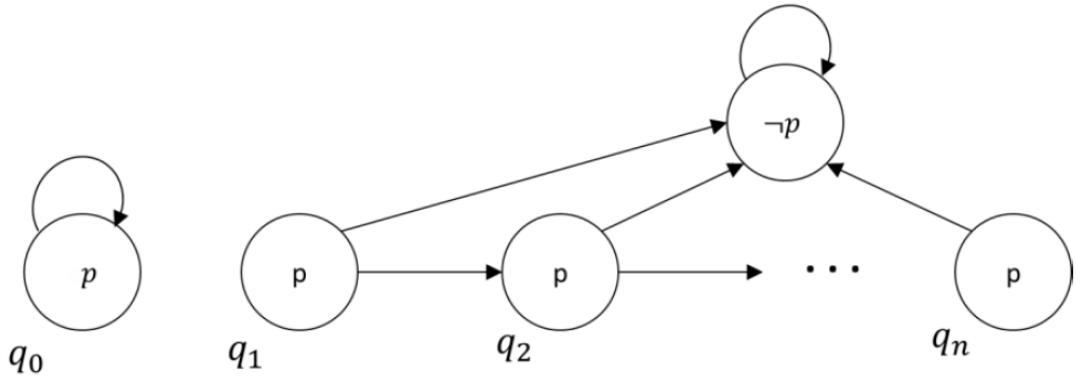
$$\bigwedge_{i=0}^{k-1} R(s_i, s_{i+1}) \wedge \bigwedge_{i=0}^{k-1} p(s_i) \wedge \neg p(s_k) \wedge \bigwedge_{i=0}^{k-1} \bigwedge_{j=i+1}^{k} s_i \neq s_j$$



thus can't prove with ~~k=2~~ k=1.

HW2

Consider the following synchronous Kripke structure K.



We wish to prove that *p* is always true.

Task 2a. [5 points]

Suppose that q_1 is the initial state. Suppose you are given formulas R, S_0, and p for the transition relation, the initial states and the property, resp.

- What is the smallest *k* such that BMC finds a counterexample?
- Show the BMC formula, using R, S_0, and p.
- Is the formula satisfiable? Explain.

Task 2c. [5 points]

Suppose that q_0 is the initial state. The new formula for the initial states is S'_0.

- What is the smallest *k* such that k-induction can prove the property correct?
- Suppose *n=2.* Choose an appropriate *k* and show the k-induction formula, using R, S'_0, and p.
- Is the formula satisfiable? Explain.